## The Probabilistic Method

Techniques

- Union bound
- Argument from expectation
- Alterations
- The second moment method
- The (Lovasz) Local Lemma

And much more

- Alon and Spencer, "The Probabilistic Method"
- Bolobas, "Random Graphs"


## The Argument from Expectation: Main Idea

- $X$ a random variable with $\mathrm{E}[X]=\mu$, then
- There must exist a sample point $\omega$ with $X(\omega) \geq \mu$
- There must exist a sample point $\omega$ with $X(\omega) \leq \mu$
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Have we seen this?

## Example 1: Large Cuts in Graphs

## Intuition \& Question

Intuition: every graph must have a "sufficiently large" cut $(A, B)$.
Question: How large?

## Line of thought

On average, a random cut has size $\mu$, hence there must exist a cut of size $\geq \mu$.

- Put a vertex in either $A$ or $B$ with probability $1 / 2$
- Expected number of edges $X$ with one end point in each is

$$
\mathrm{E}[X]=\mathrm{E}\left[\sum_{e} X_{e}\right]=\sum_{e} \operatorname{Prob}\left[X_{e}\right]=|E| / 2
$$

Theorem
For every graph $G=(V, E)$, there must be a cut with $\geq|E| / 2$ edges

## Example 2: $\pm 1$ Linear Combinations of Unit Vectors

## Theorem

Let $\mathbf{v}_{1}, \cdots, \mathbf{v}_{n}$ be $n$ unit vectors in $\mathbb{R}^{n}$.
There exist $\alpha_{1}, \cdots, \alpha_{n} \in\{-1,1\}$ such that

$$
\left|\alpha_{1} \mathbf{v}_{1}+\cdots+\alpha_{n} \mathbf{v}_{n}\right| \leq \sqrt{n}
$$

and, there exist $\alpha_{1}, \cdots, \alpha_{n} \in\{-1,1\}$ such that

$$
\left|\alpha_{1} \mathbf{v}_{1}+\cdots+\alpha_{n} \mathbf{v}_{n}\right| \geq \sqrt{n}
$$

Simply because on average these combinations have length $\sqrt{n}$. Specifically, choose $\alpha_{i} \in\{-1,1\}$ independently with prob. $1 / 2$

$$
\mathrm{E}\left[\left|\alpha_{1} \mathbf{v}_{1}+\cdots+\alpha_{n} \mathbf{v}_{n}\right|^{2}\right]=\sum_{i, j} \mathbf{v}_{i} \cdot \mathbf{v}_{j} \mathrm{E}\left[\alpha_{i} \alpha_{j}\right]=\sum_{i} \mathbf{v}_{i}^{2}=n .
$$

## Example 3: Unbalancing Lights

## Theorem

For $1 \leq i, j \leq n$, we are given $a_{i j} \in\{-1,1\}$. Then, there exist $\alpha_{i}, \beta_{j} \in\{-1,1\}$ such that

$$
\sum_{i} \sum_{j} a_{i j} \alpha_{i} \beta_{j} \geq\left(\sqrt{\frac{2}{\pi}}+o(1)\right) n^{3 / 2}
$$

- Choose $\beta_{j} \in\{-1,1\}$ independently with prob. $1 / 2$.
- $R_{i}=\sum_{j} a_{i j} \beta_{j}$, then

$$
\mathrm{E}\left[\left|R_{i}\right|\right]=2 \frac{n\left(\left\lfloor\binom{n-1}{\lfloor n-1) / 2\rfloor}\right.\right.}{2^{n}} \approx\left(\sqrt{\frac{2}{\pi}}+o(1)\right) n^{1 / 2}
$$

- Choose $\alpha_{i}$ with the same sign as $R_{i}$, for all $i$

