Example 1: Probabilistic Packet Marking (PPM)

The Setting

- A stream of packets are sent $S = R_0 \to R_1 \to \cdots \to R_{n-1} \to D$
- Each R_i can overwrite the SOURCE IP field F of a packet
- D wants to know the set of routers on the route

The Assumption

- For each packet D receives and each i, Prob$[F = R_i] = 1/n$ (*)

The Questions

1. How does the routers ensure (*)?
2. How many packets must D receive to know all routers?

Coupon Collector Problem

The setting

- n types of coupons
- Every cereal box has a coupon
- For each box B and each coupon type t,

$$\text{Prob}\left[B \text{ contains coupon type } t\right] = \frac{1}{n}$$

Coupon Collector Problem

How many boxes of cereal must the collector purchase before he has all types of coupons?

The Analysis

- $X =$ number of boxes he buys to have all coupon types.
- For $i \in [n]$, let X_i be the additional number of cereal boxes he buys to get a new coupon type, after he had collected $i-1$ different types

$$X = X_1 + X_2 + \cdots + X_n, \quad \mathsf{E}[X] = \sum_{i=1}^{n} E[X_i]$$

- After $i-1$ types collected,

$$\text{Prob[A new box contains a new type]} = p_i = 1 - \frac{i-1}{n}$$

- Hence, X_i is *geometric* with parameter p_i, implying

$$\mathsf{E}[X_i] = \frac{1}{p_i} = \frac{n}{n-i+1}$$

$$\mathsf{E}[X] = n \sum_{i=1}^{n} \frac{1}{n-i+1} = nH_n = n \ln n + \Theta(n)$$

PTCF: Geometric Distribution

- A coin turns head with probability p, tail with $1-p$
- $X =$ number of flips until a head shows up
- X has geometric distribution with parameter p

$$
\begin{aligned}
\text{Prob}[X=n] &= (1-p)^{n-1}p \\
\text{E}[X] &= \frac{1}{p} \\
\text{Var}[X] &= \frac{1-p}{p^2}
\end{aligned}
$$

Additional Questions

- We can't be sure that buying nH_n cereal boxes suffices
- Want Prob$[X \geq C]$, i.e. *what's the probability that he has to buy C boxes to collect all coupon types?*
- Intuitively, X is far from its mean with small probability
- Want something like

$$\text{Prob}[X \geq C] \leq \text{some function of } C\text{, preferably} \ll 1$$

i.e. (large) deviation inequality or tail inequalities

Central Theme

The more we know about X, the better the deviation inequality we can derive: Markov, Chebyshev, Chernoff, etc.

PTCF: Markov's Inequality

Theorem

If X is a r.v. taking only non-negative values, $\mu = \mathsf{E}[X]$, then $\forall a > 0$

$$\mathsf{Prob}[X \geq a] \leq \frac{\mu}{a}.$$

Equivalently,

$$\mathsf{Prob}[X \geq a\mu] \leq \frac{1}{a}.$$

If we know $\mathsf{Var}[X]$, we can do better!

PTCF: (Co)Variance, Moments, Their Properties

- Variance: $\sigma^2 = \mathsf{Var}[X] := \mathsf{E}[(X - \mathsf{E}[X])^2] = \mathsf{E}[X^2] - (\mathsf{E}[X])^2$
- Standard deviation: $\sigma := \sqrt{\mathsf{Var}[X]}$
- kth moment: $\mathsf{E}[X^k]$
- Covariance: $\mathsf{Cov}[X, Y] := \mathsf{E}[(X - \mathsf{E}[X])(Y - \mathsf{E}[Y])]$
- For any two r.v. X and Y,

$$\mathsf{Var}[X + Y] = \mathsf{Var}[X] + \mathsf{Var}[Y] + 2\,\mathsf{Cov}[X, Y]$$

- If X and Y are independent (define it), then

$$\begin{aligned}
\mathsf{E}[X \cdot Y] &= \mathsf{E}[X] \cdot \mathsf{E}[Y] \\
\mathsf{Cov}[X, Y] &= 0 \\
\mathsf{Var}[X + Y] &= \mathsf{Var}[X] + \mathsf{Var}[Y]
\end{aligned}$$

- In fact, if $X_1, \dots, X_n$ are mutually independent, then

$$\mathsf{Var}\left[\sum_i X_i\right] = \sum_i \mathsf{Var}[X_i]$$

PTCF: Chebyshev's Inequality

Theorem (Two-sided Chebyshev's Inequality)

If X *is a r.v. with mean* μ *and variance* σ^2*, then* $\forall a > 0$*,*

$$\text{Prob}\big[|X-\mu| \geq a\big] \leq \frac{\sigma^2}{a^2} \text{ or, equivalently } \text{Prob}\big[|X-\mu| \geq a\sigma\big] \leq \frac{1}{a^2}.$$

Theorem (One-sided Chebyshev's Inequality)

Let X *be a r.v. with* $\mathsf{E}[X] = \mu$ *and* $\mathsf{Var}[X] = \sigma^2$*, then* $\forall a > 0$*,*

$$\begin{aligned}\text{Prob}[X \geq \mu + a] &\leq \frac{\sigma^2}{\sigma^2 + a^2}\\ \text{Prob}[X \leq \mu - a] &\leq \frac{\sigma^2}{\sigma^2 + a^2}.\end{aligned}$$

Back to the Additional Questions

- Markov's leads to,

$$\text{Prob}[X \geq 2nH_n] \leq \frac{1}{2}$$

- To apply Chebyshev's, we need $\text{Var}[X]$:

$$\text{Prob}[|X - nH_n| \geq nH_n] \leq \frac{\text{Var}[X]}{(nH_n)^2}$$

- Key observation: the X_i are independent (why?)

$$\text{Var}[X] = \sum_i \text{Var}[X_i] = \sum_i \frac{1-p_i}{p_i^2} \leq \sum_i \frac{n^2}{(n-i+1)^2} = \frac{\pi^2 n^2}{6}$$

- Chebyshev's leads to

$$\text{Prob}[|X - nH_n| \geq nH_n] \leq \frac{\pi^2}{6H_n^2} = \Theta\left(\frac{1}{\ln^2 n}\right)$$

Example 2: PPM with One Bit

The Problem

Alice wants to send to Bob a message $b_1 b_2 \cdots b_m$ of m bits. She can send only **one** bit at a time, but always forgets which bits have been sent. Bob knows m, nothing else about the message.

The solution

- Send bits so that the fraction of bits 1 received is within ϵ of $p = B/2^m$, where $B = b_1 b_2 \cdots b_m$ as an integer
- Specifically, send bit 1 with probability p, and 0 with $(1 - p)$

The question

How many bits must be sent so B can be decoded with high probability?

The Analysis

- One way to do decoding: round the fraction of bits 1 received to the closest multiple of of $1/2^m$
- Let $X_1, \dots, X_n$ be the bits received (independent Bernoulli trials)
- Let $X = \sum_i X_i$, then $\mu = \mathsf{E}[X] = np$. We want, say

$$\text{Prob}\left[\left|\frac{X}{n} - p\right| \leq \frac{1}{3 \cdot 2^m}\right] \geq 1 - \epsilon$$

which is equivalent to

$$\text{Prob}\left[|X - \mu| \leq \frac{n}{3 \cdot 2^m}\right] \geq 1 - \epsilon$$

This is a kind of concentration inequality.

PTCF: The Binomial Distribution

- n independent trials are performed, each with success probability p.
- $X =$ number of successes after n trials, then

$$\text{Prob}[X = i] = \binom{n}{i} p^i (1-p)^{n-i}, \ \forall i = 0, \dots, n$$

- X is called a binomial random variable with parameters (n, p).

$$\begin{aligned} \mathsf{E}[X] &= np \\ \mathsf{Var}[X] &= np(1-p) \end{aligned}$$

PTCF: Chernoff Bounds

Theorem (Chernoff bounds are just the following idea)

Let X be any r.v., then

1. *For any $t > 0$*

$$\mathsf{Prob}[X \geq a] \leq \frac{\mathsf{E}[e^{tX}]}{e^{ta}}$$

   In particular,

$$\mathsf{Prob}[X \geq a] \leq \min_{t>0} \frac{\mathsf{E}[e^{tX}]}{e^{ta}}$$

2. *For any $t < 0$*

$$\mathsf{Prob}[X \leq a] \leq \frac{\mathsf{E}[e^{tX}]}{e^{ta}}$$

   In particular,

$$\mathsf{Prob}[X \geq a] \leq \min_{t<0} \frac{\mathsf{E}[e^{tX}]}{e^{ta}}$$

(E^{tX} is called the moment generating function of X)

PTCF: A Chernoff Bound for sum of Poisson Trials

Above the mean case.
Let $X_1, \dots, X_n$ be independent Poisson trials, $\text{Prob}[X_i = 1] = p_i$, $X = \sum_i X_i$, $\mu = \text{E}[X]$. Then,

- For any $\delta > 0$,

$$\text{Prob}[X \geq (1+\delta)\mu] < \left(\frac{e^\delta}{(1+\delta)^{1+\delta}}\right)^\mu;$$

- For any $0 < \delta \leq 1$,

$$\text{Prob}[X \geq (1+\delta)\mu] \leq e^{-\mu\delta^2/3};$$

- For any $R \geq 6\mu$,

$$\text{Prob}[X \geq R] \leq 2^{-R}.$$

Below the mean case.

Let $X_1, \dots, X_n$ be independent Poisson trials, $\mathsf{Prob}[X_i = 1] = p_i$, $X = \sum_i X_i$, $\mu = \mathsf{E}[X]$. Then, for any $0 < \delta < 1$:

1. $$\mathsf{Prob}[X \leq (1-\delta)\mu] \leq \left(\frac{e^{-\delta}}{(1-\delta)^{1-\delta}}\right)^{\mu};$$

2. $$\mathsf{Prob}[X \leq (1-\delta)\mu] \leq e^{-\mu\delta^2/2}.$$

PTCF: A Chernoff Bound for sum of Poisson Trials

A simple (two-sided) deviation case.
Let $X_1, \dots, X_n$ be independent Poisson trials, $\text{Prob}[X_i = 1] = p_i$, $X = \sum_i X_i$, $\mu = \mathsf{E}[X]$. Then, for any $0 < \delta < 1$:

$$\text{Prob}[|X - \mu| \geq \delta\mu] \leq 2e^{-\mu\delta^2/3}.$$

Chernoff Bounds Informally

The probability that the sum of independent Poisson trials is far from the sum's mean is exponentially small.

Back to the 1-bit PPM Problem

$$
\begin{aligned}
\text{Prob}\left[|X-\mu| > \frac{n}{3\cdot 2^m}\right] &= \text{Prob}\left[|X-\mu| > \frac{1}{3\cdot 2^m p}\mu\right] \\
&\leq \frac{2}{\exp\{\frac{n}{18\cdot 4^m p}\}}
\end{aligned}
$$

Now,

$$\frac{2}{\exp\{\frac{n}{18\cdot 4^m p}\}} \leq \epsilon$$

is equivalent to

$$n \geq 18p\ln(2/\epsilon)4^m.$$

Example 3: A Statistical Estimation Problem

The Problem

We want to estimate $\mu = \mathsf{E}[X]$ for some random variable X (e.g., X is the income in dollars of a random person in the world).

The Question

How many samples must be take so that, given $\epsilon, \delta > 0$, the estimated value $\bar{\mu}$ satisfies

$$\mathsf{Prob}[|\overline{\mu} - \mu| \leq \epsilon\mu] \geq 1 - \delta$$

- δ: confidence parameter
- ϵ: error parameter

Intuitively: Use "Law of Large Numbers"

- law of large numbers (there are actually 2 versions) basically says that the sample mean tends to the true mean as the number of samples tends to infinity
- We take n samples $X_1, \dots, X_n$, and output

$$\bar{\mu} = \frac{1}{n}(X_1 + \cdots + X_n)$$

- But, how large must n be? ("Easy" if X is Bernoulli!)
- Markov is of some use, but only gives upper-tail bound
- Need a bound on the variance $\sigma^2 = \mathsf{Var}[X]$ too, to answer the question

Applying Chebyshev

- Let $Y = X_1 + \cdots + X_n$, then $\overline{\mu} = Y/n$ and $\mathsf{E}[Y] = n\mu$
- Since the X_i are independent, $\mathsf{Var}\left[Y\right] = \sum_i \mathsf{Var}\left[X_i\right] = n\sigma^2$
- Let $r = \sigma/\mu$, Chebyshev inequality gives

$$\begin{aligned}\mathsf{Prob}[|\overline{\mu} - \mu| > \epsilon\mu] &= \mathsf{Prob}\left[|Y - \mathsf{E}[Y]| > \epsilon\mathsf{E}[Y]\right] \\ &< \frac{\mathsf{Var}\left[Y\right]}{(\epsilon\mathsf{E}[Y])^2} = \frac{n\sigma^2}{\epsilon^2 n^2 \mu^2} = \frac{r^2}{n\epsilon^2}.\end{aligned}$$

- Consequently, $n = \frac{r^2}{\delta\epsilon^2}$ is sufficient!
- We can do better!

Finally, the Median Trick!

- If confident parameter is $1/4$, we only need $\Theta(r^2/\epsilon^2)$ samples; the estimate is a little "weak"
- Suppose we have w weak estimates $\mu_1, \dots, \mu_w$
- Output $\bar{\mu}$: the **median** of these weak estimates!
- Let I_j indicates the event $|\mu_j - \mu| \leq \epsilon\mu$, and $I = \sum_{j=1}^{w} I_j$
- By Chernoff's bound,

$$\begin{aligned}
\text{Prob}[|\overline{\mu} - \mu| > \epsilon\mu] &\leq \text{Prob}\left[Y \leq w/2\right] \\
&\leq \text{Prob}\left[Y \leq (2/3)\mathsf{E}[Y]\right] \\
&= \text{Prob}\left[Y \leq (1 - 1/3)\mathsf{E}[Y]\right] \\
&\leq \frac{1}{e^{\mathsf{E}[Y]/18}} \leq \frac{1}{e^{w/24}} \leq \delta
\end{aligned}$$

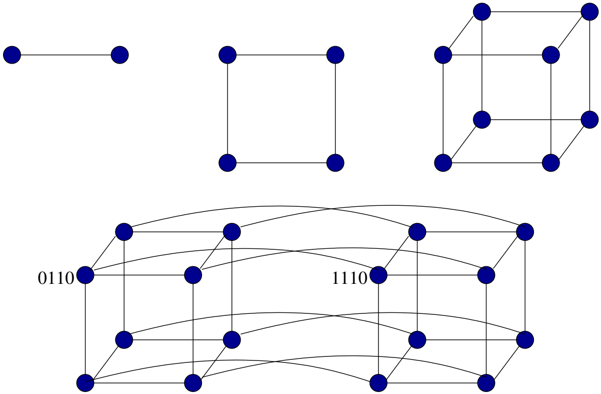
  whenever $w \geq 24\ln(1/\delta)$.
- Thus, the total number of samples needed is $n = O(r^2 \ln(1/\delta)/\epsilon^2)$.

- Directed graph $G = (V, E)$: network of parallel processors
- Permutation Routing Problem
  - Each node v contains one packet P_v, $1 \le v \le N = |V|$
  - Destination for packet from v is π_v, $\pi \in S_n$
  - Time is discretized into unit steps
  - Each packet can be sent on an edge in one step
  - Queueing discipline: FIFO
- Oblivious algorithm: route R_v for P_v depends on v and π_v only
- Question: in the worst-case (over π), how many steps must an oblivious algorithm take to route all packets?

Theorem (Kaklamanis et al, 1990)

Suppose G has N vertices and out-degree d. For any deterministic oblivious algorithm for the permutation routing problem, there is an instance π which requires $\Omega(\sqrt{N/d})$ steps.

The (Directed) Hypercube



- The n-cube: $|V| = N = 2^n$, vertices $\mathbf{v} \in \{0,1\}^n$, $\mathbf{v} = v_1 \cdots v_n$
- $(\mathbf{u}, \mathbf{v}) \in E$ iff their Hamming distance is 1

The Bit-Fixing Algorithm

- Source $\mathbf{u} = u_1 \cdots u_n$, target $\pi_u = v_1 \cdots v_n$
- Suppose the packet is currently at $\mathbf{w} = w_1 \cdots w_n$, scan $\mathbf{w}$ from left to right, find the first place where $w_i \neq v_i$
- Forward packet to $w_1 \cdots w_{i-1} v_i w_{i+1} \cdots w_n$

Source	010011
	110010
	100010
	100110
Destination	100111

- There is a π requiring $\Omega(\sqrt{N/n})$ steps

Valiant Load Balancing Idea

Les Valiant, *A scheme for fast parallel communication*, SIAM J. Computing, 11: 2 (1982), 350-361.

Two phase algorithm (input: π)

- **Phase 1:** choose $\sigma \in S_N$ uniformly at random, route P_v to σ_v with bit-fixing
- **Phase 2:** route P_v from σ_v to π_v with bit-fixing

This scheme is now used in designing Internet routers with high throughput!

Phase 1 Analysis

- P_u takes route $R_u = (e_1, \dots, e_k)$ to σ_u
- Time taken is k ($\leq n$) plus queueing delay

Lemma

If R_u and R_v share an edge, once R_v leaves R_u it will not come back to R_u

Theorem

Let S be the set of packets other than packet P_u whose routes share an edge with R_u, then the queueing delay incurred by packet P_u is at most $|S|$

Phase 1 Analysis

- Let H_{uv} indicate if R_u and R_v share an edge
- Queueing delay incurred by P_u is $\sum_{v \neq u} H_{uv}$.
- We want to bound

$$\text{Prob}\left[\sum_{v \neq u} H_{uv} > \alpha n\right] \geq \; ??$$

- Need an upper bound for $\mathsf{E}\left[\sum_{v \neq u} H_{uv}\right]$
- For each edge e, let T_e denote the number of routes containing e

$$\sum_{v \neq u} H_{uv} \leq \sum_{i=1}^{k} T_{e_i}$$

$$\mathsf{E}\left[\sum_{v \neq u} H_{uv}\right] \leq \sum_{i=1}^{k} \mathsf{E}[T_{e_i}] = k/2 \leq n/2$$

Conclusion

- By Chernoff bound,

$$\text{Prob}\left[\sum_{v \neq u} H_{uv} > 6n\right] \leq 2^{-6n}$$

- Hence,

Theorem

With probability at least $1 - 2^{-5n}$, every packet reaches its intermediate target (σ) in Phase 1 in $7n$ steps

Theorem (Conclusion)

With probability at least $1 - 1/N$, every packet reaches its target (π) in $14n$ steps