

Chapter 3: Order Notation

$\Theta, O, \Omega, o, \omega, \lfloor, \lceil, \log, \lg, \ln, !, \Sigma, \Pi$

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Asymptotic Behavior

- Let $f(n)$ and $g(n)$ be nonnegative functions when $n \geq 0$.
- Define $f(n) \prec g(n)$
 (read “ $f(n)$ is asymptotically less than $g(n)$ ”)
 if there exists a positive constant k such that
 $f(n) < g(n)$ when $n \geq k$.
 (This definition is not in the book.)
- Examples:
 $2n + 100 \prec 3n - 100$ (consider $n \geq 200$)
 $1000n^2 \prec 2^n/1000$ (consider $n \geq 30$)
- $\Theta(g(n))$ includes all functions bounded by constant factors of $g(n)$.

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Order Notation

- Theta-notation:
 $f(n) \in \Theta(g(n))$ if $c_1 g(n) \prec f(n) \prec c_2 g(n)$
 for some positive constants c_1 and c_2 .
- Big-Oh-notation:
 $f(n) \in O(g(n))$ if $f(n) \prec c g(n)$
 for some positive constant c .
- Big-Omega-notation:
 $f(n) \in \Omega(g(n))$ if $c g(n) \prec f(n)$
 for some positive constant c .
- little-oh-notation:
 $f(n) \in o(g(n))$ if $f(n) \prec c g(n)$
 for all positive constants c .

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Order Notation Continued

- little-omega-notation:
 $f(n) \in \omega(g(n))$ if $c g(n) \prec f(n)$
 for all positive constants c .

If $f(n)$ is $f(n)$ also ...?
 is ... $o(g(n))$ $O(g(n))$ $\Theta(g(n))$ $\Omega(g(n))$ $\omega(g(n))$

$o(g(n))$
 $O(g(n))$
 $\Theta(g(n))$
 $\Omega(g(n))$
 $\omega(g(n))$

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More Notation

$f(n)$ is ...	if ...
monotonically increasing	$m \leq n \rightarrow f(m) \leq f(n)$
monotonically decreasing	$m \leq n \rightarrow f(m) \geq f(n)$
strictly increasing	$m < n \rightarrow f(m) < f(n)$
strictly decreasing	$m < n \rightarrow f(m) > f(n)$

- $\lfloor x \rfloor$ is the floor of x . $\lceil x \rceil$ is the ceiling of x .
- $\sum_{i=0}^d a_i n^i$ is polynomial in n of degree d .
- a^n is exponential in n .
 $\log_a n$ is logarithmic in n .

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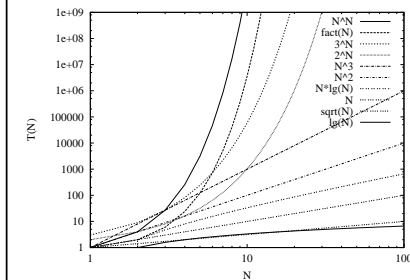
More Notation Continued

- Also, $\log_a n = \frac{\log_b n}{\log_b a}$
 $\lg n = \log_2 n$
 $\ln n = \log_e n$ where $e = 2.71828 \dots$
- $n!$ is n factorial.
- Let $0 < a < 1 < b < c$.
 $\Theta(1) \prec \Theta(\lg n) \prec \Theta(n^a) \prec \Theta(n) \prec \Theta(n \lg n) \prec \Theta(n^b) \prec \Theta(n^c) \prec$
 $\Theta(b^n) \prec \Theta(c^n) \prec \Theta(n!) \prec \Theta(n^n)$

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Log-Log Plot of Some Functions



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Summation Notation

- $\sum_{i=1}^n a_i = a_1 + a_2 + \dots + a_n$
- $\sum_{i=1}^n c a_i = c \sum_{i=1}^n a_i$
- $\sum_{i=1}^n (a_i + b_i) = \sum_{i=1}^n a_i + \sum_{i=1}^n b_i$
- If a_1, a_2, \dots, a_n is increasing, then:
 $\sum_{i=1}^n a_i$ is $O(n a_n)$
 $\sum_{i=1}^n a_i$ is $\Omega(n a_{n/2})$

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Some Miscellaneous Stuff

- $\prod_{i=1}^n a_i = a_1 a_2 \dots a_n$
- $\log_b \left(\prod_{i=1}^n a_i \right) = \sum_{i=1}^n (\log_b a_i)$
- Sets, Relations, Functions, Graphs, Trees
 See Appendix B

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