

Lecture 12

CSE 331

Sep 22, 2014

Today's agenda

Finish Proving n vertex tree has $n-1$ edges

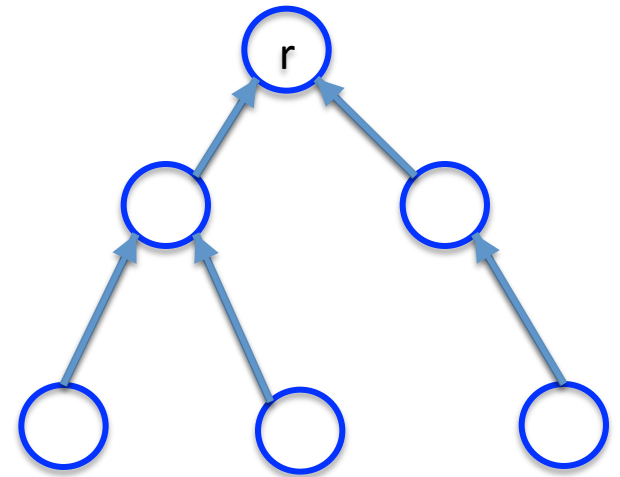
Algorithms for checking connectivity

n vertex tree has $n-1$ edges

Proof idea (recap):

Root tree T at vertex r

Direct edges from child to parent



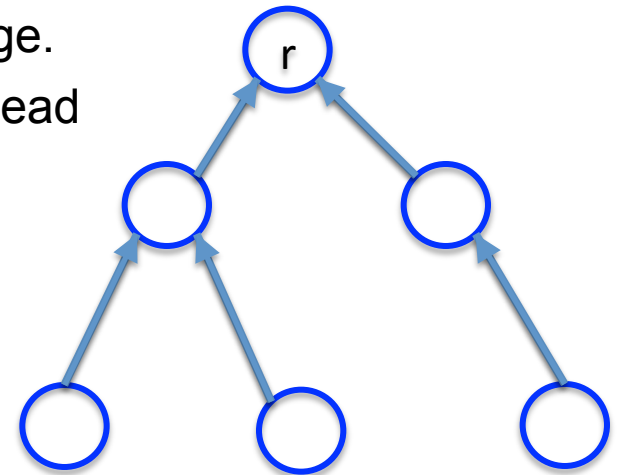
n vertex tree has $n-1$ edges

Proof idea (recap):

Claim 1: r is not the head of any directed edge.

Claim 2: Every directed edge has a unique head

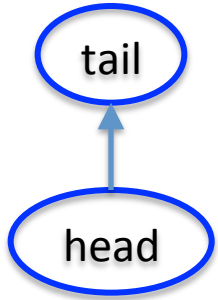
Claim 3: Every non-root $u \in V$ is the head of some directed edge.



Claim 4: Every $u \in V$ is the head of ≤ 1 edges.

Claim 1-4 \Rightarrow 1-to-1 correspondence between E and $V \setminus \{r\}$

$$\Rightarrow |E| = |V \setminus \{r\}| = n - 1$$



Questions?

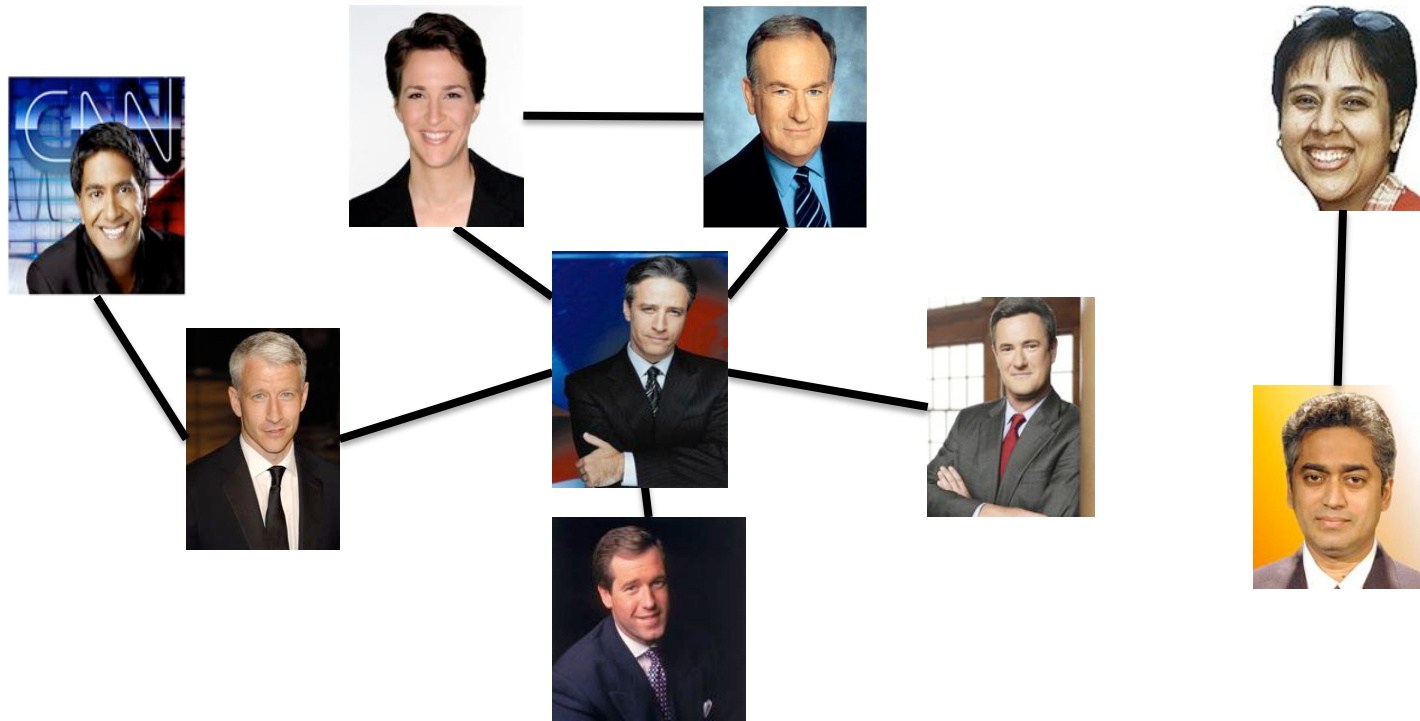


Rest of Today's agenda

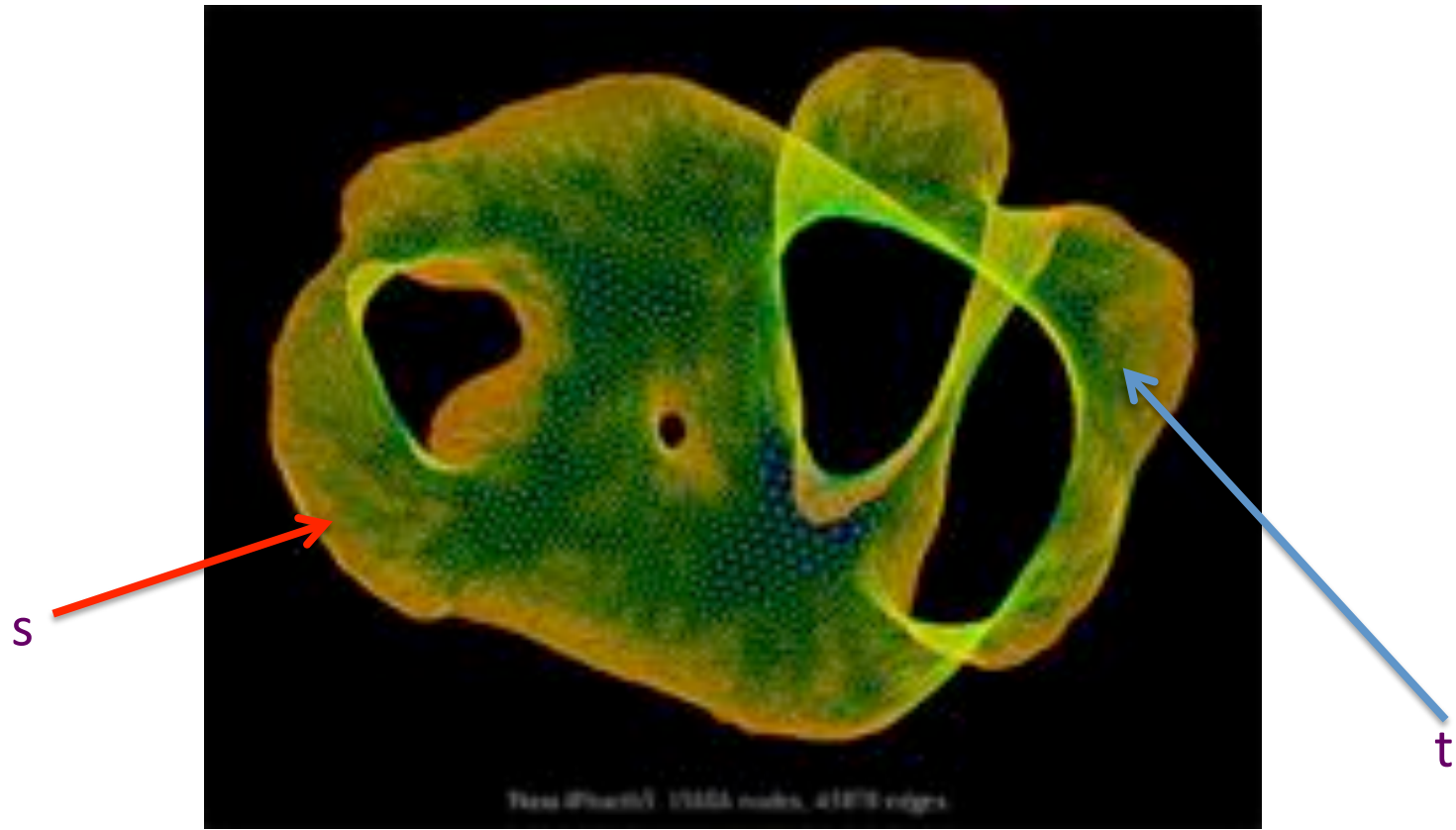
Finish Proving n vertex tree has $n-1$ edges

Algorithms for checking connectivity

Checking by inspection



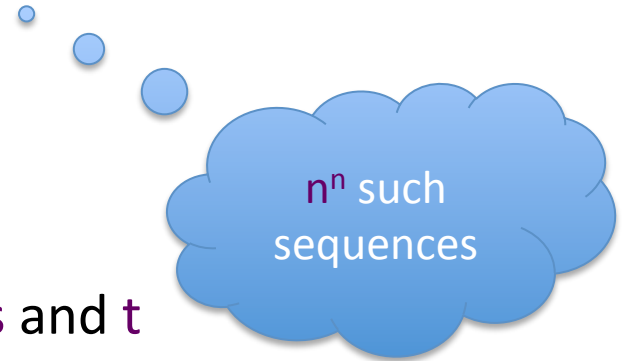
What about large graphs?



Are s and t connected?

Brute-force algorithm?

List all possible vertex sequences between s and t



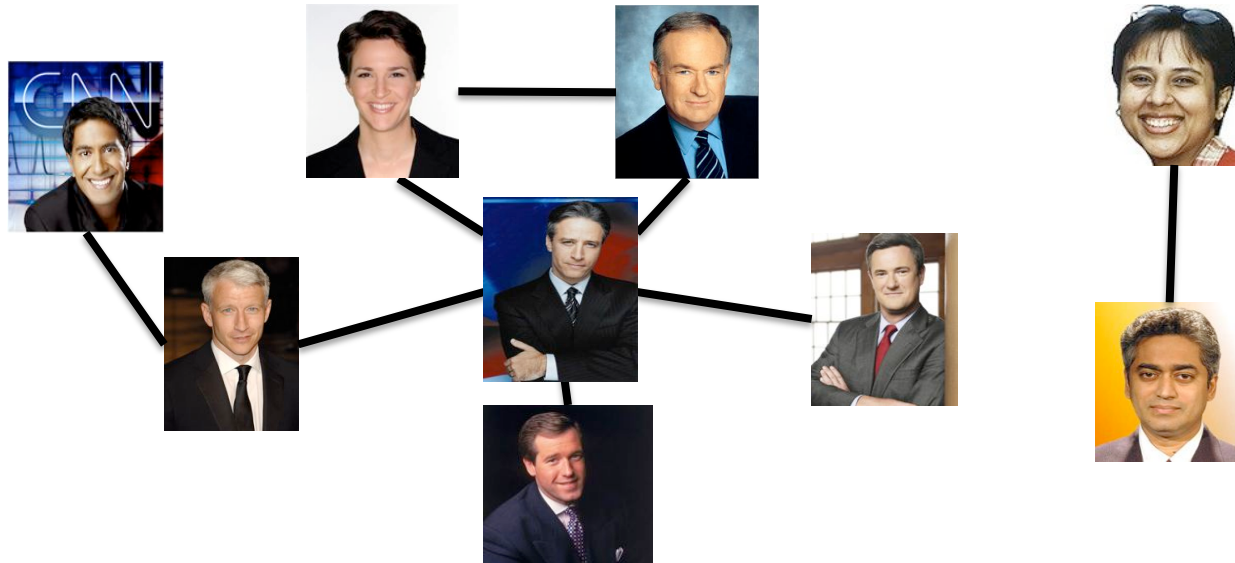
Check if any is a path between s and t

Algorithm motivation



Distance between **u** and **v**

Length of the shortest length path between **u** and **v**



Distance between RM and BO? 1

Questions?



Breadth First Search (BFS)

Is s connected to t ?

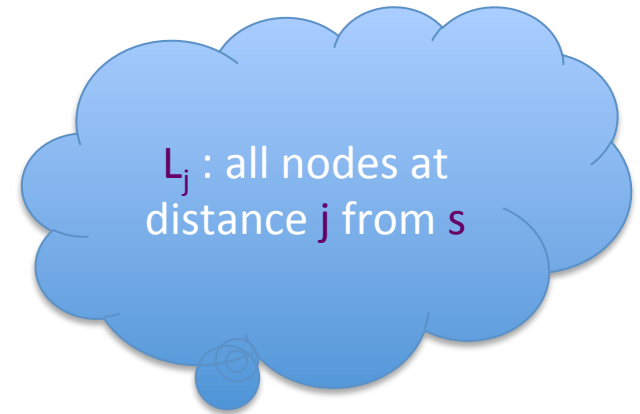
Build layers of vertices connected to s

$$L_0 = \{s\}$$

Assume L_0, \dots, L_j have been constructed

L_{j+1} set of vertices not chosen yet but are connected to L_j

Stop when new layer is empty



Exercise for you



Prove that L_j has all nodes at distance j from s

BFS Tree

BFS naturally defines a tree rooted at s

L_j forms the j th “level” in the tree

u in L_{j+1} is child of v in L_j from which it was “discovered”

Add non-tree edges

