

Lecture 7

CSE 331

Sep 10, 2014

Read the rubric for HW 1

note ☆ stop following 63 views

Homework 1 initial Grading Rubric

(Am pinning this post up till the deadline on Friday. --Atn)

The stuff below meant to illustrate how much points each part of the solution is worth
Follow the answer format provided here :

Problem 1 : (Exercise 1 & 2 in Chapter 1)

Statement : True (T) or False (F).
Choose True or False Correctly. (4 points).
If you choose True or False incorrectly, you will get zero in both **Proof Idea** and **Proof Detail** part.

Proof idea:
Here express clearly why you think it is True or False. Then explain how you are going to prove or disprove the statement. (6 points)

Proof details:
Here provide the proof if you choose True, otherwise provide (details of) a counterexample to disprove the statement. Note that a counterexample means a clearly specify the example but also argue why your example is actually a counterexample. (10 points)

Please review posted questions before posting yours

note stop following 1 views

Please read previous posts before posting your question

Hi guys,

I am really glad that you guys are asking questions: this is great!

However, there have been multiple questions now that have been already answered in a previous post. So please check existing posts **first** to make sure your question has not been answered already. If not, please go ahead and feel free to ask your question. If you see an answer that is related but does not quite answer your question, post a followup question on the earlier post.

Having too many posts with the same question just clutters the feed.

So please do keep your questions coming, but first verify if they have already been answered.

Thanks!
Atri

homework1

edit good note 0

23 seconds ago by Atri Rudra

Online Office Hours

Tonight: 9:00pm to 10:00pm

 note 

[stop following](#) **47** views
 [Actions](#) ▼

Online Office Hour #1 on Wednesday, Sep 10

I will hold the first online office hour this Wednesday, September 10 from 9:00pm to 10:00pm eastern. During that hour, I will be online and available to answer any question on piazza immediately.

Make sure to tag your question with the folder "onlineoh1" as I'll be only monitoring questions in that folder.

Any and all questions related to 331 are welcome!

(I have also added this office hours to the Google calendar for CSE 331 events.)

[onlineoh1](#)

[edit](#) - good note 

2 days ago by Abri Rudra

Gale-Shapley Algorithm

Initially all men and women are **free**

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to

w proposes to m

If m is free

(m,w) get **engaged**

Else (m,w') are engaged

If m prefers w' to w

w remains **free**

Else

(m,w) get **engaged** and w' is **free**

Output the engaged pairs S as the final output

Observation 1

Initially all men and women are **free**

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to

w proposes to m

If m is free

(m,w) get **engaged**

Else (m,w') are engaged

If m prefers w' to w

w remains **free**

Else

(m,w) get **engaged** and w' is **free**

Once a man gets engaged, he remains engaged (to “better” women)

Output the engaged pairs S as the final output

Observation 2

Initially all men and women are **free**

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to

w proposes to m

If m is free

(m, w) get **engaged**

Else (m, w') are engaged

If m prefers w' to w

w remains **free**

Else

(m, w) get **engaged** and w' is **free**

If w proposes to m after m' , then she prefers m' to m

Output the engaged pairs S as the final output

Proof via “progress”

Initially all men and women are **free**

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to

w proposes to m

If m is free

(m,w) get **engaged**

Else (m,w') are engaged

If m prefers w' to w

w remains **free**

Else

(m,w) get **engaged** and w' is **free**

End of iteration t ,
define progress

$P(t)$ s.t.:

1. $1 \leq P(t) \leq n^2$
2. $P(t+1) = P(t)+1$

Output the engaged pairs as the final output

Today's lecture

GS algorithms always outputs a stable marriage

The Lemmas

Lemma 1: The GS algorithm has at most n^2 iterations

Lemma 2: S is a perfect matching

Lemma 3: S has no instability

Questions/Comments?



Extensions

Fairness of the GS algorithm

Different executions of the GS algorithm

Main Steps in Algorithm Design

Problem Statement



Problem Definition



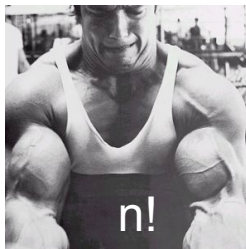
Algorithm



“Implementation”



Analysis



Correctness Analysis

Definition of Efficiency

An algorithm is efficient if, when implemented, it runs quickly on real instances

Implemented where?



What are real instances?

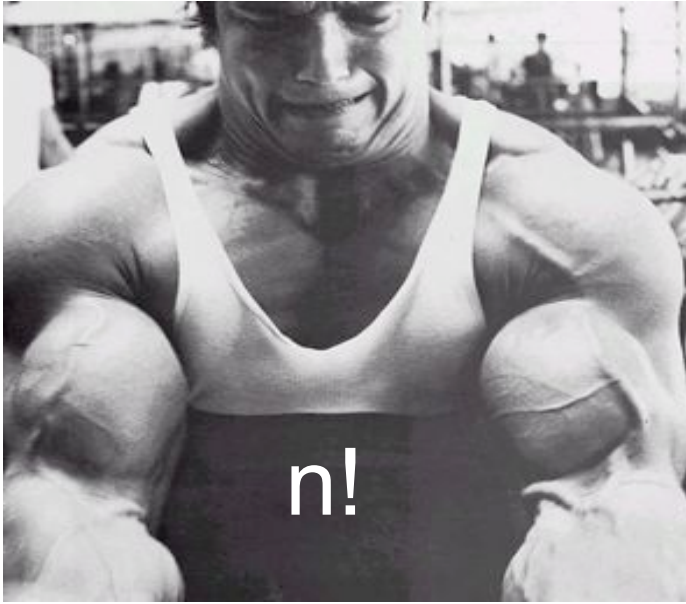
Worst-case Inputs

$$N = 2n^2 \text{ for SMP}$$

Efficient in terms of what?

Input size N

Definition-II



Analytically better than brute force

How much better? By a factor of 2?

Definition-III

Should scale with input size

If N increases by a constant factor,
so should the measure



Polynomial running time

At most $c \cdot N^d$ steps ($c > 0$, $d > 0$ absolute constants)

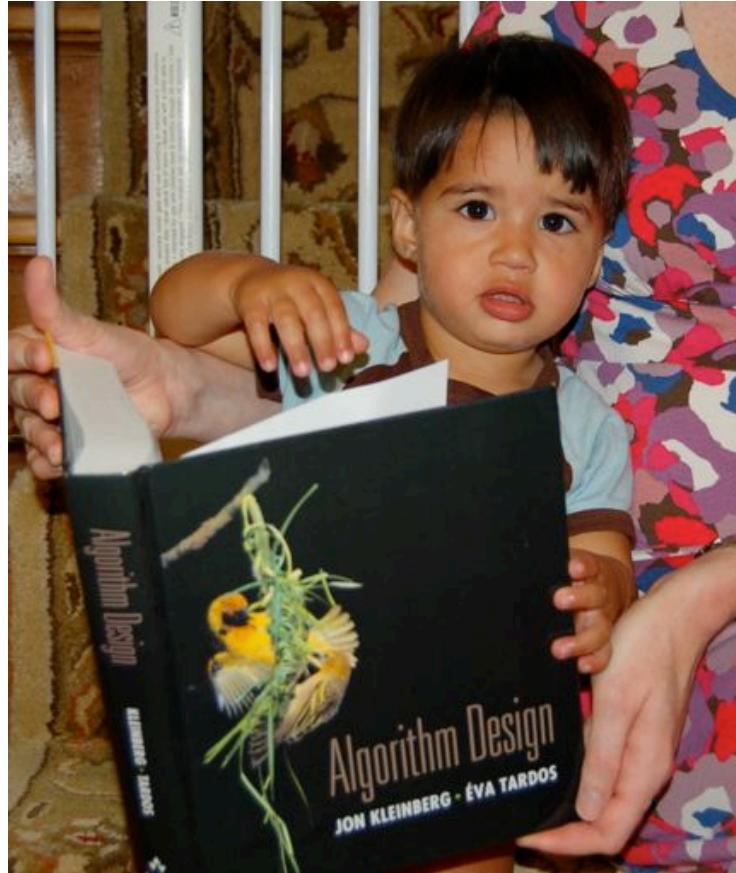
Step: “primitive computational step”

More on polynomial time

Problem centric tractability

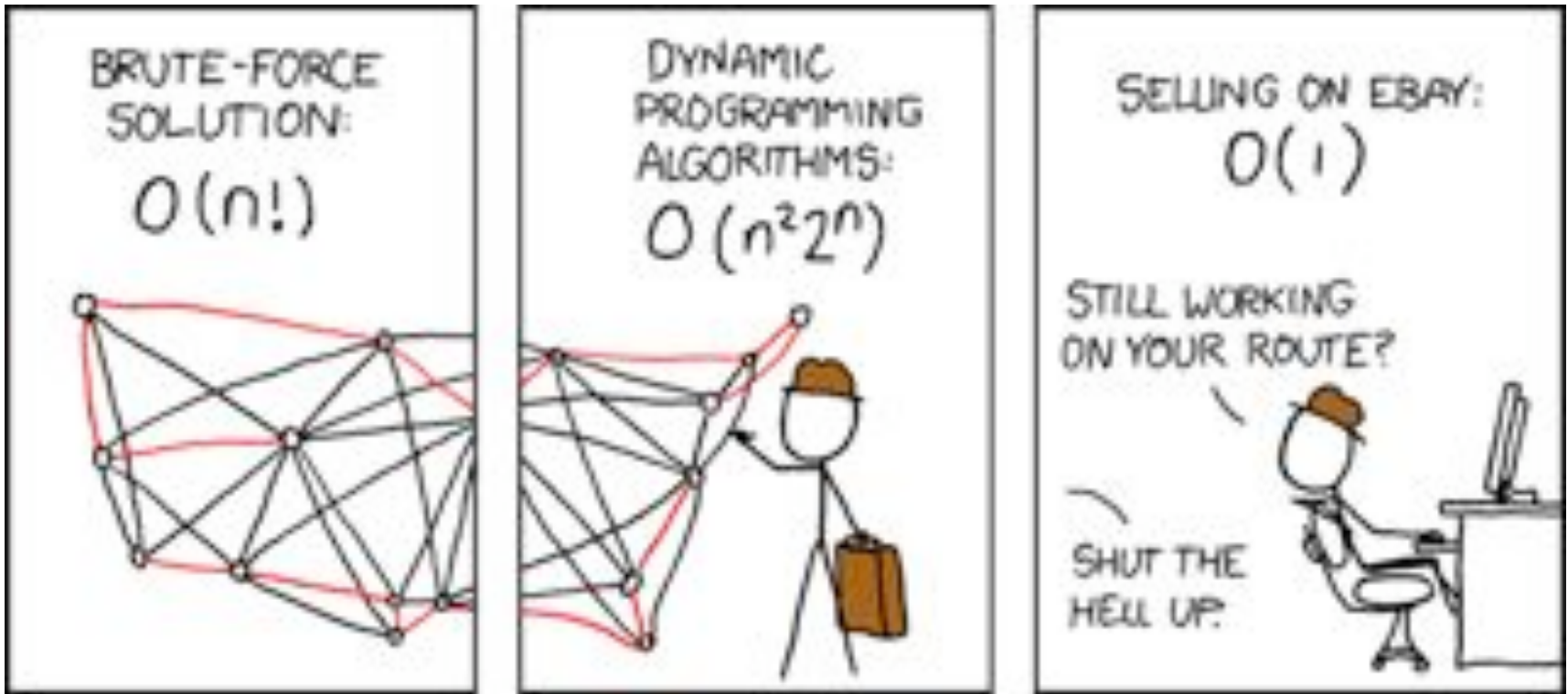
Can talk about problems that are not efficient!

Reading Assignments



Sections 1.2, 2.1, 2.2 and 2.4 in [KT]

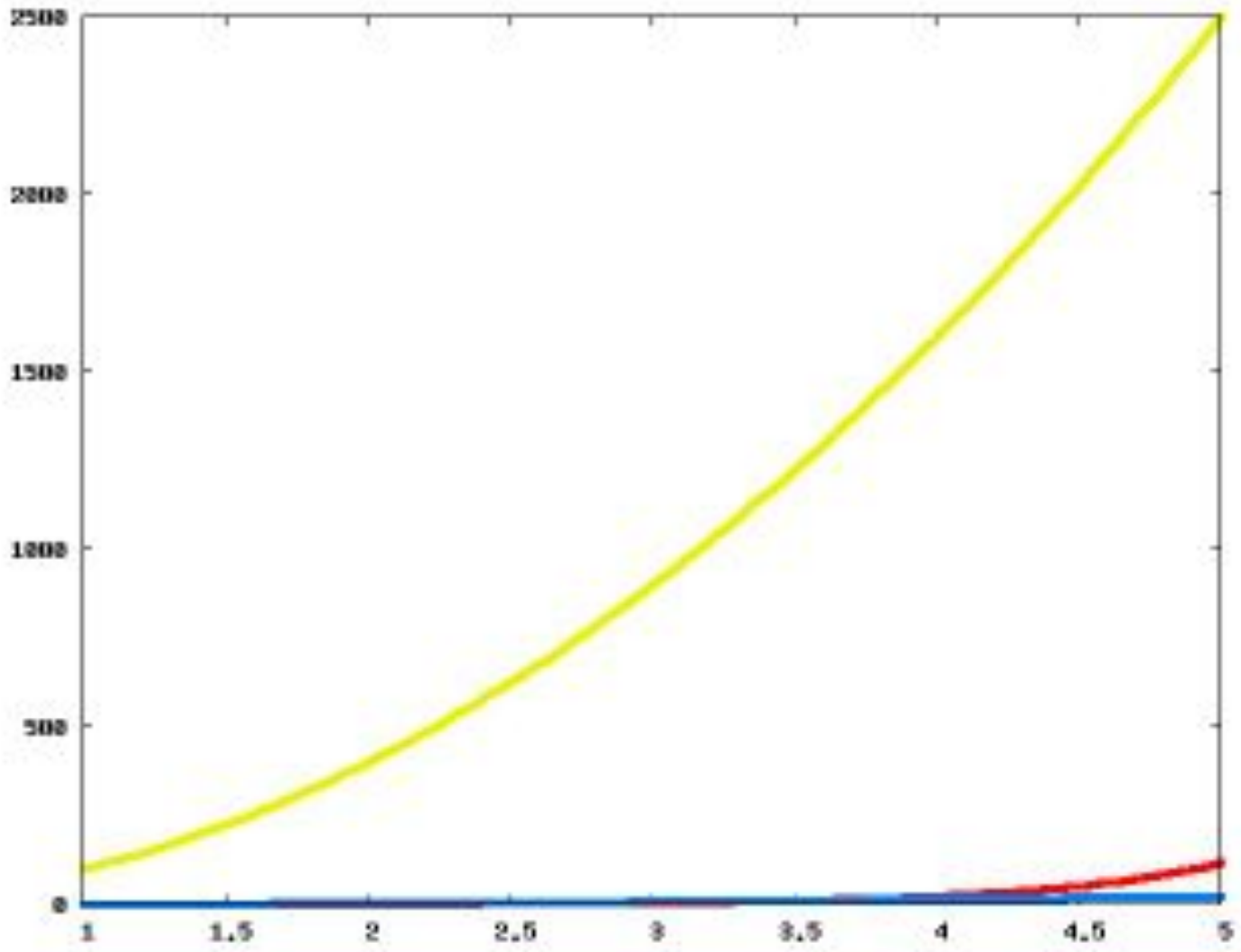
Asymptotic Analysis



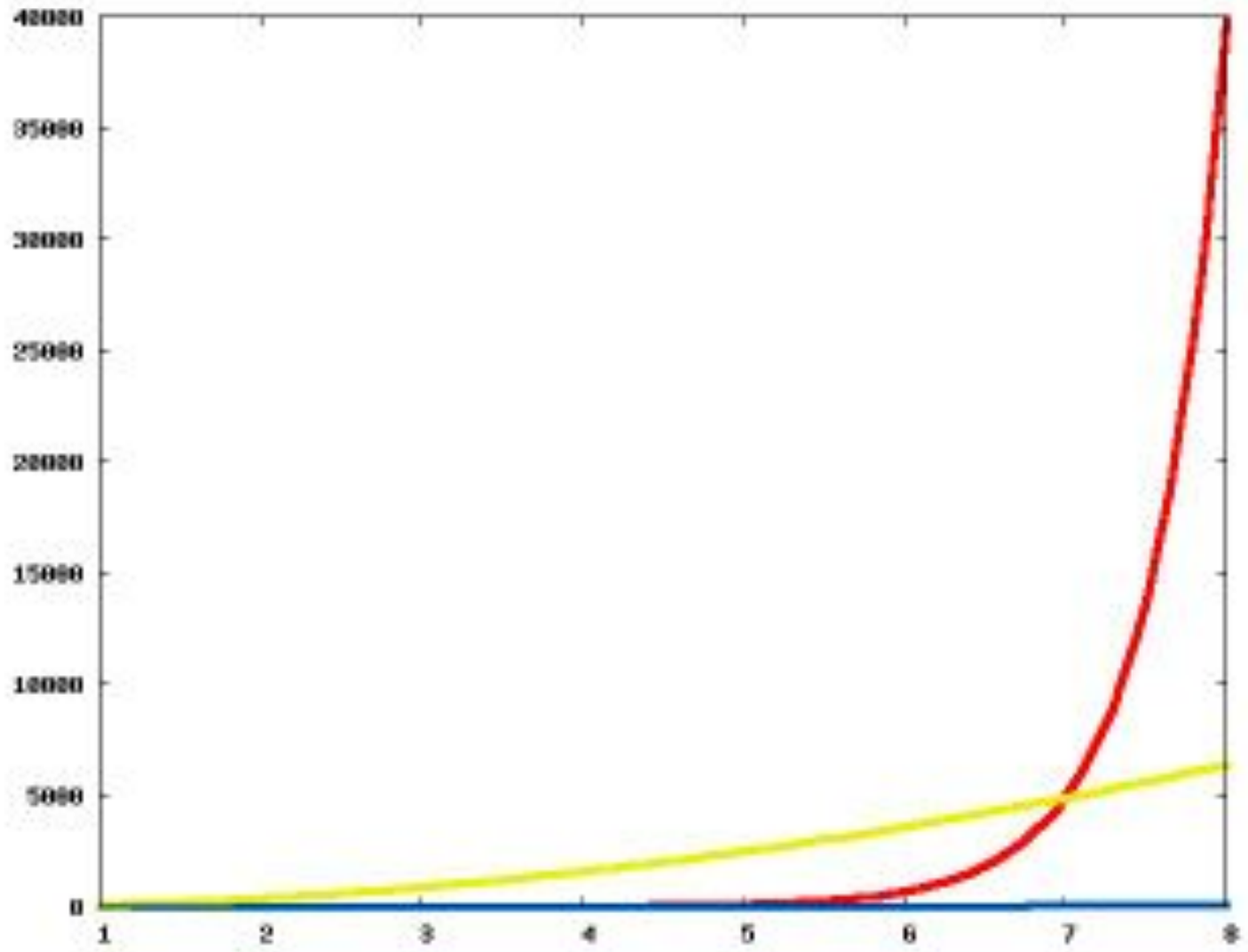
Travelling Salesman Problem

(<http://xkcd.com/399/>)

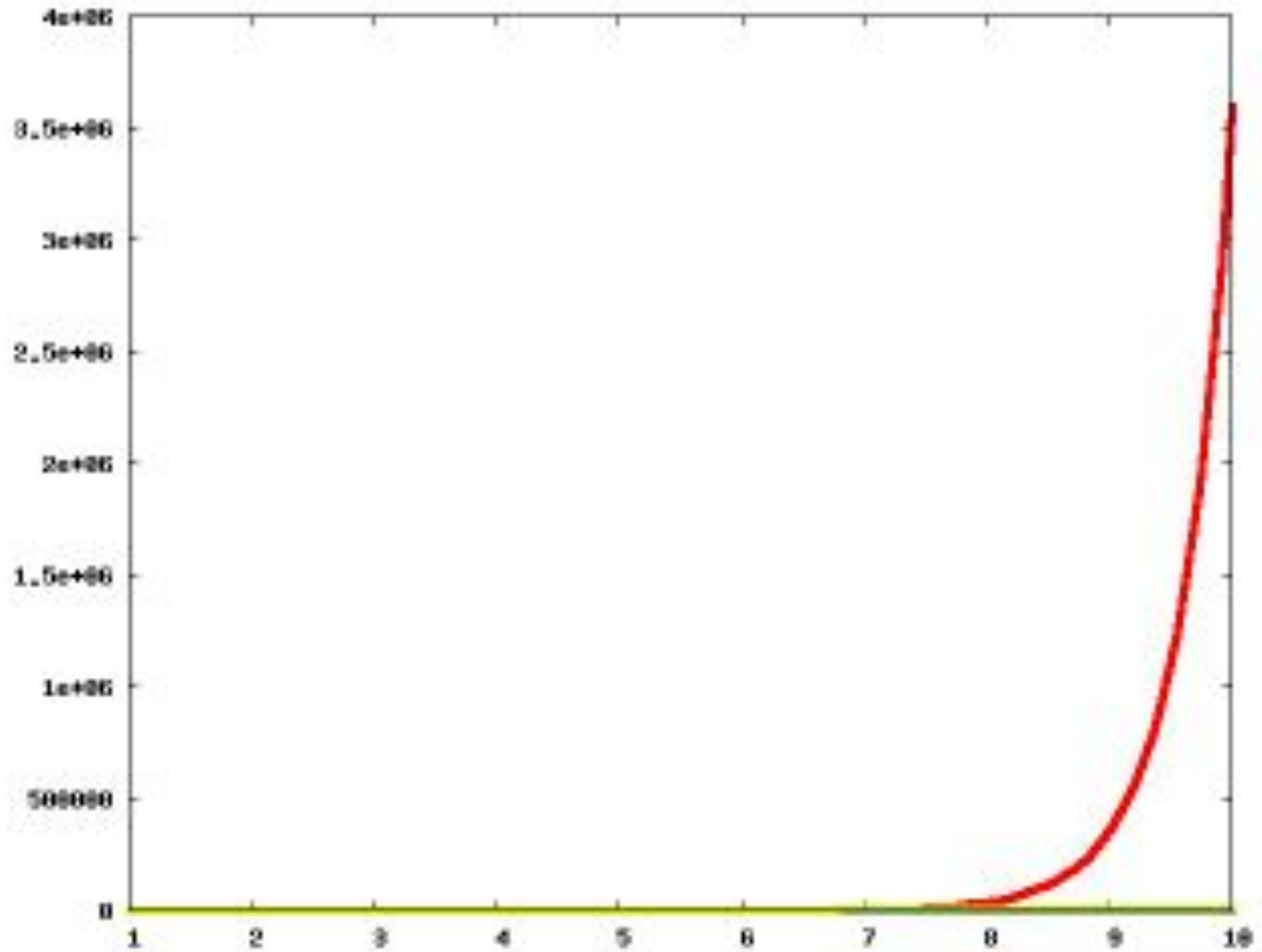
Which one is better?



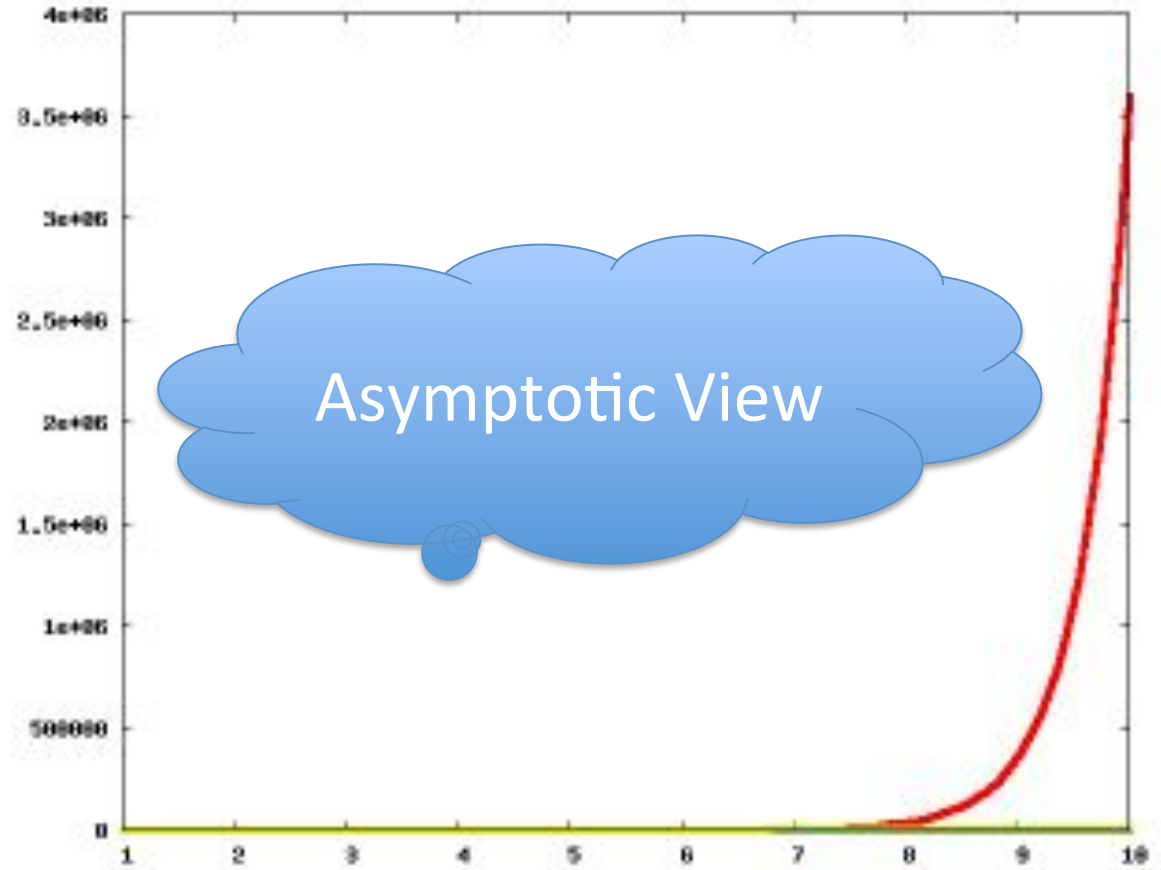
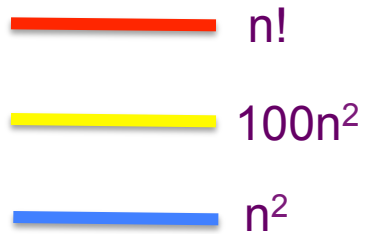
Now?



And now?



The actual run times



Asymptotic Notation



\leq is O with glasses

\geq is Ω with glasses

$=$ is Θ with glasses