

Oct 15, 2014

Lemma: We can convert an optional solution \mathcal{Q}' with 0 idle time to another optimal solution \mathcal{Q} with 0 idle time & 0 # inversions

If: $\text{idle}(\mathcal{Q}_i) = 0$

$\mathcal{Q}' = \mathcal{Q}_1 \rightarrow \mathcal{Q}_2 \rightarrow \dots \rightarrow \mathcal{Q}_i \rightarrow \mathcal{Q}_{i+1}$

Properties: $\forall i \rightarrow \mathcal{Q}_m = \mathcal{Q}$

(i) $\text{idle}(\mathcal{Q}_{i+1}) = 0$

(ii) If $\# \text{inv}(\mathcal{Q}_i) > 0 \Rightarrow \# \text{inv}(\mathcal{Q}_{i+1}) = \# \text{inv}(\mathcal{Q}_i) - 1$

(iii) $L(\mathcal{Q}_{i+1}) \leq L(\mathcal{Q}_i) \xrightarrow{i=m-1} L(\mathcal{Q}_m) \leq L(\mathcal{Q}_{m-1}) \leq L(\mathcal{Q}_{m-2}) \dots \leq L(\mathcal{Q}_1) \leq L(\mathcal{Q}')$

$\forall i \Rightarrow L(\mathcal{Q}_m) \leq L(\mathcal{Q}_1) = L(\mathcal{Q}')$
 $\leq L(\mathcal{Q})$
 $\Rightarrow L(\mathcal{Q}) = L(\mathcal{Q}')$ (hence, opt.)
as \mathcal{Q}' is optimal

Conversion: $\mathcal{Q}_i \rightarrow \mathcal{Q}_{i+1}$ $\# \text{inv}(\mathcal{Q}_i) > 0$

~~idle~~ $\text{idle}(\mathcal{Q}_i) = 0$

Property (a) (Since $\# \text{inv}(\mathcal{Q}_i) > 0$) $\exists (j, k) s.t.$

i) (j, k) is an inversion ii)



Property (b) Swap j & k to get \mathcal{Q}_{i+1}

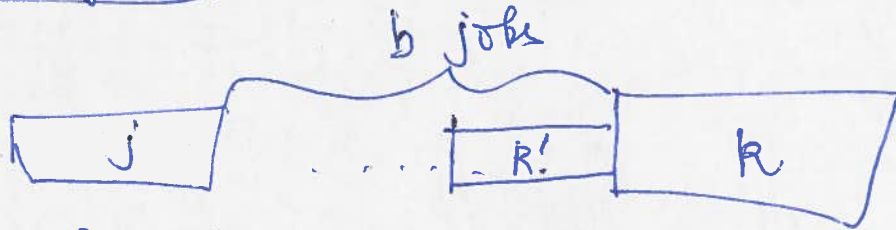
$\# \text{inv}(\mathcal{Q}_{i+1}) = \# \text{inv}(\mathcal{Q}_i) - 1$
 $\text{idle}(\mathcal{Q}_{i+1}) = 0$

$\left. \begin{array}{l} \text{swap to get } \mathcal{Q}_{i+1} \\ \text{follows by construction of } \mathcal{Q}_{i+1} \end{array} \right\}$

Property (c) $L(\mathcal{Q}_{i+1}) \leq L(\mathcal{Q}_i)$

If of property (a) know.

d_k d_j

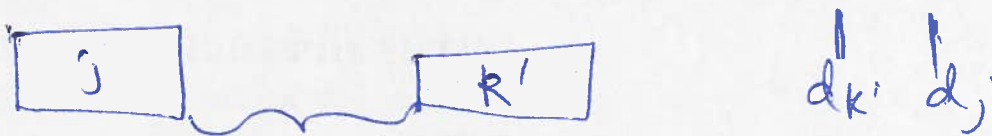


Case 1: If j & k are right next to each other: ✓

Case 2: k' next to k closer to j

Case 2.1: $d_{k'} > d_k \Rightarrow (k', k)$ is our pair.

Case 2.2: $d_{k'} \leq d_k < d_j \Rightarrow d_{k'} < d_j$



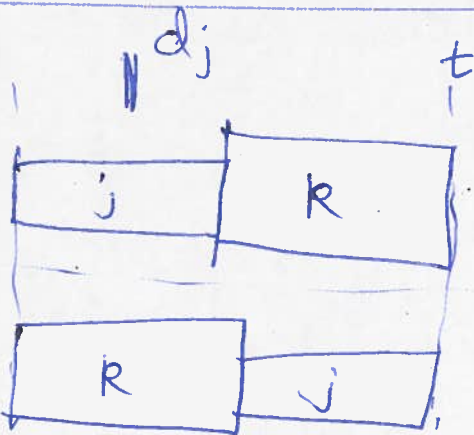
Continue till $b=0 \Rightarrow$ Case 1 ✓

Property (c)

Q_i

Q_{i+1}

$\forall m \notin \{j, k\}$



same

same

$$l(m) = l'(m)$$

lateness in Q_i lateness in Q_{i+1}

$$L(Q_{i+1}) \leq L(Q_i)$$

\Rightarrow Enough to argue

$$\max\{l'(j), l'(k)\} \leq \max\{l(j), l(k)\}$$

$$\max\{l'(j), l'(k)\} \leq l(k)$$

(i) $l'(k) \leq l(k)$ (as in Q_{i+1} k finishes $< t$)

~~Now if we could pr~~

→ Possible that $l'(j) > l(j)$

$$(ii) \quad l'(j) = t - d_j$$

$$\leq t - d_k$$

$$= l(k)$$

(assume $d_j < t$)

$$d_j > d_k$$

$$\Rightarrow -d_j < -d_k$$

as (j, k) an
inv

□