Using MPI to Break the Data Encryption Standard

CSE 633: Parallel Algorithms
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What is DES?

- DES (or Data Encryption Standard) is an algorithm used to encrypt electronic data.
- Encryption takes place as bitwise data transformation over several stages.
- Developed in the late 1970s and implemented by the early 1980s.
- Still used today, many variants are out there such as AES (Advanced Encryption Standard).
DES Overview

```
Encryption

“Hello, how are you?”
(Cleartext)

Encryption key
00111010001001
1110100001010

Decryption key

Decryption

(Ciphertext)
```
Structure of DES: Single iteration
Short-Key Weakness

- Cryptographers built a machine to break DES in 1998.
  - Cost: About $250,000
  - Took about 56 hours to find a key

- DES uses a 64 bit key but uses only 56 of those bits for encryption.
  - When it was first implemented, this was good enough
  - Nowadays, this is too short. More computing power means keys can be obtained by an exhaustive search, “brute force”
Known-Plaintext Attack

- We know a corresponding plaintext with ciphertext
- We run DES encryption on each 56-bit key in the key space and check if it matches the ciphertext.
- You can perform both ways because it is symmetric.
- Our goal is to find the key.
- Each processor will receive a subset of the key space to search.
- This is different from chosen-plaintext attack.
Total KeySpace to Search

- The total keyspace is $2^{56}$
- $2^{56} = 7.2057594037927936 \times 10^{16}$
- 72 quadrillion... = 72 million billion
- This is reduced by one half since taking the complement of the plaintext and complement of key will result in the same ciphertext. Therefore we can just disregard the complement of a key we already attempted.
- $(1/2) \times 2^{55} = 3.6028797018963968 \times 10^{16}$
- 36 quadrillion... = 36 million billion
Breaking DES (Logic)

- Distribute the key space among workers.
- Master is responsible for synchronization of the program.
- Each worker is given the cipher text and corresponding plaintext.
- The worker tries to find the key by encrypting the plain text and comparing it with cipher text using its own key space.
- As soon as it breaks the key it notifies the master and the master notifies all other workers to stop processing.
MPI Implementation Results
<table>
<thead>
<tr>
<th>NUMBER OF WORKERS</th>
<th>Keys Checked per sec</th>
</tr>
</thead>
<tbody>
<tr>
<td>36 workers</td>
<td>19097562</td>
</tr>
<tr>
<td>48 workers</td>
<td>25461504</td>
</tr>
<tr>
<td>72 workers</td>
<td>38192256</td>
</tr>
<tr>
<td>98 workers</td>
<td>51983904</td>
</tr>
<tr>
<td>120 workers</td>
<td>63653760</td>
</tr>
<tr>
<td>168 workers</td>
<td>89115264</td>
</tr>
</tbody>
</table>
Keys Searched Per Sec (approximately)

Keys per sec

0 20000000 40000000 60000000 80000000 100000000 120000000 140000000 160000000 180000000

36 48 72 98 120 168 200 300

workers workers workers workers workers workers workers workers workers
Importance of Key position

- In this program the running time is extremely affected by the position of the key.
- A particular key will be located at the different positions according to the total number of workers chosen.
- Due to this it is very important to calculate the results based on a specific position of the key for a particular number of workers.
Graphs without taking KEY POSITION into consideration
Implementation

- These graphs are calculated using random keys.
- The running time is in SECONDS.
- As we will see further that we do not get any kind of consistent results due to it.
- These graphs are included just to highlight the importance of key position.
We observed the number of nodes vs running time to be as follows: (1
N Cores/Node

![Graph showing the relationship between number of nodes and running time. The graph has a y-axis labeled 'Running Time (s)' ranging from 0 to 700 and an x-axis labeled 'Number of Nodes' ranging from 0 to 35. The graph shows a peak around 10 nodes with running time increasing before decreasing.]
Graphs obtained taking KEY POSITION into consideration
Implementation

- We used 0.001 percent location of key consistently for variable number of Workers.
- This is done due to the long running time constraint.
- To select the key we wrote a small function.
- The function takes input as total key space and total number of workers.
- It also takes a particular location of key as input in percentage form. E.g. Find the key located at 10 percent of each worker when key space is $2^{56}$ and total number of workers are 32.
Implementation

- The function provides output as the keys in each worker at that particular user defined location as well as approximate time required to break the key.
- The graphs are calculated using one of the random keys provided by the function.
# Observation Table for 1 Core Per Node

<table>
<thead>
<tr>
<th>Number of Workers</th>
<th>Time Required in Hours to Break the Cipher Text</th>
</tr>
</thead>
<tbody>
<tr>
<td>36</td>
<td>8.23</td>
</tr>
<tr>
<td>72</td>
<td>5.57</td>
</tr>
<tr>
<td>98</td>
<td>3.56</td>
</tr>
<tr>
<td>168</td>
<td>2.24</td>
</tr>
<tr>
<td>200</td>
<td>2.38</td>
</tr>
<tr>
<td>300</td>
<td>1.03</td>
</tr>
</tbody>
</table>
1 Core Per Node

Time Required in Hours

- 35 workers
- 71 workers
- 97 workers
- 168 workers
- 200 workers
- 300 workers

Time Required in Hours
# Observation Table for N Tasks Per Node

<table>
<thead>
<tr>
<th>Number of Workers and Their Configuration</th>
<th>Time Required in Hours to Break the Cipher Text</th>
</tr>
</thead>
<tbody>
<tr>
<td>12nodes 3 tasks per node</td>
<td>8.32</td>
</tr>
<tr>
<td>12nodes 6 tasks per node</td>
<td>5.69</td>
</tr>
<tr>
<td>14nodes 7 tasks per node</td>
<td>3.35</td>
</tr>
<tr>
<td>14nodes 12 tasks per node</td>
<td>2.27</td>
</tr>
<tr>
<td>20nodes 10 tasks per node</td>
<td>2.38</td>
</tr>
<tr>
<td>60nodes 5 tasks per node</td>
<td>1.18</td>
</tr>
<tr>
<td>N</td>
<td>Tasks Per Node</td>
</tr>
<tr>
<td>--------</td>
<td>----------------</td>
</tr>
<tr>
<td>12 nodes</td>
<td>3 tasks per node</td>
</tr>
<tr>
<td>12 nodes</td>
<td>6 tasks per node</td>
</tr>
<tr>
<td>14 nodes</td>
<td>7 tasks per node</td>
</tr>
<tr>
<td>14 nodes</td>
<td>12 tasks per node</td>
</tr>
<tr>
<td>20 nodes</td>
<td>10 tasks per node</td>
</tr>
<tr>
<td>60 nodes</td>
<td>5 tasks per node</td>
</tr>
</tbody>
</table>
Things Learned.....

- The most important part of this project was the enormous data size.
- So we learned how to handle as well as compute such a huge data size and make efficient use of computing power provided to us.
- It also painfully taught us to debug long running parallel programs eg. Getting errors after the program runs for a long time.
- One of the important features particular to this project was the impact of key position and the calculation required to find ideal keys prior to computation.
- So in short it did teach some basic techniques of Parallel Programming even though sometimes the lessons were hard.
References & Resources


- Center for Concurrent Research (CCR) Systems

  - [http://w2.eff.org/Privacy/Crypto/Crypto_misc/DESCracker/HTML/19980716_eff_des_faq.html](http://w2.eff.org/Privacy/Crypto/Crypto_misc/DESCracker/HTML/19980716_eff_des_faq.html)

- *The DES Algorithm Illustrated* by J. Orlin Grabbe
Thank you for your attention! 😊