## PRIME FACTORIZATION

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## Background

- Prime factorization is the process of breaking down a composite number into its prime factors, which are the prime numbers that multiply together to equal the original number.
- Some of the applications are :
- Cryptography(RSA, Computer Security)
- Physics(Study properties of materials \& their electronic structures)
- Database Design(Unique Identifiers)
- Optimization Problems(Computationally Intensive)
- Cryptocurrency(Proof-Of-Work System)


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## Naive Sequential Algorithm

- Iterate through 2 to n
- Divide the number $n$ until its evenly divisible.
- Evenly divisible number is one of the factor of the number $n$.



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## WORKING -

Sieve(n):

```
Input: an integer n}>1
```

Let $A$ be an array of Boolean values, indexed by integers 2 to $n$,
initially all set to true.
for $i=2,3,4, \ldots$, not exceeding $\sqrt{n}$ :
if $A[i]$ is true:
for $j=i^{2}, i^{2}+i, i^{2}+2 i, i^{2}+3 i, \ldots$ not exceeding $n$ :
$A[j]$ : $=$ false.
Output: all $i$ such that $A[i]$ is true.
primes = Sieve(n)
for primeNumber in primeNumbers:
while n is divisible by primeNumber:
add primeNumber to the factor list divide n by primeNumber and update

## NEED FOR PARALLELIZATION

- Improve Factorization Performance
- Faster Execution time
- Scaling for larger inputs


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## PARALLELIZATION STRATEGY



## PARALLELIZATION STRATEGY

- Parallelize sieve to find prime factors
- Then, distribute the primes equally into number of processors except the last processor
- Machine processes only the subset of the primes and find factors from them


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## PARALLEL ALGORITHM

- Divide input range across all the processors
- Apply sieve by marking multiples of prime in the range
- On completion, broadcast prime number to other processors to eliminate non-primes
- Consolidate(MPI_Reduce)
- Distribute the primes equally across the processors
- $\mathrm{n}^{\text {th }}$ rank * chunksize $<=$ nth processor $<\mathrm{n}^{\text {th }}$ rank * chunksize + chunksize
- Find prime factors which evenly divide $\mathbf{n}$ in the range
- Terminate


## Improvements/Observations

- Parallelized Sieve of Eratosthenes(Previous Bottleneck) $V$
- Scaled up the number of nodes $\sqrt{ }$
- Increase in input size and increase in number of nodes/processors is directly proportional w.r.t performance
- Decrease in input size and increase in number of nodes/processors is inversely proportional w.r.t performance
- Implement Lowest Prime Factor for per processor optimization (LPF) X
- LPF[18] $=2$ => 18/2 = 9;
- LPF[9] $=3$ => 9/3 = 3;
$-\operatorname{LPF}[3]=3=>3 / 3=1$;



## Results with equal distribution of prime factors

| Processors | Time(in seconds) |  |
| ---: | ---: | ---: |
|  | 1 |  |
|  | 2 |  |
|  | 4 |  |
|  | 8 | 779.814 |
|  | 16 | 455.916 |
| 32 | 234.902 |  |
| 64 | 91.448 |  |


| Processors |  | Time(in seconds) |
| ---: | ---: | ---: |
|  | 16 |  |
|  | 32 |  |
|  | 64 | 52.600724 |
| 128 | 26.33962 |  |
| 256 |  | 13.821954 |
| 512 | 6.867855 |  |
| 1024 | 3.586585 |  |
| 2048 |  | 1.898545 |

## Results with equal distribution of prime factors



Time vs Processors ( $\mathrm{n}=12345678945$ )


## Results with exponential distribution of prime factors

| Processors |  | Time(in seconds) |
| ---: | ---: | ---: |
|  | 1 |  |
|  | 2 | 1023.654 |
|  | 4 | 800.814 |
| 8 | 480.655 |  |
| 16 | 290.659 |  |
| 32 | 150.598 |  |
| 64 | 100.665 |  |


| Processors | Time(in seconds) |
| :---: | :---: |
| 16 | 114.598 |
| 32 | 60.598 |
| 64 | 40.566 |
| 128 | 25.526 |
| 256 | 15.565 |
| 512 | 10.889 |
| 1024 | 5.233 |
| 2048 | 2.265 |

## Results with exponential distribution of prime factors



Time vs Processors ( $\mathrm{n}=12345678945$ )


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## Variation for $n$



Time vs Processors ( $\mathrm{n}=5689721545$ )


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## Variation for $n$

Time vs Processors ( $\mathrm{n}=562116359841$ )


Time vs Processors ( $\mathrm{n}=562116359841$ )


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## Variation for n



Time vs Processors ( $\mathrm{n}=8989454632$ )


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## Variation for n



Time vs Processors ( $\mathrm{n}=965612456285$ )


## Results for small n

| Processors |  | Time(in seconds) |
| ---: | ---: | ---: |
|  | 1 |  |
|  | 2 |  |
|  | 4 | 913.254 |
| 8 | 739.414 |  |
| 16 | 415.516 |  |
| 32 | 194.501 |  |
| 64 | 51.047 |  |


| Processors |  | Time(in seconds) |
| ---: | ---: | ---: |
|  | 16 |  |
|  | 32 |  |
|  | 48 |  |
|  | 64 | 29.565 |
| 128 |  | 25.111 |
| 256 |  | 20.569 |
|  | 512 |  |
| 1024 |  | 32.789 |
|  |  | 45.598 |

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## Small n

Time vs Processors ( $\mathrm{n}=800$ )


Time vs Processors ( $n=800$ )


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## Small n

Time vs Processors ( $\mathrm{n}=1500$ )


Time vs Processors ( $n=1500$ )


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## REFERENCES

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## Thank You! Questions?

