

PARALLELIZING PAGE RANK

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History

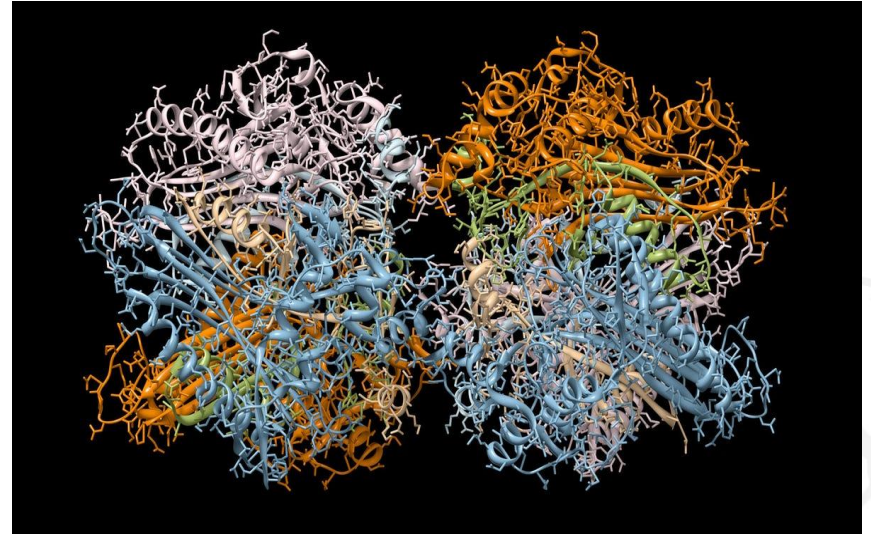
The heart of the Google search engine is the PageRank algorithm, which was described in the paper titled *The PageRank Citation Ranking: Bringing Order to the Web*.



Sergie Brin & Larry Page

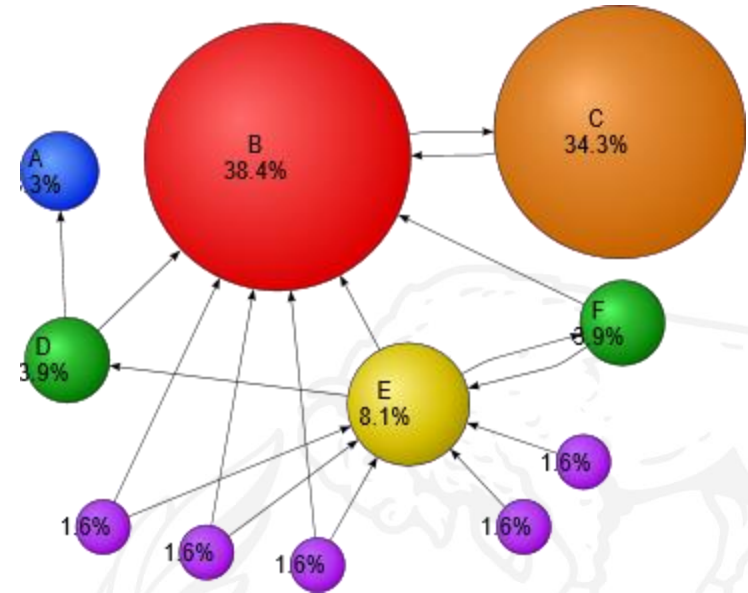
Web Graph: The World Wide Web is a graph where web pages are nodes, and hyperlinks between pages are edges. This graph considers a web page with many incoming links more important or authoritative.

Applications of Page Rank



Algorithm

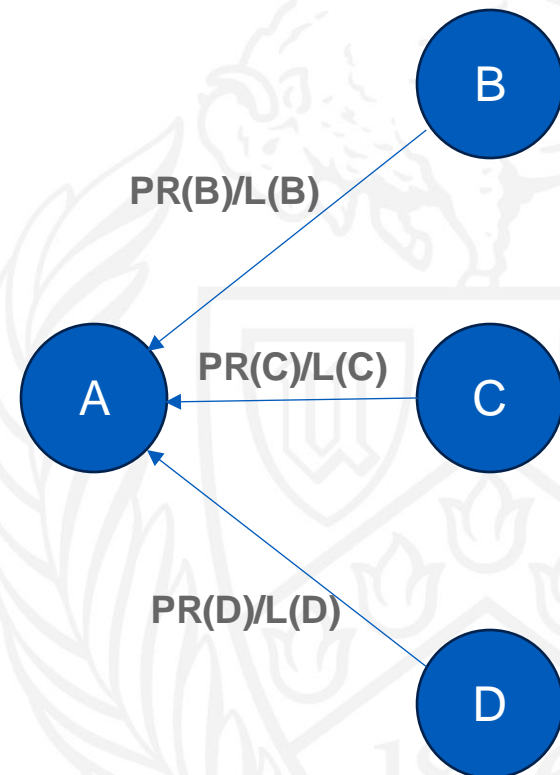
- The PageRank algorithm gives each page a rating of its importance, which is a recursively defined measure whereby a page becomes important if important pages link to it. This definition is recursive because the importance of a page refers back to the importance of other pages that link to it.
- One way to think about PageRank is to imagine a random surfer on the web, following links from page to page. The page rank of any page is roughly the probability that the random surfer will land on a particular page. Since more links go to the important pages, the surfer is more likely to end up there.



Simplified Sequential implementation

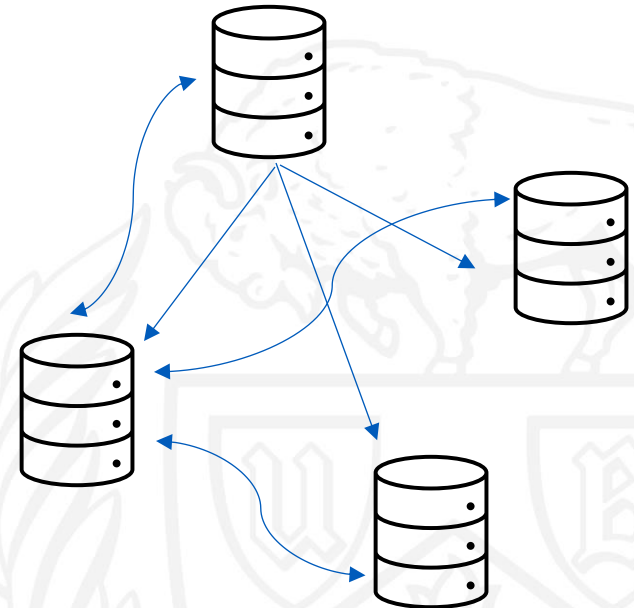
- Iterate over each of the vertices and update individual ranks based on incoming edges.
- All these fractions of votes are added together but, to stop the other pages having too much influence, this total vote is damped down by multiplying it with d .

$$PR(A) = \frac{1-d}{N} + d \left(\frac{PR(B)}{L(B)} + \frac{PR(C)}{L(C)} + \frac{PR(D)}{L(D)} + \dots \right).$$



Parallelized implementation

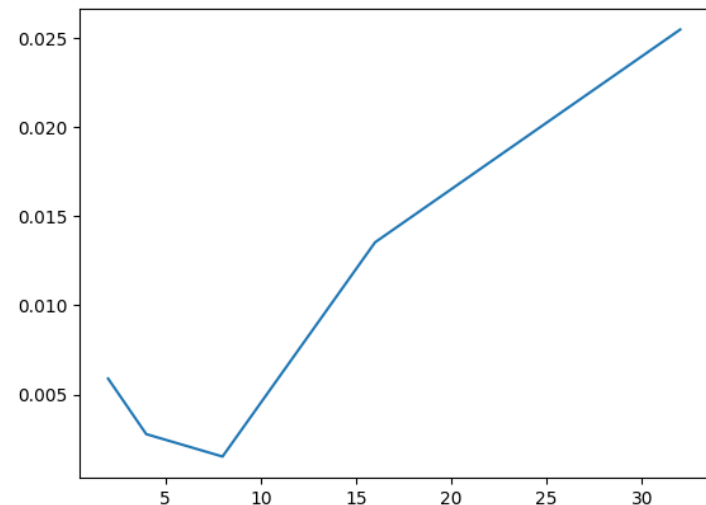
- Root processor reads from a file.
- Divide the $N \times N$ graph into (p) partitions and broadcast it to each of the nodes. So $[N/p]$ rows for each.
- Iteratively (*timer starts here*)
 - Calculate the latest PR value
 - Share (the local start to end row in page rank array) with rest of the nodes and update each of its values.
- Perform above 2 steps for k times.
- End the timer & save the rank values



Results

Time taken for different number of processors for a graph with 100 vertices and 2000 edges.
Lowest: 8 nodes.

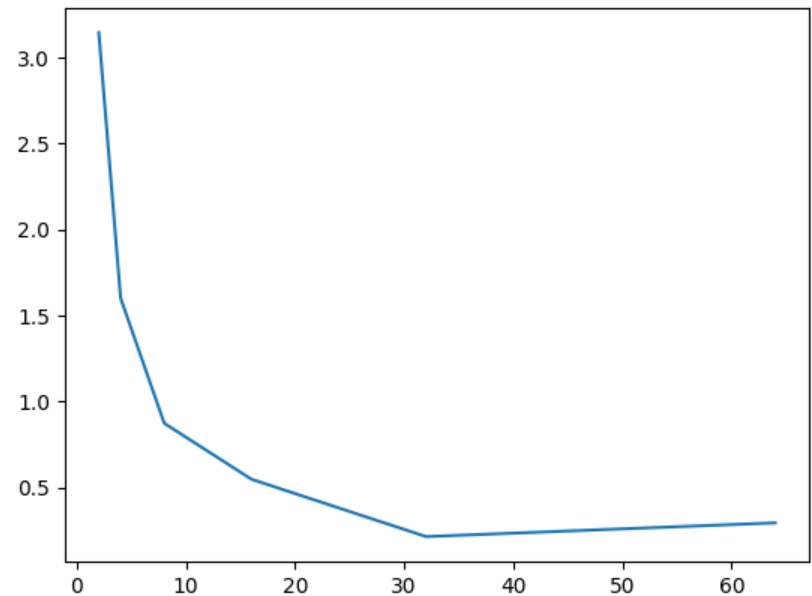
Processors	Time
2	0.005897
4	0.002780
8	0.001521
16	0.013540
32	0.025477



Results

Time taken for different number of processors for a graph with 1000 vertices and 20000 edges.
Lowest: 32 nodes.

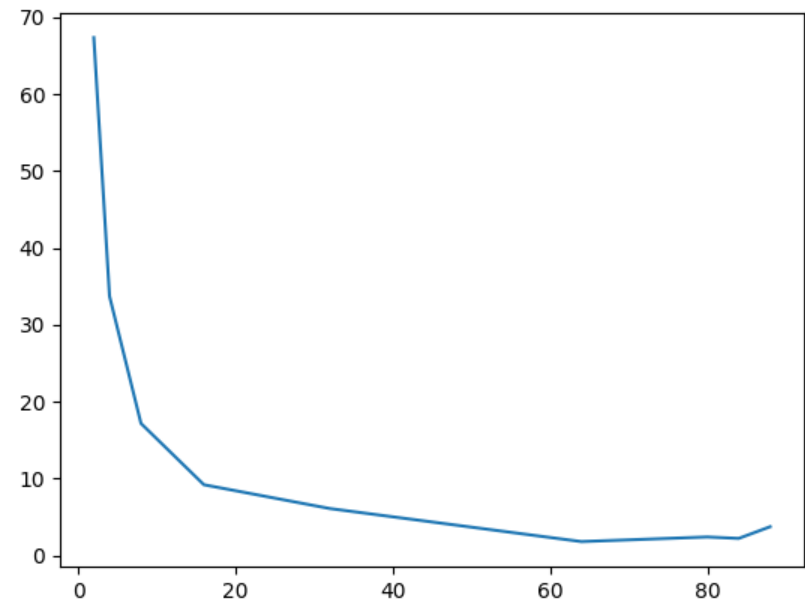
Processors	Time
2	3.145767
4	1.601854
8	0.874434
16	0.548852
32	0.215619
64	0.295514



Results

Time taken for different number of processors for a graph with 10000 vertices and 200,000 edges.
 Lowest: 64 nodes.

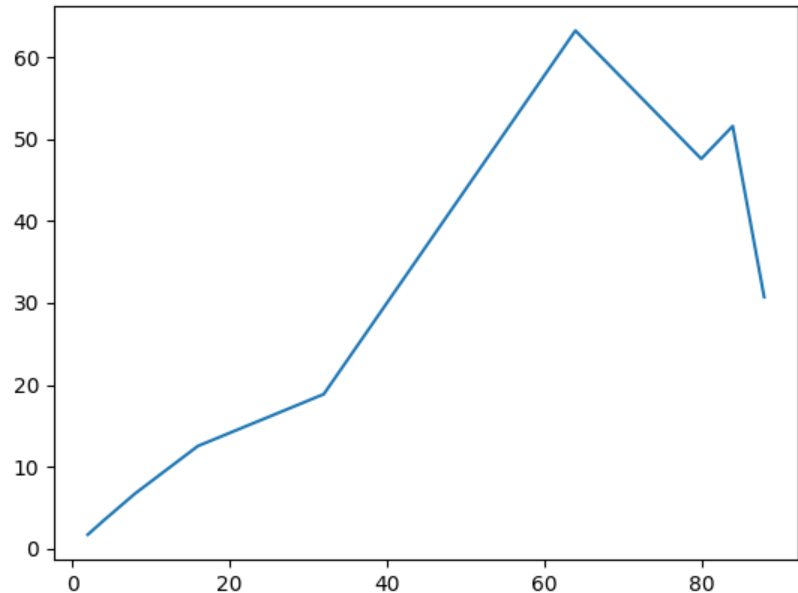
Processors	Time
1	115.460798
2	67.363267
4	33.729089
8	17.190452
16	9.219321
32	6.125911
64	1.826219
80	2.42698
82	2.23899
84	3.76102



Speedup

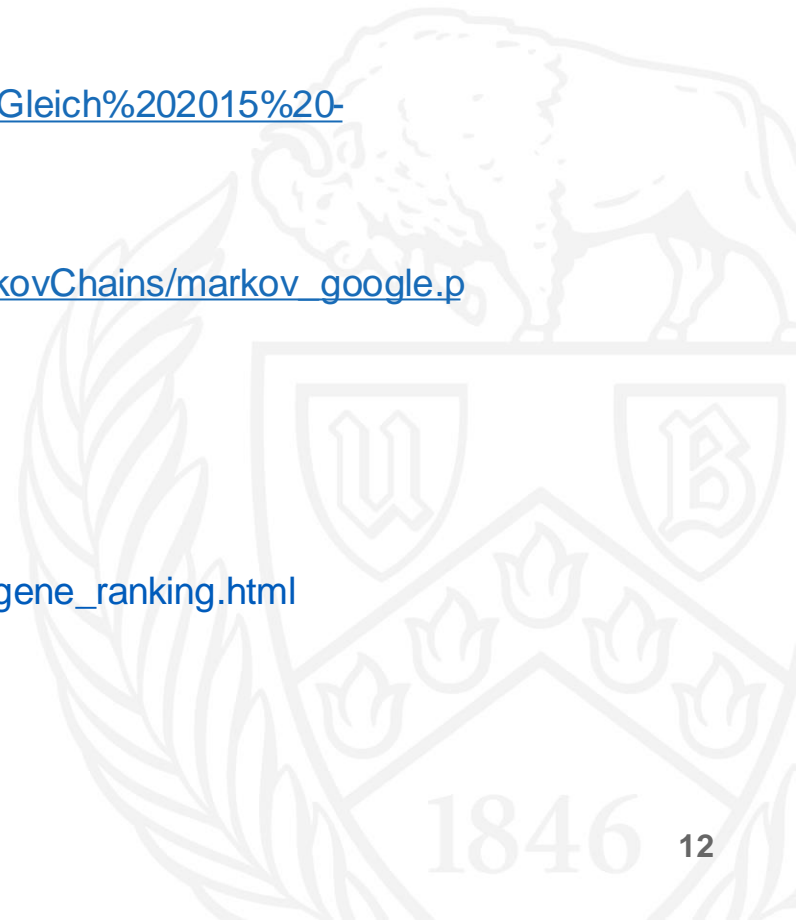
The speedup is the ratio of time taken sequentially vs parallel for the same program.

$$\text{Speedup} = T_{\text{sequential}} / T_{\text{parallel}}$$



References

1. <https://www.cis.upenn.edu/~mkearns/teaching/NetworkedLife/pagerank.pdf>
2. <https://en.wikipedia.org/wiki/PageRank>
3. <https://www.cs.purdue.edu/homes/dgleich/publications/Gleich%202015%20-%20prbeyond.pdf>
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QUESTIONS?

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