

# FACTORIZATION OF A LARGE NUMBER

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# Outline

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- ▣ Expected Values
  - Input; Output;
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- ▣ Solution (Executed on Device)
- ▣ Distributed Solution (Multiple Devices)
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# Goal

- ▣ To find the optimal amount of computation a thread should compute when using CUDA devices
- ▣ Idea (Steps)
  - 1: Implement an easy solution that requires minimal computation
    - ▣ The problem must be scalable relatively easy with respect to threads
  - 2: Run tests for a set number of threads with varying input size
  - 3: Change thread count and repeat step 2 until complete range of thread domain is exhausted
  - 5: Analyze results for conclusion

# Problem Set

- ▣ Factoring a number
  - Simple task but can get very time consuming once the number becomes very large
  - If the number only contains two factors then you can assume it is the product of two prime numbers
- ▣ For cryptography it's essential to use a one-way, or trapdoor mathematical function.
  - A mathematical function that's easy to do in one direction but very difficult, or impossible to reverse.
- ▣ Factoring of prime numbers
  - Easy to find product of two large prime numbers.
  - Difficult to factor large product to two prime numbers.
  - Very large prime number used, because larger the prime number, the more difficult factoring becomes.

# Expected Values

- ▣ Input
  - Very large integer
- ▣ Output
  - Array of values which are factors of the input

# Plan of Attack

- ▣ Brute force solution
  - Iterate over all values in range
  - Do Modulus operation to test for factor
  - Remember proper factors
  - Return results
- ▣ Data Structure Storage (Results)
  - Array of integers that hold the results
  - Size of array is  $\text{sqrt}(\text{"input number"}) * 2$ 
    - Array multiplied by two to hold pair of the factor
  - Each successful modulus operation sets appropriate location in array to integer value found
    - Also sets adjacent value of array which is offset by  $\text{sizeCount}(\text{number of iterations required per device})$

# Array Data Structure

- In this example assuming a device had to compute Factors for the number 20
  - sizeCount would be  $\text{sqrt}(20) + 1$  which is 5
  - factor found at Index = 1 would mean it needs to set its adjacent factor at (Index + sizeCount) which would be position 6 in the array.
    - Ex.  $\text{Array}[1] * \text{Array}[6] = 20$  or  $10 * 2 = 20$

Index

0 1 2 3 4 5 6 7 8 9

Value

1 2 0 4 5 20 10 0 5 4

# Solution (Executed on Device)

- ▣ Input number N
- ▣ Each Iteration
  - Perform Modulus operation for each index on N looking for resulting 0 value and set according Factor position
  - Device's starting number is set according to starting location passed in plus it's own threadId.x
    - ▣ Ex.  $\text{tempNumber} \% N = \text{start} + \text{threadId.x}$
  - If  $((\text{tempNumber} \% N) == 0)$   $\text{Array}[\text{Index}] = \text{tempNumber} \% N$   
 $\text{Array}[\text{Index} + \text{sizeCount}] = N / \text{tempNumber}$





# Solution (Executed on Device)

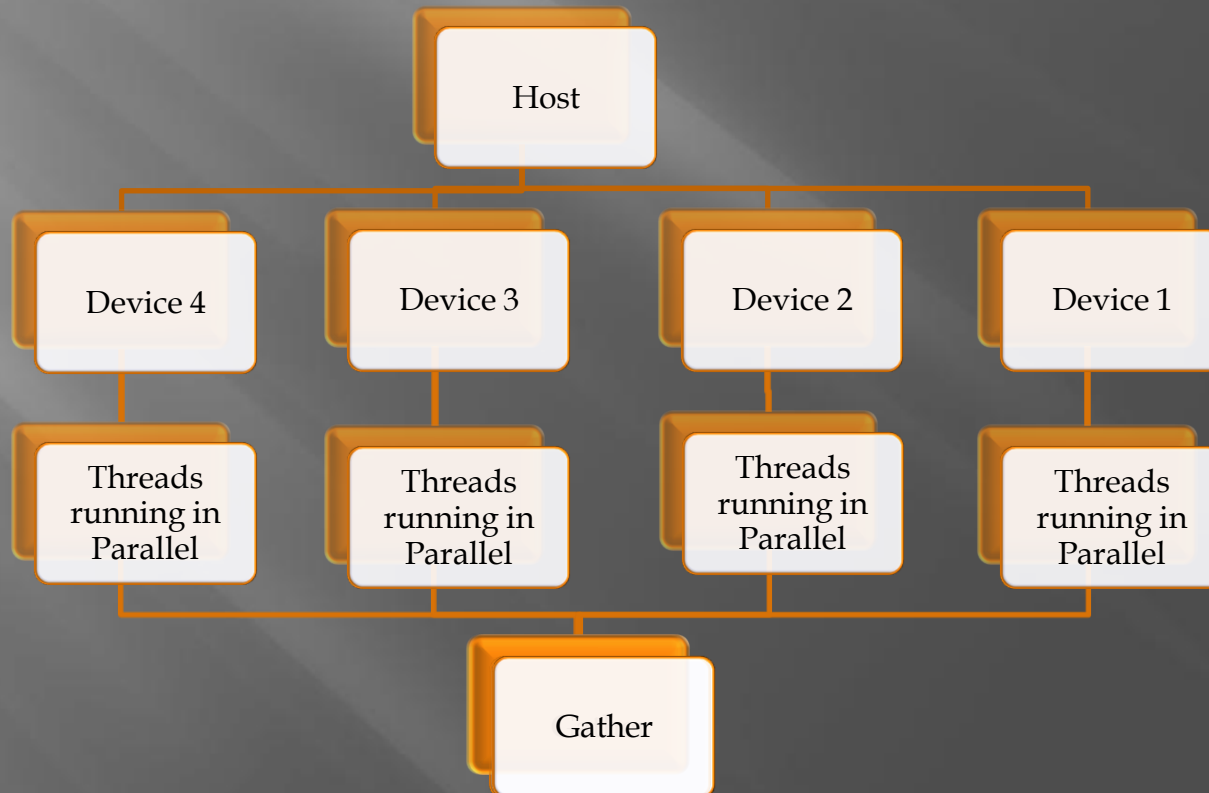
```
__global__ static void kernel(int num,int start,int sizeCount,int numThreads,int
*value)
{
    int tx = threadIdx.x;
    int insertPosition = tx;

    int tempNum = tx + start;

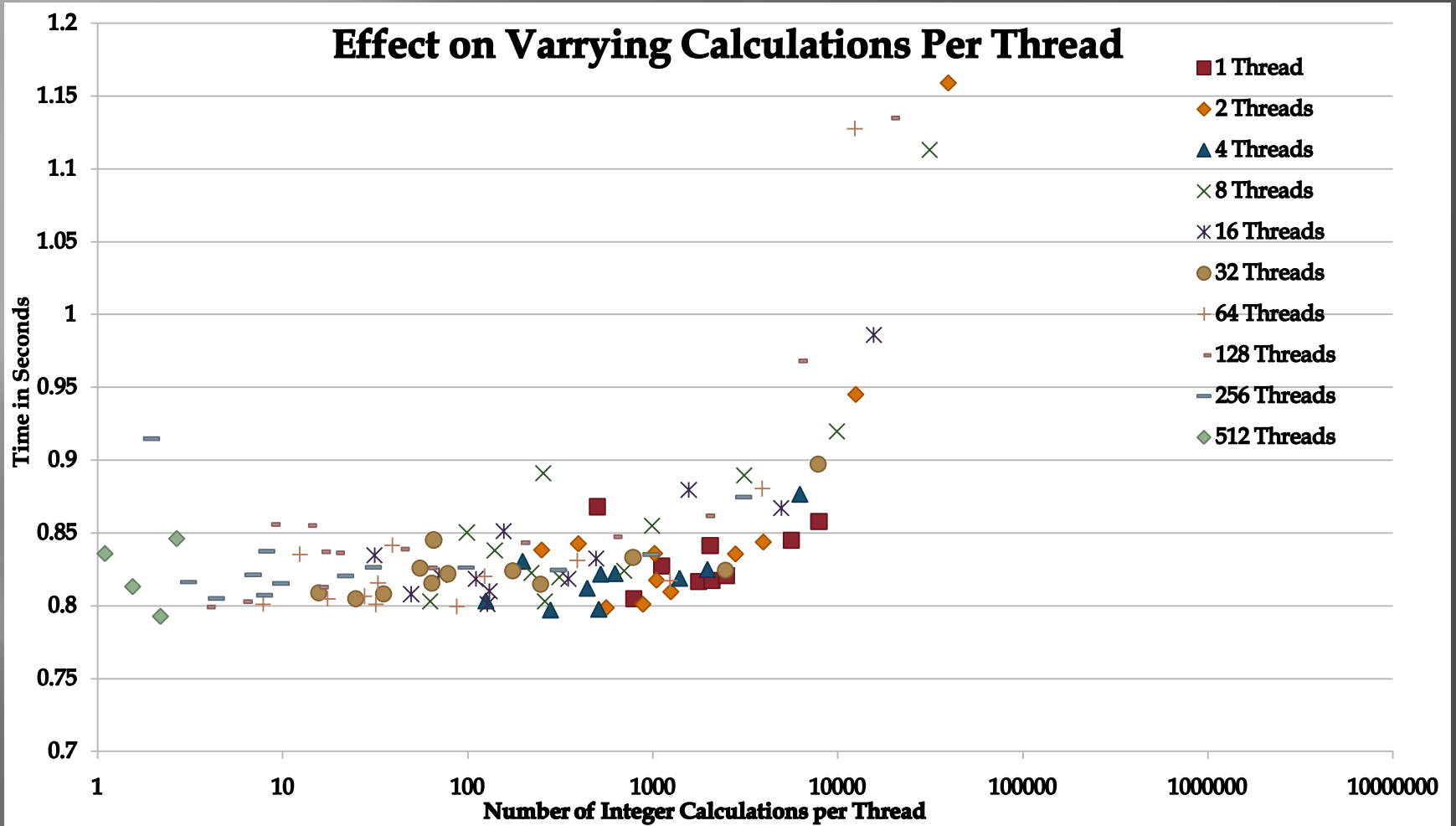
    while( insertPosition < sizeCount ){
        if((num%tempNum) == 0){
            value[insertPosition] = tempNum;
            value[insertPosition+sizeCount] = num/tempNum;
        }
        insertPosition = insertPosition + numThreads;
        tempNum = tempNum + numThreads;
    }
}
```

# Distributed Solution (Multiple Devices)

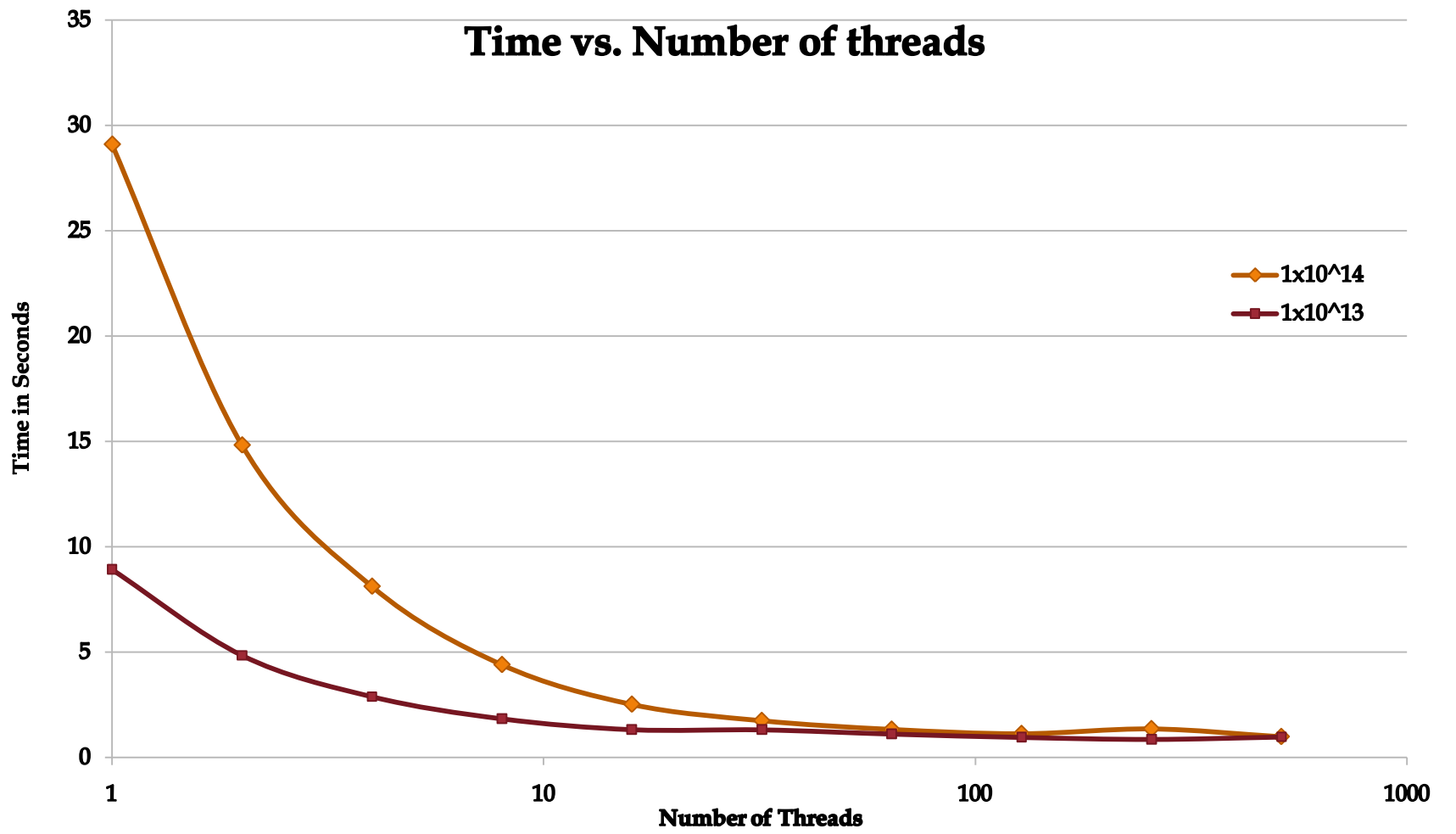
- ▣ Host cudaMalloc's on all CUDA devices
- ▣ Split Array by Device then between each process
  - ▣ Amount of Iterations =  $(\text{sqrt}(\text{NUM}) + 1) / \text{devCount}$



# Results

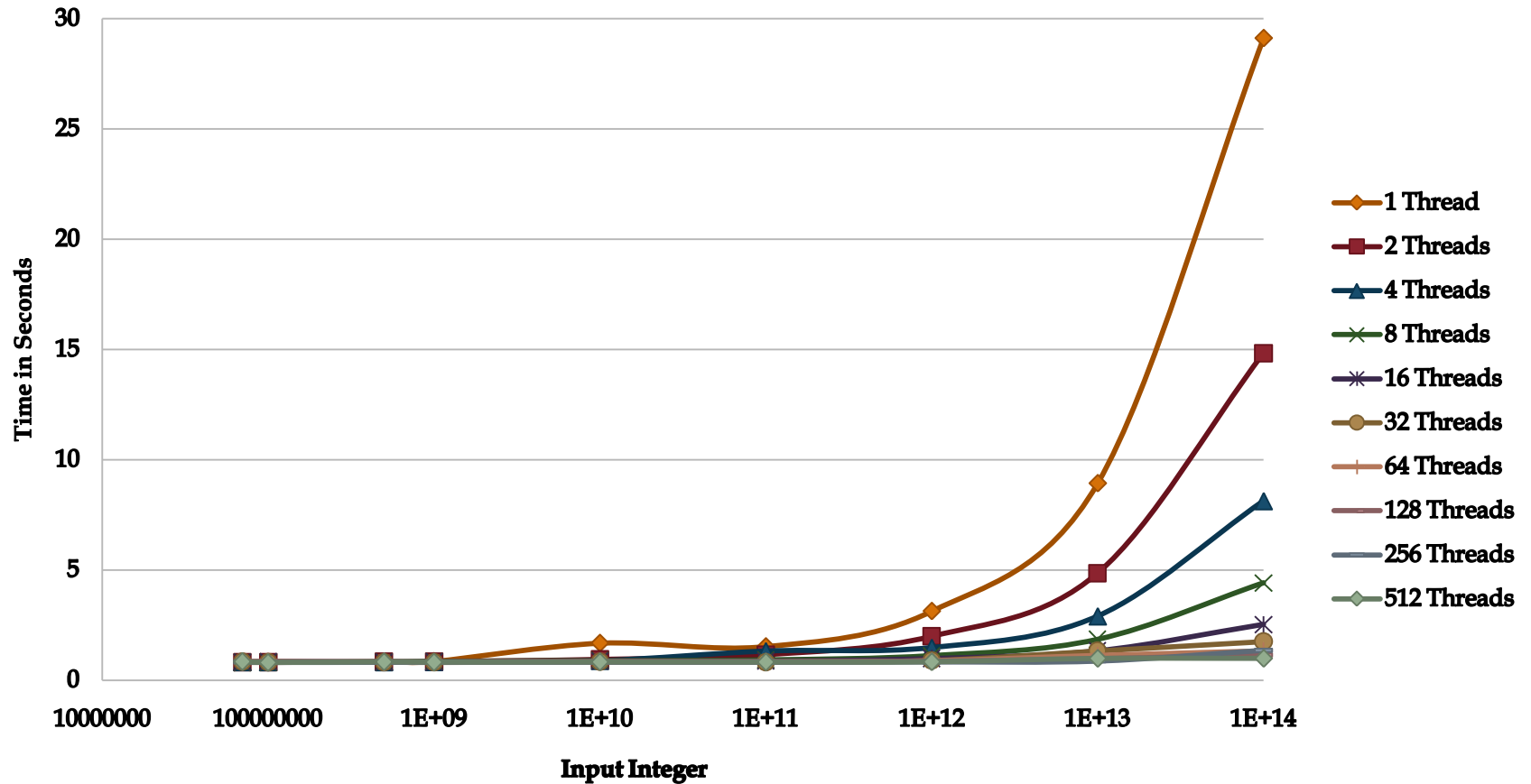


# Results



# Results

## Varying Integer Input Size



# Conclusion

- ▣ Thread creation is very minimal
- ▣ Most time spent in device initialization
- ▣ Threads computing up to 1000 computations seems optimal with including device initialization for timing
  
- ▣ Further research
  - Time kernel execution to obtain direct relation to spawning threads with excluding device time