CSE 421/521 - Operating Systems
Fall 2011

LECTURE - XVI

VIRTUAL MEMORY - II

Tevfik Koşar

University at Buffalo
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Roadmap

- Virtual Memory
  - Page Replacement Algorithms
    - Optimal Algorithm
    - Least Recently Used (LRU)
    - LRU Approximations
    - Counting Algorithms
  - Allocation Policies
  - Thrashing
  - Working Set Model

FIFO

- FIFO is obvious, and simple to implement
  - when you page in something, put it on the tail of a list
  - evict page at the head of the list
- Why might this be good?
  - maybe the one brought in longest ago is not being used
- Why might this be bad?
  - then again, maybe it *is* being used
  - have absolutely no information either way
- In fact, FIFO's performance is typically lousy
- In addition, FIFO suffers from Belady's Anomaly
  - there are reference strings for which the fault rate *increases* when the process is given more physical memory

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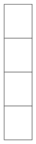
Optimal Algorithm

- Replace page that will not be used for the longest time in future
- 4 frames example

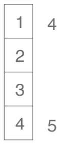
1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



Optimal Algorithm

- Replace page that will not be used for longest period of time
- 4 frames example

1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



6 page faults

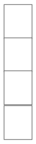
- How would you know this in advance?

Optimal (Belady's) Algorithm

- Provably optimal: lowest fault rate (remember SJF?)
  - evict the page that won't be used for the longest time in future
  - problem: impossible to predict the future
- Why is Belady's Optimal algorithm useful?
  - as a yardstick to compare other algorithms to optimal
    - if Belady's isn't much better than yours, yours is pretty good
      - how could you do this comparison?
- Is there a best practical algorithm?
  - no; depends on workload
- Is there a worst algorithm?
  - no, but random replacement does pretty badly
    - there are some other situations where OS's use near-random algorithms quite effectively!

Least Recently Used (LRU)

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



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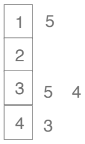
Least Recently Used (LRU)

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



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Least Recently Used (LRU)

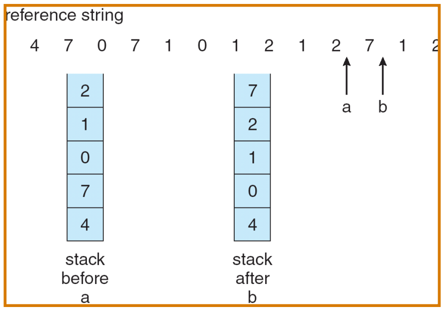
- LRU uses reference information to make a more informed replacement decision
  - idea: past experience gives us a guess of future behavior
  - on replacement, evict the page that hasn't been used for the longest amount of time
    - LRU looks at the past, Belady's wants to look at future
    - *How is LRU different from FIFO?*
- Implementation
  - to be perfect, must grab a timestamp on every memory reference, put it in the PTE, order or search based on the timestamps …
  - way too costly in memory bandwidth, algorithm execution time, etc.
  - so, we need a cheap approximation …

LRU Implementations

- **Counter implementation** (Needs hardware assistance)
  - Every page entry has a counter; every time page is referenced through this entry, copy the clock into the counter
  - When a page needs to be changed, look at the counters to determine which are to change

- **Stack implementation** - keep a stack of page numbers in a double link form:
  - Page referenced:
    - move it to the top
    - requires 6 pointers to be changed
  - No search for replacement

Use Of A Stack to Record The Most Recent Page References



LRU Approximation Algorithms

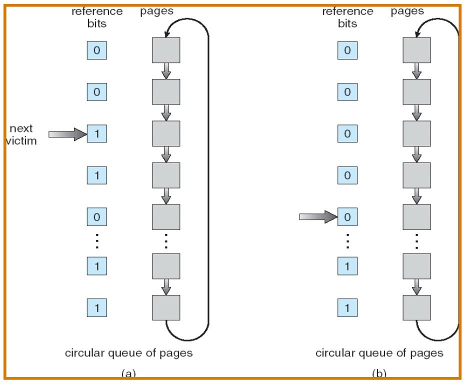
- Reference bit
  - With each page associate a bit, initially = 0
  - When page is referenced bit set to 1
  - Replace the one which is 0 (if one exists). We do not know the order, however.

- Additional Reference bits
  - 1 byte for each page: eg. 00110011
  - Shift right at each time interval

LRU Clock Algorithm

- AKA Not Recently Used (NRU) or Second Chance
  - replace page that is “old enough”
  - logically, arrange all physical page frames in a big circle (clock)
    - just a circular linked list
  - a “clock hand” is used to select a good LRU candidate
    - sweep through the pages in circular order like a clock
    - if ref bit is off, it hasn’t been used recently, we have a victim
      - so, what is minimum “age” if ref bit is off?
    - if the ref bit is on, turn it off and go to next page
  - arm moves quickly when pages are needed
  - low overhead if have plenty of memory
  - if memory is large, “accuracy” of information degrades
    - add more hands to fix

Second-Chance (clock) Page-Replacement Algorithm



Counting Algorithms

- Keep a counter of the number of references that have been made to each page
- **LFU Algorithm**: replaces page with smallest count
- **MFU Algorithm**: based on the argument that the page with the smallest count was probably just brought in and has yet to be used

Allocation of Frames

- Each process needs *minimum* number of pages
- Two major allocation schemes
  - fixed allocation
  - priority allocation

Fixed Allocation

- Equal allocation – For example, if there are 100 frames and 5 processes, give each process 20 frames.
- Proportional allocation – Allocate according to the size of process

s_i = size of process p_i

$S = \sum s_i$

m = total number of frames

a_i = allocation for $p_i = \frac{s_i}{S} \times m$

$m = 64$

$s_i = 10$

$s_2 = 127$

$a_1 = \frac{10}{137} \times 64 \approx 5$

$a_2 = \frac{127}{137} \times 64 \approx 59$

Priority Allocation

- Use a proportional allocation scheme using priorities rather than size
- If process P_i generates a page fault,
  - select for replacement one of its frames
  - select for replacement a frame from a process with lower priority number

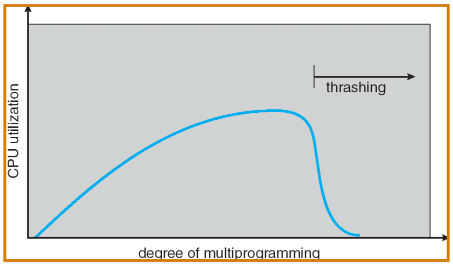
Global vs. Local Allocation

- **Global replacement** – process selects a replacement frame from the set of all frames; one process can take a frame from another
- **Local replacement** – each process selects from only its own set of allocated frames

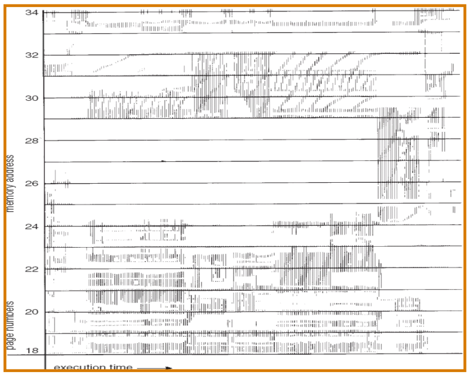
Thrashing

- If a process does not have “enough” frames, the page-fault rate is very high. This leads to:
  - Replacement of active pages which will be needed soon again
  - ➔ **Thrashing** ≡ a process is busy swapping pages in and out
- Which will in turn cause:
  - low CPU utilization
  - operating system thinks that it needs to increase the degree of multiprogramming
  - another process added to the system

Thrashing (Cont.)



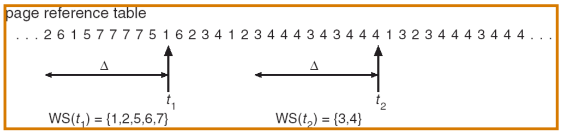
Locality in a Memory-Reference Pattern



Working-Set Model

- $\Delta \equiv$ working-set window ≡ a fixed number of page references
  Example: 10,000 instruction
- WSS_i (working set of Process P_i) =
  total number of pages referenced in the most recent Δ (varies in time)
  - if Δ too small will not encompass entire locality
  - if Δ too large will encompass several localities
  - if $\Delta = \infty \Rightarrow$ will encompass entire program
- $D = \Sigma\, WSS_i \equiv$ total demand frames
- if $D > m \Rightarrow$ Thrashing
- Policy if $D > m$, then suspend one of the processes

Working-set model



Summary

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    - Optimal Algorithm
    - Least Recently Used (LRU)
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  - Working Set Model

- Next Lecture: Project 2 & 3 Discussion
- Reading Assignment: Chapter 9 from Silberschatz.

Acknowledgements

- “Operating Systems Concepts” book and supplementary material by A. Silberschatz, P. Galvin and G. Gagne
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