

Disk Scheduling

Based on the slides supporting the text

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User-Space I/O Software

The diagram shows five layers of the I/O system, each with associated functions:

- User processes:** Make I/O call; format I/O; spooling
- Device-independent software:** Naming, protection, blocking, buffering, allocation
- Device drivers:** Set up device registers; check status
- Interrupt handlers:** Wake up driver when I/O completed
- Hardware:** Perform I/O operation

Arrows indicate the flow of I/O requests from user processes down to hardware, and I/O replies from hardware back up to user processes.

Layers of the I/O system and the main functions of each layer

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Disk Structure

- Disk drives are addressed as large 1-dimensional arrays of *logical blocks*, where the logical block is the smallest unit of transfer.
- The 1-dimensional array of logical blocks is mapped into the sectors of the disk sequentially.
 - Sector 0 is the first sector of the first track on the outermost cylinder.
 - Mapping proceeds in order through that track, then the rest of the tracks in that cylinder, and then through the rest of the cylinders from outermost to innermost.

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Disk Scheduling

- The operating system is responsible for using hardware efficiently — for the disk drives, this means having a fast access time and disk bandwidth.
- Access time has two major components
 - *Seek time* is the time for the disk are to move the heads to the cylinder containing the desired sector.
 - *Rotational latency* is the additional time waiting for the disk to rotate the desired sector to the disk head.
- Minimize seek time
- Seek time \approx seek distance
- Disk bandwidth is the total number of bytes transferred, divided by the total time between the first request for service and the completion of the last transfer.

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Disk Scheduling (Cont.)

- Several algorithms exist to schedule the servicing of disk I/O requests.
- We illustrate them with a request queue (0-199).

98, 183, 37, 122, 14, 124, 65, 67

Head pointer 53

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FCFS

The diagram illustrates the First-Come, First-Served (FCFS) disk scheduling algorithm. A request queue is shown with the following cylinder numbers: 0, 14, 37, 53, 65, 67, 98, 122, 124, 183, 199. The head starts at cylinder 53. The path of the head is shown as a series of arrows: 53 → 98 → 183 → 124 → 122 → 67 → 65 → 37 → 14 → 0. The total head movement is 640 cylinders.

Illustration shows total head movement of 640 cylinders.

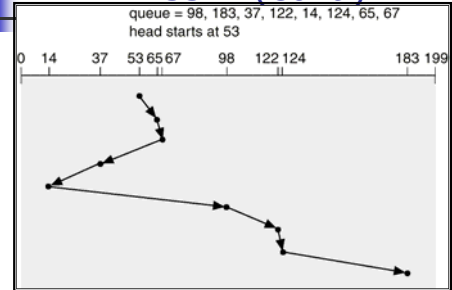
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SSTF

- Selects the request with the minimum seek time from the current head position.
- SSTF scheduling is a form of SJF scheduling; may cause starvation of some requests.
- Illustration shows total head movement of 236 cylinders.

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SSTF (Cont.)



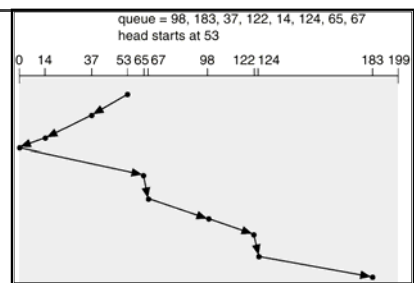
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SCAN

- The disk arm starts at one end of the disk, and moves toward the other end, servicing requests until it gets to the other end of the disk, where the head movement is reversed and servicing continues.
- Sometimes called the *elevator algorithm*.
- Illustration shows total head movement of 208 cylinders.

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SCAN (Cont.)



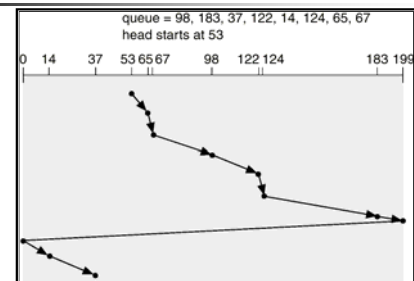
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C-SCAN

- Provides a more uniform wait time than SCAN.
- The head moves from one end of the disk to the other, servicing requests as it goes. When it reaches the other end, however, it immediately returns to the beginning of the disk, without servicing any requests on the return trip.
- Treats the cylinders as a circular list that wraps around from the last cylinder to the first one.

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C-SCAN (Cont.)



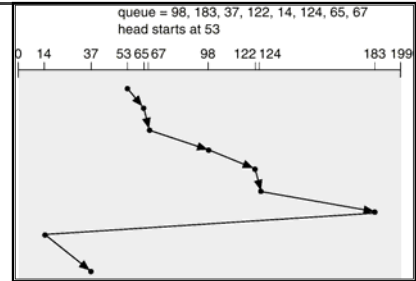
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C-LOOK

- Version of C-SCAN
- Arm only goes as far as the last request in each direction, then reverses direction immediately, without first going all the way to the end of the disk.

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C-LOOK (Cont.)



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Selecting a Disk-Scheduling Algorithm

- SSTF is common and has a natural appeal
- SCAN and C-SCAN perform better for systems that place a heavy load on the disk.
- Performance depends on the number and types of requests.
- Requests for disk service can be influenced by the file-allocation method.
- The disk-scheduling algorithm should be written as a separate module of the operating system, allowing it to be replaced with a different algorithm if necessary.
- Either SSTF or LOOK is a reasonable choice for the default algorithm.

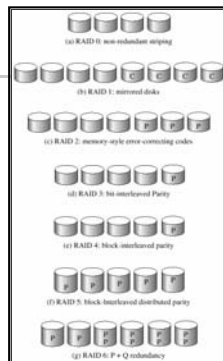
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RAID

- RAID** – multiple disk drives provides **reliability** via **redundancy**.
 - Redundant Array of Independent Disks (RAID)
 - RAID is arranged into six different levels.
 - Several improvements in disk-use techniques involve the use of multiple disks working cooperatively.
 - Disk striping uses a group of disks as one storage unit.
 - RAID schemes improve performance and improve the reliability of the storage system by storing redundant data.
 - Mirroring* or *shadowing* keeps duplicate of each disk.
 - Block interleaved parity* uses much less redundancy.

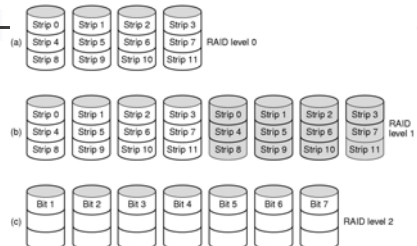
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RAID Levels



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Raid Levels(continued)



- Raid levels 0 through 2
- Backup and parity drives are shaded

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Raid Levels (continued)



- Raid levels 3 through 5
- Backup and parity drives are shaded

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Raid Levels 6 and 7

- RAID 6:P and Q redundancy
- RAID 7: heterogeneous disks array

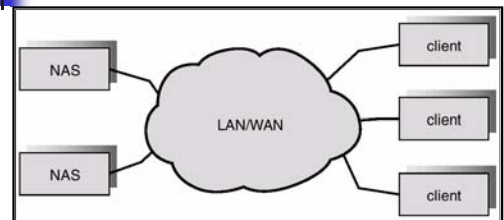
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Disk Attachment

- Disks may be attached one of two ways:
 1. **Host attached** via an I/O port
 2. **Network attached** via a network connection

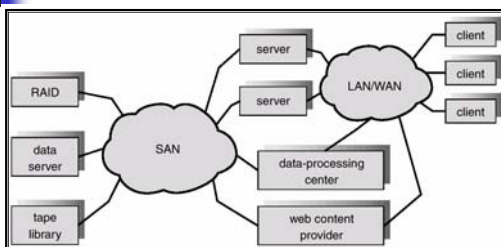
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Network-Attached Storage(NAS)



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Storage-Area Network (SAN)



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Summary

- We looked various disk scheduling methods to optimize secondary storage access.
- We also studies basic principles of RAID.
- NAS and SAN are two recent advances in secondary storage.

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