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A NET STRUCTURE BASED RELATIONAL QUESTION ANSWERER: DESCRIPTION AND EXAMPLES.*

by

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Abstract

A question answering system is described which uses a net structure for storage of information. The net structure consists of nodes and labelled edges, which represent relations between the nodes. The labels are also nodes, and therefore definitions of relations may be stored in the net. It is demonstrated that the generality and complexity of this memory structure allows a surprisingly powerful question answering system to be constructed using comparatively simple executive routines. Output from the question answerer, which is currently running on an interactive, time sharing system, is included, showing its range of applicability including question answering, inductive and deductive inference, simple theorem proving and problem solving.

Key words and phrases: relational question answering, question answering, memory net, memory structure, data structure, semantic memory, semantic information retrieval, deductive inference, inductive inference, problem solving, concept formation, relational logic, learning, theorem proving, fact retrieval.

1. Introduction

Our main research interest has been in the organization of data structures for question answering systems, systems that retrieve facts and have deductive and inductive capabilities to derive new information from the facts explicitly given them. Our two main aims have been to maintain as much generality as possible so that no additional programming be needed regardless of the domain of knowledge for which the system is used and to put as much question answering power as possible into the memory structure itself rather than in the executive routines. This latter aim supports the first in that it would allow special instructions for particular domains to be entered into the memory in the same way as any other information.

SAMENIAQ II, the system described in this paper represents progress toward reaching these aims. Further progress is being made in a later system (see Section 5).

SAMENIAQ II is based upon binary relations. This was a natural starting point because of the generality of binary relations, and the fact that they provide a reasonable test environment for our ideas. Since our major interest is the memory structure, we have not used natural language input, thus avoiding the attendant problems. Instead, all statements input to the system are in the form $x R y$. We hope it will become evident that even with this restriction to binary relations and with basically simple executive routines the system attains a surprising amount of power and range of applicability. This derives from the following characteristics of the system:

1. The memory is a net structure, with the relations serving as labels on directed edges. Each statement $x R y$ is also stored in the converse form $y R(x)$. So that all the information about a name is reachable from the node in the net which represents it.

2. The relations, though used as labels on the edges, are actually also nodes themselves, so information about them may be stored in the memory structure. The major use of this capability is to define a relation in terms of other relations.

3. The system has the ability to use the information stored about a relation when searching memory. Such information may be entered at any time and in the same manner as any other type of data or it may be constructed and entered by the system itself. In our system, any relation may be used as an undefined term, may be defined in terms of other relations, or may be defined recursively. Any single relation may be used in any or all of these ways. Complex relations may be built out of simple relations using the relative product operation and node restrictions to restrict the domains or ranges of the simple relations. Thus, quite complicated relations may be defined.

The following sections give more detailed information about and examples of the SAMENIAQ II

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system. The final section describes a later system which is being developed to satisfy more completely the goals discussed above.

2. Implementation and Operation

SAMENLAQ II is a revision of SAMENLAQ, "A Semantic Association Memory Net that Learns and Answers Questions". Both programs are written in SNOSOL as interactive question answering systems, but SAMENLAQ II, unlike SAMENLAQ which was run in batch mode with simulated interaction on a CDC 3600, is fully interactive and is currently running under the University of Wisconsin B5500 time sharing system. SAMENLAQ II differs from SAMENLAQ in that it provides aids to the user, who inputs data directly via a teletype, allows for storage of input files and memories in disk files, and, most importantly, allows for recursive definitions of relations and allows the user to control in real time how much effort the system should spend searching its memory to discover more information for use in answering a question.

Figures 2a and 2b depict the overall flow of control in SAMENLAQ II. Figure 1 demonstrates its operation. At the top level of operation, three types of input are allowable: statements, questions and requests to the system executive. [B1,2a]* Representative examples are given in fig. 1 by lines "a", "f", and "b" respectively. Notice that user inputs are indicated by terminal left arrows whereas unterminated lines indicate SAMENLAQ II responses. Simple statements such as line "c" result in the construction of a net substructure containing the nodes "BOSTON", "EAST. OF", "EAST. OF(CNV)", and "WORCESTER" in which "BOSTON" and "WORCESTER" are tied together via the "EAST. OF" node and "WORCESTER" and "BOSTON" are tied together via the "EAST. OF(CNV)" node. This structure is considered in more detail in the next section. More complicated statements such as line "d" are interpreted by the system as a series of simple statements. The system provides the user with various types of feedback, some of which may be turned off by appropriate requests to the executive.[B7,2a] Lines prior to "f" demonstrate input in the full and limited response modes. Line "b" requests that the full response mode be turned off. Requests are identified by a terminating "Ea". Other request options are indicated in fig. 2.[B2,2a]

Several relation words are built into the system. "MEMBER" allows a particular net search to be limited to a subclass of all the relations represented in the net. Subclass definition may take place at any point during a conversation and is determined solely by the user. Line "d" represents the introduction of the subclass "COMPASS.RELS". Such classes are useful for handling questions involving paths in the net which connect prescribed nodes.

"IMPLYB" allows a given relation to be defined in terms of other relations. It is one of the most important features of SAMENLAQ II. The system has the ability to utilize "IMPLYB" information about a relation during the question answering process by using it as a generalization of substitution rule.

Line "e" demonstrates the use of IMPLYB to introduce NORTH_OF(CNV) as an acceptable replacement for SOUTH_OF. The line also results in SOUTH_OF(CNV) IMPLYB NORTH_OF being incorporated in memory.

To enhance readability, the system allows the user to introduce his own interrogatives by means of the form "x IS QUESTION".

Questions are terminated by "*". There are four possible types -- one verification type (x R y *) and three fill in the blank types (x R _ *), (R y *) and (x _ y *).[B3,2a]. Line "e" illustrates the x R _ * type, line "j" the X _ Y* type and line "k" the _ R Y* type. Notice that non-simple relations can be handled by the system. There are three types of relations used in SAMENLAQ II, simple, compound and complex. A simple relation is of the form R or R(CNV) where the character string R is not meaningfully decomposable and Y R(CNV) X if and only if X R Y. A compound relation is a simple relation or a relative product of simple relations and is of the form R1/R2/R3, (or sometimes, as a stylistic variant, R1/R2/R3). X R1/R2/R3 Y holds if and only if there exists some z and w such that X R1 z, z R2 w and w R3 Y. A complex relation is a compound relation or a compound relation with node restrictions. A node restriction is of the form (R - B) where B is a simple relation and B is a string of names, each one followed by a comma. Range restrictions are used to restrict the domain or range of a simple relation which forms part of the complex relation. For example, (R - B)R1 is the relation R1 with a restricted domain, and R1(R - B) is R1 with a restricted range.

Domain of R3 Domain of R6
x (R1 - B1)(R2 - B2)R3(R4 - B3)(R5 - B4)R6(R7 - B5) y

A name, x, satisfies the node restriction (R - B) if for all y in the string B xRy is explicitly stored in memory. Thus the relational statement above holds if x satisfies (R1 - B1) and (R2 - B2), y satisfies (R7 - B5) and there exists some z such
Fig. 1 A Conversation With SAMENLAQ II.
OTHERWISE TYPE AN INTEGER FOR FURTHER SEARCH
OR - LIST TROUBLES OR AN INPUT STATEMENT OR - EXAMPLES
OK--
WAIT...

1 ANSWER - WORCESTER
   AND ALBANY
   AND BUFFALO
   READY

j BUFFALO HAS WHAT RELATIONSHIP TO NEW YORK?
   ENTER COMMA LIST OF RELATION CLASSES TO BE USED OR "ANY"
   ANY--
   BUFFALO IN NEW YORK STATE IN (CNV) NEW YORK
   ENTER - OK OR AN INTEGER INDICATING NUMBER OF ADDITIONAL PATH LINKS.
   1--
   WAIT...
   BUFFALO EAST OP (CNV) ALBANY NORTH OP POUGHKEEPSIE NORTH OP
   NEW YORK
   ENTER - OK OR AN INTEGER INDICATING NUMBER OF ADDITIONAL PATH LINKS.
   OK--
   READY
   FLAG ASK OFF
   READY
   FLAG TRACE OFF
   READY

k WHAT (IN NEW YORK STATE ) SOUTHWEST OP BOSTON?
   I FIGURE BOSTON SOUTHWEST OP (CNV) / (IN NEW YORK STATE ) UNKNOWN
   WHAT NOW?
   LIST TROUBLES--
   AT THE FOLLOWING NAMES COULD NOT APPLY THE LISTED RELATIONS.
   BOSTON - SOUTHWEST OP (CNV)
   WHAT NOW?
   SOUTHWEST OP IMPLY SOUTH OP WEST OP WEST OP SOUTH OP--
   STATEMENT - SOUTHWEST OP IMPLY SOUTH OP WEST OP WEST OP SOUTH OP
   OK?...YES--
   WAIT...
   WHAT NOW?
   3--
   I FIGURE BOSTON SOUTHWEST OP (CNV) / (IN NEW YORK STATE ) POUGHKEEPSIE
   WHAT NOW?
   SOUTH OP IMPLY SOUTH OP SOUTH OP--
   STATEMENT - SOUTH OP IMPLY SOUTH OP SOUTH OP
   OK?...YES--
   WAIT...
   WHAT NOW?
   3--
   I FIGURE BOSTON SOUTHWEST OP (CNV) / (IN NEW YORK STATE ) POUGHKEEPSIE
   AND NEW YORK
   WHAT NOW?
   OK--
   ANSWER - POUGHKEEPSIE
   AND NEW YORK
   READY

Fig. 1 (cont.) A Conversation With SAMENLAQ II
Fig. 2a  Overall Flow of Control

*See figure 2b.*
Fig. 2b  Flow Chart of Main Question Answering Routines

*This routine is used at [B4,2a], [B5,2a] and [B6,2a].
that \( z \) satisfies (R4-\( \delta \)) and (R5-\( \delta \)), \( xR3z \), and \( zR6y \).

A complex relation may also consist solely of node restrictions, in which case it is an identity relation on a restricted domain (viz. the set of all names which satisfy all the node restrictions) and is the statement of a conjunctive concept. The executive routines have built into them the ability to deal with compound relations and with compound and complex relations. When a relation serves to label an edge of the net, i.e., in its appearance in the value of a name, it is treated as a simple relation. Compound and complex relations are used when defining other relations and may themselves have definitions. Lines beginning at "k" illustrate these ideas.

Line "f" represents one of the four question types. To answer it, SAMENLAAQ II attempts to apply "EAST,OF" to "BOSTON". Since the system has the statement "BOSTON EAST,OF WORCESTER" represented explicitly in its memory, EAST,OF can be successfully applied - yielding "WORCESTER". Since the system does not now explicitly that "BOSTON EAST,OF ALBANY" or that EAST,OF is a transitive relation, it is incapable of finding further nodes satisfying "BOSTON EAST,OF X". This is illustrated in line "g" where "1" indicates that SAMENLAAQ II is to execute one substitution cycle - i.e. substitute for each relation it is currently attempting to apply to a node, all acceptable replacements contained on the relations IMPLBY list.[B1,2b] (The IMPLBY list for a relation \( R \) is a list of all relations, \( x, \) such that \( R \) IMPLBY \( x \).) In line "h" the system is informed that EAST,OF is transitive. Two additional IMPLBY substitution cycles are then requested by typing in "2".

Figure 3 illustrates an application of SAMENLAAQ to rudimentary conjunctive concept formation. The statements presented in Fig. 3a supply the system with a small data base concerning the relations of various objects in a room. Note, in line "a" two of these are grouped into the class "SPIREL". In Fig. 3b SAMENLAAQ is asked to answer a question concerning a relation it has never seen before. Failing to apply this relation, it is given a series of examples whose relationship to one another is arbitrarily designated CONCEPTI. Restricting its search to the relation class SPIREL, the system obtains all paths of length 2 or less connecting the example pairs. The properties common to each class of nodes at the same level along the path are also calculated. For example (CONSTRUCTION-ELECTRICAL), (CONSTRUCTION-WOOD), and (HAS.PART-ELEVATED.HORIZONTAL.SUPPORTING SURFACE,) are node properties common to the first, second, and third levels respectively. The path connecting the example pairs is "ON, TOP, ELEMENT". This relation is then placed upon the IMPLBY list for CONCEPTI. Finally CONCEPTI is used to define the original concept IS.AN.ELECTRICAL.COMMUNICATION.DEVICE, THAT.IS,ON, TOP, OF, A.PIECE, OF, WOODEN. The system then has sufficient information to answer the question and responds with the correct answer.

3. SAMENLAAQ II Description

Statements entered into the memory are of the form "NAMEI RELATION NAME2" where NAMEI and NAME2 are non-decomposable names and RELATION is a simple relation. Information contained in such statements is stored on parent lists associated with NAMEI and NAME2. Thus, the above statement produces the following paren list for "NAME1;" (RELATION/-1)" where the contents of the slash name "/1" is the comma list "NAME2." In the example below, the paren pair (WEST,OF/-3) on the paren list for WORCESTER indicates that Worcester is West of each of the elements found on the comma list named /3. Note that the value of a slash name is a comma list, similarly the value of a name is its paren list.

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Fig. 3a A Conjunctive Concept Formation Example: Data Base Input.
WHAT IS AN ELECTRICAL COMMUNICATION DEVICE THAT IS ON TOP OF A PIECE OF WOODEN FURNITURE?
I FIGURE FURNITURE IS AN ELECTRICAL COMMUNICATION DEVICE THAT IS ON TOP OF A PIECE OF WOODEN FURNITURE (CNV)
WHAT NOW?
EXAMPLES:
RELATION WHOSE DEFINITION IS TO BE FOUND CONCEPT 1
RELATION CLASS CONCEPT 1 IS A MEMBER OF SPLREL
DEPTH OF ALTERNATE DEFINITIONS 2
TYPE PAIRS X Y SUCH THAT X CONCEPT 1 Y
WHEN FINISHED, TYPE - END.
TEL1 TABLE
RADIO1 DRESSER
SIGNAL, LIGHT1 SHELF
END
OK?...YES
THANK YOU.
WAIT...
FROM THE EXAMPLES YOU HAVE GIVEN ME, I WOULD GUESS THAT CONCEPT 1 IS THE SAME AS (CONSTRUCTION-ELECTRICAL)(USED-FOR-COMMUNICATION,)(ON-TOP-(CONSTRUCTION-WOOD,))ELEMENT/(HAS-PART-ELEVATED, HORIZONTAL, SUPPORTING, SURFACE,)
WHAT NOW?
IS AN ELECTRICAL COMMUNICATION DEVICE THAT IS ON TOP OF A PIECE OF WOODEN IMPLY CONCEPT 1/ELEMENT
STATEMENT
IS AN ELECTRICAL COMMUNICATION DEVICE THAT IS ON TOP OF A PIECE OF WOODEN IMPLY CONCEPT 1/ELEMENT
OK?...YES
WAIT...
WHAT NOW?
2
1 I FIGURE FURNITURE IS AN ELECTRICAL COMMUNICATION DEVICE THAT IS ON TOP OF A PIECE OF WOODEN FURNITURE (CNV) TTY1 AND SIGNAL, LIGHT1 AND TEL1 AND TV1 AND RADIO1
WHAT NOW?
OK
WAIT...
ANSWER - TTY1
AND SIGNAL, LIGHT1
AND TEL1
AND TV1
AND RADIO1
READY

Fig. 3b A Conjunctive Concept Formation Example: Interrogation.
changing the statement to its converse form for storage under the second argument allows all statements about a name to be stored in the same place (the name's value) regardless of whether the name was the first or second argument in the original statements. This contrasts with the methods for retrieving a relational statement from either argument used by the Relational Data File (RDF) and by DEACON. In RDF, statements are stored in one direction, but in different files, ordered on different parts of the statement. Thus to get all information about a single name, one file must be searched exhaustively or the name must be looked up in all files. In DEACON, the statements are stored in the form of closed "connecting rings" through the three parts of the statement. Thus the statement is reachable from any part of it without recourse to several files, but it is impossible to tell from a connecting ring where the statement should begin, i.e., whether the ring through x, R, and y represents the statement xRy, Ryx, or yRx.

The generality of SAMENL AQ derives largely from the ability to introduce new relations at any time, to introduce definitions of new relations or relations that had previously been undefined and to extend the definition of a relation. Definitions may also be added at any time and are stored in the memory net structure just like any other data. Definitions are not given in terms of relation properties that have been built into the system, but in terms of other relations. Nevertheless, various standard relation properties can be dealt with, for example:

1. Entering the statement, "RI IMPLY R1/R1 causes R1 to be transitive.
2. Entering "RI IMPLY R1(CNV)\) causes R1 to be symmetric.
3. It was previously pointed out (section 2) that the complex relation consisting only of a node restriction (R-$\theta$) serves as an identity relation for all names x such that for every name y in the string $\theta$, xRy is explicitly in memory. If $\theta$ were the string consisting only of the delimiter ",", (R-$\theta$) would be the identity relation for all x such that for any y, xRy were explicitly stored. If this were true for all x in memory, (R-$\tau$) would be the universal identity relation. In that case entering, the statement "RI IMPLY (R-$\tau$)" would cause R1 to be reflexive.

Thus, R1 might be defined to be an equivalence relation by entering "RI IMPLY (MEMBER-$\tau$), RI(CNV), R1/R1/\). In addition to these forms of definition, a relation may be defined a) by using only other relations, and b) by combining other relations along with the relation being defined thus forming a general recursive definition. A wide range of relations may thus be used without programming them into the executable routines as it is done in Raphael's SIR. This method of defining relations also contrasts with that used by Elliott in "GRAIS". In "GRAIS" relational properties are built into the executable routines. In fact, this is done in such a way as to provide specific routines for 32 classes of relations. The user introduces a new relation by specifying which of the 32 classes it belongs in and this determines how it will be handled. This does not allow a user to use a relation whose properties he either does not know completely or does not wish to make specific initially. A user may also define a relation in terms of a Boolean function of previously introduced relations. However, these relations may not be stored in the data structure, only used for question answering.

It is interesting to note that the user builds his own logic system into SAMENLAQ II when he specifies IMPLY \(I\) information (used as the rules of inference) and other statements (the axioms). The only logical structure imposed on the user is the metatheoretic substitution rule embodied in the procedures which apply IMPLY and the limits on the form of a rule of inference imposed by the syntax of complex relations. If the user specifies a strange or even self-contradictory "logic" SAMENLAQ II will produce deductions that are equally strange or contradictory: interpretation is in the mind of the user. (For example, in section 4, deductions arising from the relation "IS,PART,OF" only make sense if the interpretation of "x IS,PART,OF y" motivating the rule of inference "IS,PART,OF IMPLY IS,PART,OF/IS(CNV)" is that every member of y has a part which is a member of x and the interpretation of "IS" is "is subset of." )

It would be possible to append an executive to SAMENLAQ II which would constrain the type of logical system to one with certain prescribed properties.

4. Some Varied Applications of SAMENLAQ II

The following sample conversations demonstrate SAMENLAQ II's ability to deal with relations arising from a variety of problem areas.

The first conversation (Figure 4a,b,c) involves a subdialogue from SIR. It demonstrates SAMENLAQ's ability to handle relations such as part, subset, owns and the interdependence between such relations. Although specific relations and relational properties are not built into the system, SAMENLAQ can utilize the information \((x)(\forall)(y)(\exists)(x IS,PART,OF Y & Z<x)\) to construct a statement IS,PART,OF IMPLY IS,PART,OF/IS(CNV). (See comments at the end of section 3.)
IS.PART.OF INLYB IS.PART.OF/IS.PART.OF,IS.PART.OF/IS(CNV)—
READY
SOMETIMES INLYB IS(CNV)—
READY
IS INLYB IS/IS—
READY
NOSE IS.PART.OF PERSON—
READY
NOSTRIL IS.PART.OF NOSE—
READY
PROFESSOR IS TEACHER—
READY
TEACHER IS PERSON—
READY
NOSTRIL IS.PART.OF PROFESSOR*1—
I FIGURE NOSTRIL IS.PART.OF/ NOSE AND PERSON
WHAT NOW?
1—
I FIGURE NOSTRIL IS.PART.OF/ NOSE AND PERSON AND TEACHER
WHAT NOW?
1—
I FIGURE NOSTRIL IS.PART.OF/ NOSE AND PERSON AND TEACHER AND PROFESSOR
WHAT NOW?
OK—
WAIT...
ANSWER - TRUE
READY
PERSON IS LIVING CREATURE—
BAD INPUT. TRY AGAIN.
READY
PERSON IS LIVING,CREATURE—
READY
HAS,AS.PART INLYB IS.PART.OF(CNV)—
READY
LIVING,CREATURE SOMETIMES/HAS,AS.PART NOSTRIL*3—
I FIGURE LIVING,CREATURE SOMETIMES/HAS,AS.PART/ NOSTRIL
WHAT NOW?
OK—
ANSWER - TRUE
READY
CRT IS DISPLAY,DEVICE—
READY
CRT IS.PART.OF B5500—
READY
BRUTUS IS B5500—
READY
SCREEN IS.PART.OF DISPLAY,DEVICE—
READY
SCREEN IS.PART.OF BRUTUS*1—
I FIGURE SCREEN IS.PART.OF/ DISPLAY,DEVICE AND CRT
WHAT NOW?
1—
I FIGURE SCREEN IS.PART.OF/ DISPLAY,DEVICE AND CRT AND B5500
WHAT NOW?
1—
I FIGURE SCREEN IS.PART.OF/ DISPLAY,DEVICE AND CRT AND B5500 AND
BRUTUS
WHAT NOW?

Fig. 4a Learning and Deduction Using Several Relations From SIR
OK--
WAIT...
ANSWER - TRUE
[Answer to question posed at line "a" Fig. 4a.]
READY
OWNS IMPLY IS/OWNS--
READY
FIREMAN OWNS PAIR. OF. RED. SUSPENDERS--
READY
DOCTOR OWNS PAIR. OF. RED. SUSPENDERS*1--
I FIGURE DOCTOR OWNS/ UNKNOWN
WHAT NOW?
OK--
WAIT...
ANSWER - UNKNOWN
READY
FIRECHIEF IS FIREMAN--
READY
FIRECHIEF OWNS PAIR. OF. RED. SUSPENDERS*1--
I FIGURE FIRECHIEF OWNS/ PAIR. OF. RED. SUSPENDERS
WHAT NOW?
OK--
WAIT...
ANSWER - TRUE
READY
EXAMPLE. OF IMPLY EXAMPLE. OF/IS--
READY
A IMPLY EXAMPLE. OF--
READY
STU OWNS LOG. LOG. DECITRIG1--
READY
LOG. LOG. DECITRIG1 EXAMPLE. OF LOG. LOG. DECITRIG--
READY
LOG. LOG. DECITRIG IS SLIDE. RULE--
READY
STU OWNS/A SLIDE RULE*1--
I FIGURE STU OWNS/A/ LOG. LOG. DECITRIG
WHAT NOW?
1--
I FIGURE STU OWNS/A/ LOG. LOG. DECITRIG AND SLIDE. RULE
WHAT NOW?
OK--
ANSWER - TRUE
READY
ENGINEERING. STUDENT OWNS SLIDE. RULE--
READY
GEORGE EXAMPLE. OF TECH. MAN--
READY
TECH. MAN IS ENGINEERING. STUDENT--
READY
GEORGE OWNS/A SLIDE. RULE*1--
I FIGURE GEORGE OWNS/A/ UNKNOWN
WHAT NOW?
1--
I FIGURE GEORGE OWNS/A/ UNKNOWN
WHAT NOW?
OK--
ANSWER - UNKNOWN
READY

Fig. 4b Learning and Deduction Using Several Relations From SIR

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c  ENGINEERING,STUDENT EXAMPLE,OF(CNV)/OWNS/EXAMPLE,OF LOG,LOG,DECITRIG*1-
  I FIGURE ENGINEERING,STUDENT EXAMPLE,OF(CNV)/OWNS/EXAMPLE,OF/ UNKNOWN
  WHAT NOW?

b  STU EXAMPLE,OF TECH,MAN-- [Provide additional information necessary to answer question]
  STATEMENT - STU EXAMPLE,OF TECH,MAN
  OK?...YES--
  WAIT...
  WHAT NOW?
  2--
  I FIGURE ENGINEERING,STUDENT EXAMPLE,OF(CNV)/OWNS/EXAMPLE,OF/
  LOG,LOG,DECITRIG AND SLIDE,RULE
  WHAT NOW?
  OK--
  ANSWER - TRUE [Answer to question posed at line "c".]
  READY

Fig. 4c  Learning and Deduction Using Several Relations From SIR
NEW YORK, LOS ANGELES, SANTA BARBARA, ORLANDO IS CITY—
READY
NEW YORK HAS POPULATION 7781984—
READY
LOS ANGELES HAS POPULATION 2479015—
READY
ORLANDO HAS POPULATION 88135—
READY
SANTA BARBARA HAS POPULATION 58768—
READY
7781984, 2479015 IS GREATER THAN 100000—
READY
88135, 58768 IS LESS THAN 100000—
READY
BOEING 707, BOEING 727, D C 8 IS JET PLANE—
READY
CONVAIR 240 IS PROP PLANE—
READY
JET PLANE, PROP PLANE SUBSET AIRPLANE—
READY
IS IMPLY IS SUBSET—
READY
BOEING 707, D C 8 CARRIES 150—
READY
BOEING 727 CARRIES 120—
READY
CONVAIR 240 CARRIES 50—
READY
150 IS GREATER THAN 120—
READY
50 IS LESS THAN 120—
READY
a IS, LARGE IMPLY (IS CITY,) HAS, POPULATION/(IS, GREATER THAN 100000,)
HAS, POPULATION(CNV)/IS—
READY
b IS, LARGE IMPLY (IS AIRPLANE,) CARRIES/(IS, GREATER THAN 120,) CARRIES
(CNV)/IS—
READY
WHAT IS QUESTION—
WHAT IS AIRPLANE*1—
I FIGURE AIRPLANE IS(CNV)/ BOEING 707 AND BOEING 727 AND D C 8 AND
CONVAIR 240
WHAT NOW?
OK—
WAIT...
ANSWER - BOEING 707
AND BOEING 727
AND D C 8
AND CONVAIR 240
READY
c D C 8 IS, LARGE WHAT*1—
I FIGURE D C 8 IS, LARGE/ JET PLANE AND AIRPLANE
WHAT NOW?
OK—
WAIT...
ANSWER - JET PLANE
AND AIRPLANE
READY

Fig. 5a Application to census, airplane and airline flight data with an ambiguous relation.
WHAT IS LARGE CITY*1 --
I FIGURE CITY IS LARGE(CNV)/ NEW YORK AND LOS ANGELES
WHAT NOW?
OK--
WAIT...
ANSWER - NEW YORK
    AND LOS ANGELES
READY
SANTA BARBARA, ORLANDO HAS POPULATION WHAT*--
I FIGURE SANTA BARBARA, ORLANDO HAS POPULATION/ 58768 AND 88135
WHAT NOW?
OK--
WAIT...
ANSWER - 58768
    AND 88135
READY
FLT. 207 FLIES FROM NEW YORK--
READY
FLT. 207 FLIES TO LOS ANGELES--
READY
DC. 8 USED ON FLT. 207--
READY
FLT. 207 DEPARTS AT 10:00AM--
READY
FLT. 207 ARRIVES AT 12:30PM--
READY
FLT. 308 FLIES FROM NEW YORK--
READY
FLT. 308 FLIES TO ORLANDO--
READY
BOEING 727 USED ON FLT. 308--
READY
FLT. 45 FLIES FROM LOS ANGELES--
READY
FLT. 45 FLIES TO SANTA BARBARA--
READY
CONVAIR 240 USED ON FLT. 45--
READY
FLT. 45 DEPARTS AT 1:30PM--
READY
FLT. 45 ARRIVES AT 2:15PM--
READY
1:30PM IS LATER THAN 12:30PM--
READY
CONNECTS WITH IMPLY ARRIES AT 1S LATER THAN (CNV)/ DEPARTS AT (CNV)--
READY
WHAT (FLIES FROM NEW YORK,) CONNECTS WITH / FLIES TO SANTA BARBARA*1--
I FIGURE SANTA BARBARA FLIES TO (CNV)/ CONNECTS WITH (CNV)/(FLIES FROM NEW YORK, ) FLT. 207
WHAT NOW?
OK--
ANSWER - FLT. 207
READY
FLT. 207 USED ON (CNV)/ IS LARGE AIRPLANE*--
I FIGURE FLT. 207 USED ON (CNV)/ IS LARGE/ JET PLANE AND AIRPLANE
WHAT NOW?
OK--
ANSWER - TRUE
READY
Fig. 5b Application to census, airplane and airline flight data with an ambiguous relation.
a  TRUE IMPLY IMPLIES(CNV)/TRUE--
    READY
    FALSE IMPLY IMPLIES/FALSE--
    READY
    A IMPLIES B--
    READY
    A TRUE PROPOSITION--
    READY
    B TRUE PROPOSITION*1--
    I FIGURE B TRUE/ PROPOSITION
    WHAT NOW?
    OK--
    WAIT...
    ANSWER - TRUE
    READY
    FOLLOWS FROM IMPLY IMPLIES(CNV)--
    READY

b  IMPLIES IMPLY IMPLIES/IMPLIES--
    READY
    C IMPLIES D--
    READY
    D FALSE PROPOSITION--
    READY
    C FALSE PROPOSITION*1--
    I FIGURE C FALSE/ PROPOSITION
    WHAT NOW?
    OK--
    WAIT...
    ANSWER - TRUE
    READY
    FORGET PROPOSITIONB--
    FORGET - PROPOSITION
    OK?...YES--
    WAIT...
    READY
    B IMPLIES C--
    READY
    D FALSE PROPOSITION--
    READY
    A FALSE PROPOSITION*1--
    I FIGURE A FALSE/ UNKNOWN
    WHAT NOW?
    2--
    I FIGURE A FALSE/ PROPOSITION
    WHAT NOW?
    OK--
    WAIT...
    ANSWER - TRUE

Fig. 6a  A Simple Problem in Logic Using Modus Ponens
READY
ELEMENT IMPLBY ELEMENT/SUBSET--
READY
S1 SUBSET S2--
READY
S2 SUBSET S3--
READY
E1 ELEMENT S1--
READY
WHAT IS QUESTION--
READY
E1 ELEMENT WHAT*1--
I FIGURE E1 ELEMENT/ S1 AND S2
WHAT NOW?
1--
I FIGURE E1 ELEMENT/ S1 AND S2 AND S3
WHAT NOW?
OK--
WAIT...
ANSWER - S1
    AND S2
    AND S3
READY
D FOLLOWS, FROM A*2--
I FIGURE D FOLLOWS, FROM/ C AND B
WHAT NOW?
1--
I FIGURE D FOLLOWS, FROM/ C AND B AND A
WHAT NOW?
OK--
WAIT...
ANSWER - TRUE [Answer to question posed at line "c".]
READY
F IS AXIOM--
READY
FIND, PROOF, FROM IS QUESTION--
READY
A DERIVABLE, FROM F--
READY
D FIND, PROOF, FROM AXIOM*6--
Enter comma list of relation classes to be used or - " ANY "
ANY--
D FOLLOWS, FROM A DERIVABLE, FROM F IS AXIOM
WHAT NOW?
OK--
READY

Fig. 6b Set Theory and Simple Theorem Proving
1 PLUS.1.IS 2--
READY
2 PLUS.1.IS 3--
READY
3 PLUS.1.IS 4--
READY
PLUS.2.IS IMPLY PLUS.1.IS/PLUS.1.IS--
READY
WHAT IS QUESTION--
READY
2 PLUS.2.IS WHAT*1--
1 FIGURE 2 PLUS.2.IS/ 4
WHAT NOW?
OK--
WAIT...
ANSWER - 4
READY
MINUS.2.IS IMPLY PLUS.2.IS(CNV)--
READY
 a 1 DIVISIBLE.BY,2 FALSE--
READY
 b 2 DIVISIBLE.BY,2 TRUE--
READY
c DIVISIBLE.BY,2 IMPLY MINUS.2.IS/DIVISIBLE.BY,2--
READY
4 DIVISIBLE.BY,2 WHAT*2--
1 FIGURE 4 DIVISIBLE.BY,2/ TRUE
WHAT NOW?
OK--
WAIT...
ANSWER - TRUE
READY
3 DIVISIBLE.BY,2 WHAT*3--
1 FIGURE 3 DIVISIBLE.BY,2/ FALSE
WHAT NOW?
OK--
WAIT...
ANSWER - FALSE
READY

Fig. 7 Arithmetic and Handling Recursive Definitions
3M3CL 1C 0M1CR--
READY
3M3CL 1M1C 1M1CR--
READY
3M3CL 2C 0M2CR--
READY
3M3CL 1C 0M2CR--
READY
3M2CL 1M 1M1CR--
READY
3M2CL 2C 0M3CR--
READY
3M1CL 1C 0M3CR--
READY
3M1CL 2M 2M2CR--
READY
2M2CL 1M1C 2MCR--
READY
2M2CL 2M 3M1CR--
READY
1M1CL 1M 3M2CR--
READY
1M1CL 1M1C 3M3CR--
READY
0M3CL 1C 3M1CR--
READY
0M3CL 2C 3M2CR--
READY
0M2CL 1C 3M2CR--
READY
0M2CL 2C 3M3CR--
READY
0M1CL 1C 3M3CR--
READY
GET TO IS QUESTION--
READY
3M3CL GET TO 3M3CR*11--
Enter comma list of relation classes to be used or - " ANY "
ANY-- [First solution follows:]
  3M3CL 2C 0M2CR 1C(CNV) 3M2CL 2C 0M3CR 1C(CNV) 3M1CL 2M 2M2CR
  1M1C(CNV) 2M2CL 2M 3M1CR 1C(CNV) 0M3CL 2C 3M2CR 1M(CNV) 1M1CL
  1M1C 3M3CR [Second solution follows:]
  3M3CL 2C 0M2CR 1C(CNV) 3M2CL 2C 0M3CR 1C(CNV) 3M1CL 2M 2M2CR
  1M1C(CNV) 2M2CL 2M 3M1CR 1C(CNV) 0M3CL 2C 3M2CR 1C(CNV) 0M2CL 2C
  3M3CR
WHAT NOW?
OK--
READY

Fig. 8 Solving the Missionary-Cannibal Problem
In attempting to answer a question, it may be necessary to supply further information to the system. Such a situation is illustrated by line "b" in the conversation starting at line "a".

Notice that in the case of a "What R Y" question, SAMENLAQ II proceeds by attempting to apply the R(CNV) relation to the node Y.

Figure 5 shows an application of SAMENLAQ II to census, airplane and airline flight data. Note especially that even though the relation "IS_LARGE" is defined ambiguously in lines "a" and "b" as to its application to cities or airplanes, SAMENLAQ II can disambiguate it from context (lines "c" and "d").

In figure 6, SAMENLAQ II works with simple logic and set theory. The rule of inference Modus Ponens is entered in 6a line "a", and with this and the transitivity of IMPLIES given in line "b", simple "chain implication" problems can be solved. In 6b the relations SUBSET and ELEMENT are introduced along with the rule (x)(x ∈ A & A ⊆ B → x ∈ B). Then some set membership problems are solved, and finally, a simple proof is constructed.

Figure 7 shows SAMENLAQ II being taught its first lesson in arithmetic. Although neither numbers nor arithmetic functions have been built into the SAMENLAQ II structure or executive routines, SAMENLAQ II is capable of being taught arithmetic the way school children use to be taught: by first memorizing tables, and then being taught certain rules. Notice especially that divisibility by 2 was defined recursively in lines "a", "b" and "c". Similarly, SAMENLAQ II could have been taught multiplication, division, the recursive definitions for less than and greater than as well as other arithmetic relations.

Figure 8 shows SAMENLAQ II solving the Missionary - Cannibal Problem, with three missionaries, three cannibals and a boat that holds a maximum of two people. The problem was described to SAMENLAQ II as a set of all the legal states in the problem with all the possible transitions between the states. For example, the first line represents the fact that if 3 missionaries and 3 cannibals are on the left bank with the boat on the left bank, then 1 cannibal can take the boat to the right bank, which will result in there being 0 missionaries and 1 cannibal on the right bank with the boat on the right bank. SAMENLAQ II solves the problem by showing how the boat should be used to get from the initial state, 3M3CL, to the final state, 3M3CR. Any problem solving task that can be represented as finding a path from an initial state to a final state through a state transition graph can in theory be solved similarly by SAMENLAQ II.*

5. Extension of the SAMENLAQ Structure

Work is now proceeding on the design and implementation of a memory net structure, MENS, which goes further than SAMENLAQ II toward satisfying goals discussed in the first section of this paper. The two major improvements needed in the SAMENLAQ structure are:
1. the ability to deal with a name which is
   itself a statement
2. the ability to store names which represent some unspecified other names, i.e. act as variables.

The first would facilitate the handling of n-ary relations and statements which serve to modify or give further information about other statements. The second would allow generalizations to be stored, and would also permit the storage of statements of the predicate calculus directly in the memory structure. These statements could then be interpreted by the executive and used as rules of inference to direct the memory search routines in a manner similar to the way SAMENLAQ II deals with IMPLY definitions.

The currently extant implementation of MENS (which is programmed in Burroughs Extended ALGOL and uses SLIP, a SLIP like package of list processing routines) incorporates the first improvement, and work is progressing on the design of the implementation of the second improvement.

The main generalization involved in going from SAMENLAQ to MENS was to let xRy statements be nodes in the net along with arguments and relations. The basic element of the MENS structure is called an item, which may be an unstructured unit or may be a structure consisting of a pair or triple of items. Thus, items are similar to the "events" used by Simmons et al in Prososyntax II. As in SAMENLAQ, a major characteristic of the MENS structure is that there is no duplication of items or structures; the physically same item is used everywhere that the structure it represents is referred to in a containing structure. Several implications of this uniqueness of items are: (1) two structures which have a substructure in common actually overlap in the net, (2) if there is an item representing logical implication, all structures interpretable as rules of inference will be discoverable directly from that item since it will be a central substructure of all of them, (3) in

* Since SAMENLAQ II exhaustively searches the state transition graph, the threat of exponential growth is ever present. Thus in certain interesting problems exhaustive search would be infeasible and heuristic search techniques would be necessary.
general statements involving a quantified variable, the separate occurrences of the variable will all be pointers to a single item, so that a substitution attached to that item will serve as a substitution for all occurrences of the variable.

Allowing a structure to be formed from a pair of substructures provides for the representation of unary relations such as negation and quantification. Allowing a structure to be formed from a triple of substructures provides for the representation of binary relation, and since any of the substructures may in fact be structures as well as unstructured items provides for the representation of n-ary relations. A more direct representation of n-ary relations is provided in another version of MENS being implemented*, which will allow any structure to consist of any number of substructures.

It is demonstrated in the examples given in Section 4 that SAMENLAQ II is capable of answering questions in formal logic involving simple chains of inference and basic set theory. With its ability to store statements of the predicate calculus, the MENS structure should enable simple question answering routines to perform more complicated theorem proving. Although we should perhaps not expect a high powered theorem prover to be developed in this way, the MENS extension of SAMENLAQ will provide an interesting contrast to systems, such as Green and Raphael's QA2*, which use theorem proving techniques to answer questions.

References


* A paper describing this version of MENS and the natural language question answering system which will use it is forthcoming from the RAND Corporation as a RAND Memorandum by Martin Kay, Ronald M. Kaplan, and Stuart C. Shapiro.
APPENDIX - MEMORY STRUCTURE

READY

DUMP MEMORY ON TELETYPewriter-

THE MEMORY IS ---

\[ 2 = (\text{DIV.BY.2} - /20)(\text{PLUS.2.IS} - /13)(\text{PLUS.1.IS} - /4)(\text{PLUS.1.IS(INV)} - /1) \]

/1 = 1,
/2 = 2,
/3 = (\text{DIV.BY.2} / 25)(\text{PLUS.1.IS} - /6)(\text{PLUS.1.IS(INV)} - /3)
/3 = 2,
/4 = 3,
/4 = (\text{DIV.BY.2} / 24)(\text{PLUS.2.IS(INV)} - /12)(\text{PLUS.1.IS(INV)} - /5)
/5 = 3,
/6 = 4,
\text{PLUS.1.IS} / \text{PLUS.1.IS} = (\text{IMPLY(INV)} - /7)
/7 = \text{PLUS.2.IS},
\text{PLUS.2.IS(INV)} = (\text{IMPLY(INV)} - /14)(\text{IMPLY} - /8)
/8 = \text{PLUS.1.IS(INV)} / \text{PLUS.1.IS(INV)},
\text{PLUS.2.IS} = (\text{IMPLY} - /9)
/9 = \text{PLUS.1.IS} / \text{PLUS.1.IS},
\text{QUESTION} = (\text{IS(INV)} - /10)
/10 = \text{WHAT},
\text{WHAT} = (\text{IS} - /11)
/11 = \text{QUESTION},
/12 = 2,
/13 = 4,
/14 = \text{MINUS.2.IS},
\text{MINUS.2.IS(INV)} = (\text{IMPLY} - /15)
/15 = \text{PLUS.2.IS},
\text{MINUS.2.IS} = (\text{IMPLY} - /16)
/16 = \text{PLUS.2.IS(INV)},
\text{FALSE} = (\text{DIV.BY.2(INV)} - /17)
/17 = 1,3
/18 = \text{FALSE},
\text{TRUE} = (\text{DIV.BY.2(INV)} - /19)
/19 = 2,4
/20 = \text{TRUE},
\text{MINUS.2.IS} / \text{DIV.BY.2} = (\text{IMPLY(INV)} - /21)
/21 = \text{DIV.BY.2},
\text{DIV.BY.2(INV)} = (\text{IMPLY} - /22)
/22 = \text{DIV.BY.2(INV)} / \text{MINUS.2.IS(INV)},
\text{DIV.BY.2} = (\text{IMPLY} - /23)
/23 = \text{MINUS.2.IS} / \text{DIV.BY.2},
/24 = \text{TRUE},
/25 = \text{FALSE},
\text{READY}

Fig. 9 Shows SAMENLAQ II's actual memory structure after its arithmetic lesson, which was shown in figure 7. The underlined material was learned after being discovered as implicit information while answering questions.