CSE 250 Data Structures

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Lec 19: Adjacency Lists and Matrices

Announcements

• WA3 due Sunday

Edge List Summary

<u>Graph</u>

vertices: LinkedList<Vertex> edges: LinkedList<Edge>

<u>Vertex</u>

Storing the list nodes in the edges/vertices allows us to remove by reference in $\Theta(1)$ time

<u>Edge</u>

label: T vertex: origin vertex: destination node: LinkedListNode

Edge List Summary

- addEdge, addVertex: O(1)
- removeEdge: O(1)
- removeVertex: O(m)
- vertex.incidentEdges: O(m)
- vertex.edgeTo: O(m)
- Space Used: *O*(*n*) + *O*(*m*)

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- removeVertex: O(m)
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 Involves checking every edge in the graph

• Space Used: O(n) + O(m)

How can we improve?

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Idea: Store the in/out edges for each vertex!

(Called an adjacency list)

```
1 public class Vertex<V,E> {
2    public Node<Vertex> node;
3    public List<Edge> inEdges = new CustomLinkedList<Edge>();
4    public List<Edge> outEdges = new CustomLinkedList<Edge>();
5    /*...*/
6 }
```

Each vertex stores a list of **inEdges** and **outEdges**, which are maintained as the graph is modified...

What functions need to change to maintain these lists?

```
1 public Edge addEdge(Vertex orig, Vertex dest, E label) {
2 Edge e = new Edge(orig, dest, label);
3 e.node = edges.add(e);
4 orig.outEdges.add(e);
5 dest.inEdges.add(e);
6 return e;
7 }
```

What is the complexity of addEdge now?

```
1 public Edge addEdge(Vertex orig, Vertex dest, E label) {
2 Edge e = new Edge(orig, dest, label);
3 e.node = edges.add(e);
4 orig.outEdges.add(e);
4 dest.inEdges.add(e);
5 dest.inEdges.add(e);
6 return e;
7 }
```

What is the complexity of addEdge now? Still $\Theta(1)$

1	<pre>1 public void removeEdge(Edge edge) {</pre>			
2	<pre>edges.remove(edge.node);</pre>			
3 4	<pre>edge.orig.outEdges.remove(edge); edge.dest.inEdges.remove(edge);</pre>	← When we remove an edge from the graph, also remove it from the adjacency lists		
5	}	also remove it norm the dajasensy lists		

What is the complexity of **removeEdge** now?

1	<pre>1 public void removeEdge(Edge edge) {</pre>		
2	edges.remove(edge.node);		
3 4	edge.orig.outEdges.remove(edge); ← When we remove an edge from the graph, edge.dest.inEdges.remove(edge); also remove it from the adjacency lists		
5	}		

What is the complexity of removeEdge now? O(deg(orig) + deg(dest)) :(

But how can we fix this?

1	<pre>public class Edge<v,e> {</v,e></pre>
2	<pre>public Node<edge> node;</edge></pre>
3	<pre>public Node<edge> inNode;</edge></pre>
4	<pre>public Node<edge> outNode;</edge></pre>
5	/**/
6	}

Each Edge now also stores a reference to the nodes in each adjacency list

1 public Edge addEdge(Vertex orig, Vertex dest, E label) {
2 Edge e = new Edge(orig, dest, label);
3 e.node = edges.add(e);
4 e.outNode = orig.outEdges.add(e);
5 e.inNode = dest.inEdges.add(e);
6 return e;
7 }

What is the complexity of addEdge now? Still $\Theta(1)$

1	<pre>public void removeEdge(Edge edge) {</pre>
2	edges.remove(edge.node);
3 4	edge.orig.outEdges.remove(edge.outNode); edge.dest.inEdges.remove(edge.inNode); ← When we remove an edge from the graph, also remove it from the
5	} adjacency lists (remove by reference)

What is the complexity of **removeEdge** now?

1	<pre>public void removeEdge(Edge edge) {</pre>
2	edges.remove(edge.node);
3 4	edge.orig.outEdges.remove(edge.outNode); edge.dest.inEdges.remove(edge.inNode); ← When we remove an edge from the graph, also remove it from the
5	adjacency lists (remove by reference)

What is the complexity of **removeEdge** now? $\Theta(1)$

So, we are able to store and maintain adjacency lists in each vertex while still keeping a $\Theta(1)$ runtime for addVertex, addEdge, and removeEdge

How much extra space is used?

So, we are able to store and maintain adjacency lists in each vertex while still keeping a $\Theta(1)$ runtime for addVertex, addEdge, and removeEdge

How much extra space is used? $\Theta(1)$ per edge

Each edge only appears in 3 lists:

- The edge list
- One vertices inList
- One vertices outList

So, we are able to store and maintain adjacency lists in each vertex while still keeping a $\Theta(1)$ runtime for addVertex, addEdge, and removeEdge

How much extra space is used? $\Theta(1)$ per edge

Each edge only appears in 3 lists:

- The edge list
- One vertices inList
- One vertices outList

But now what have we gained?

```
1 public void removeVertex(Vertex v) {
2 for(edge : v.getIncidentEdges()) {
3 removeEdge(edge.node)
4 }
5 vertices.remove(v.node);
6 }
```

What is the complexity of **removeVertex** now?

```
1 public void removeVertex(Vertex v) {
2 for(edge : v.getIncidentEdges()) {
3 0(1)
4 }
5 0(1)
6 }
```

What is the complexity of **removeVertex** now?

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What is the complexity of **removeVertex** now? $\Theta(deg(v))$

Adjacency List Summary

Graph Vertex vertices: LinkedList[Vertex] label: LinkedListNode edges: LinkedList[Edge] node: inEdges: LinkedList[Edge] outEdges: LinkedList[Edge] Storing the list of incident edges in Edge the vertex saves us the time of checking every edge in the graph. label: node: LinkedListNode

inNode:

outNode:

LinkedListNode

LinkedListNode

The edge now stores additional nodes to ensure removal is still $\Theta(1)$

Adjacency List Summary

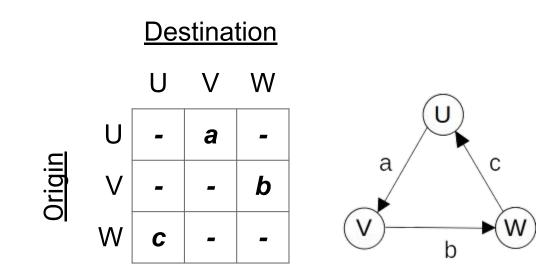
- addEdge, addVertex: $\Theta(1)$
- removeEdge: $\Theta(1)$
- removeVertex: $\Theta(deg(vertex))$
- vertex.incidentEdges: $\Theta(deg(vertex))$
- vertex.edgeTo: Θ(deg(vertex))
- Space Used: $\Theta(n) + \Theta(m)$

Adjacency List Summary

- addEdge, addVertex: $\Theta(1)$
- removeEdge: $\Theta(1)$
- removeVertex: $\Theta(deg(vertex))$
- vertex.incidentEdges: @(deg(vertex))
- vertex.edgeTo: Θ(deg(vertex))
- Space Used: $\Theta(n) + \Theta(m)$

Now we already know what edges are incident without having to check them all

Adjacency Matrix



- addEdge, removeEdge:
- addVertex, removeVertex:
- vertex.incidentEdges:
- vertex.edgeTo:
- Space Used:

Just change a single entry of the matrix

- addEdge, removeEdge: ⊖(1)
- addVertex, removeVertex:
- vertex.incidentEdges:
- vertex.edgeTo:
- Space Used:

- addEdge, removeEdge: $\Theta(1)$
- addVertex, removeVertex: $\Theta(n^2)$
- vertex.incidentEdges:
- vertex.edgeTo:
- Space Used:

Resize and copy the whole matrix

- addEdge, removeEdge: Θ(1)
- addVertex, removeVertex: $\Theta(n^2)$
- vertex.incidentEdges: $\Theta(n)$
- vertex.edgeTo:
- Space Used:

Check the row and column for that vertex

- addEdge, removeEdge: ⊖(1)
- addVertex, removeVertex: $\Theta(n^2)$
- vertex.incidentEdges: $\Theta(n)$
- vertex.edgeTo: $\Theta(1)$
- Space Used:

Check a single entry of the matrix

- addEdge, removeEdge: ⊖(1)
- addVertex, removeVertex: $\Theta(n^2)$
- vertex.incidentEdges: $\Theta(n)$
- vertex.edgeTo: $\Theta(1)$
- Space Used: $\Theta(n^2)$

How does this relate to space of edge/adjacency lists?

- addEdge, removeEdge: ⊖(1)
- addVertex, removeVertex: $\Theta(n^2)$
- vertex.incidentEdges: $\Theta(n)$
- vertex.edgeTo: $\Theta(1)$
- Space Used: $\Theta(n^2)$

How does this relate to space of

edge/adjacency lists? If the matrix is "dense" it's about the same

So...what do we do with our graphs? ...next lecture