

# CSE 250 Recitation

March 4-5 - Recursion



# Binary Search

The binary search algorithm let's us effective search a List

To work correctly and efficiently the List must:

- Be sorted
- Allow constant time random access (ie an Array)

It works by comparing our target to the midpoint, then searching **only** the left half or the right half

## Binary Search Code

```
1 int binarySearch(ArrayList<T> list, T target) {
2     return binarySearch(list, target, 0, list.size() - 1);
3 }
4
5 int binarySearch(ArrayList<T> list, T target, int start, int end) {
6     if(start == end) { return start; }
7     int mid = (start + end) / 2;
8     T guess = list.get(mid);
9     if(guess.equals(target)){ return mid; } // We found our target!
10    else if(target.compareTo(guess) < 0) { // Target is in the left half
11        return binarySearch(list, target, start, mid);
12    } else { // Target is in the right half
13        return binarySearch(list, target, mid+1, end);
14    }
15 }
```

# Runtime Growth Function

$$T(N) = \begin{cases} T\left(\frac{N}{2}\right) + \theta(1) & \text{if target is not found} \\ \theta(1) & \text{else} \end{cases}$$

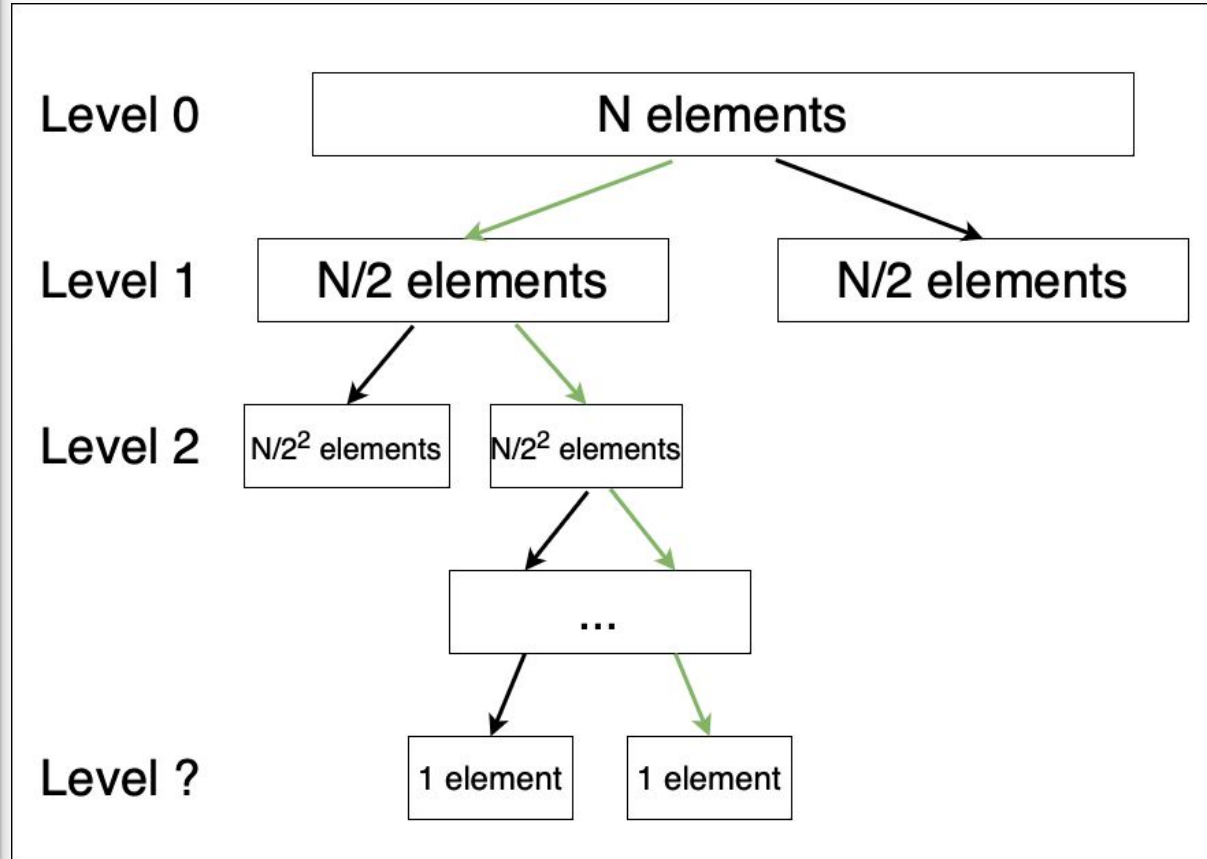
# Runtime Growth Function

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**Exercise:** Draw the recursion tree for this growth function

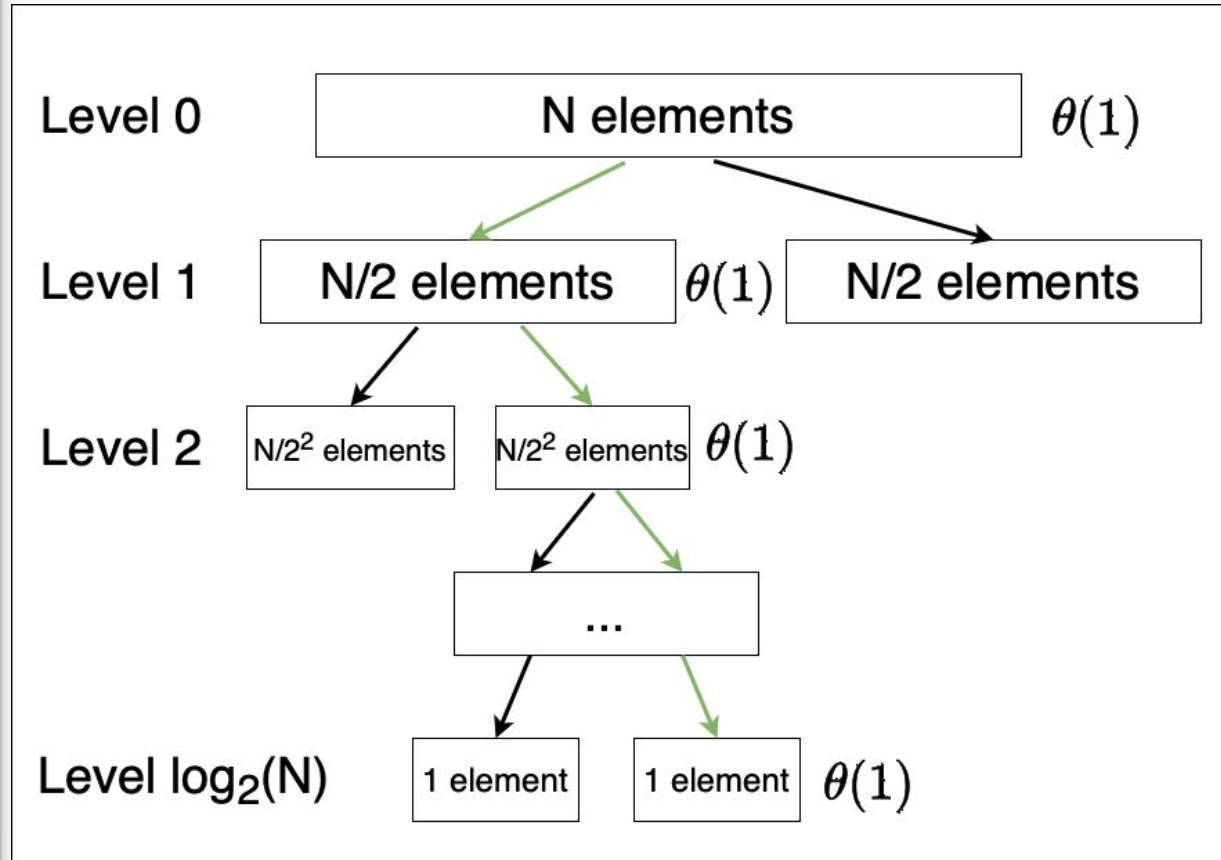
# Hypothesis

- To form your hypothesis for the runtime of the algorithm, you should ask yourself two questions:
  - How much work are you doing on each level?
  - How many levels are there?



# Hypothesis

- We are doing a constant amount of work on each level
- We are guaranteed to have a maximum of  $\log_2(N)$  levels



# Hypothesis Summation

$$O(\log_2(N)) \sum_{i=1} \theta(1)$$



## Inductive Hypothesis

$$T(N) \in O(\log_2(N))$$

Base Case

$$T(1) \leq c \cdot \log_2(1)$$

Base Case

$$T(\overset{2}{\cancel{1}}) \leq c \cdot \log_2(\overset{2}{\cancel{1}})$$

## Inductive Case

Assume:  $T\left(\frac{N}{2}\right) \leq c \cdot \log_2\left(\frac{N}{2}\right)$

Show:  $T(N) \leq c \cdot \log_2(N)$