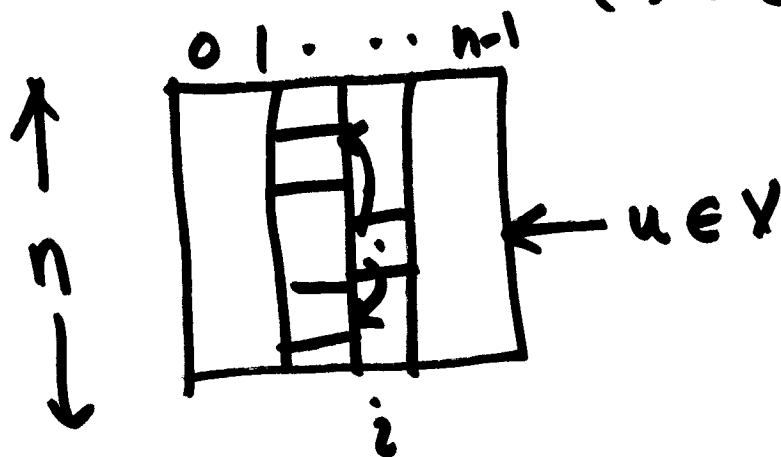


DECEMBER 2, 2011

$$\text{OPT}(i, u) = \min \{ \text{OPT}(i-1, u),$$

$$\min_{(u, w) \in E} \text{OPT}(i-1, w) + c_{uw} \}$$



Column i
only depends
on column
 $i-1$

$$M[u, i] (= \text{OPT}(i, u)) \leftarrow (\text{prove by induction})$$

Bellman-Ford $(G = (V, E), \{c_e\}, t)$

$$1. M[t, 0] = 0, M[u, 0] = \infty \quad \forall u \neq t$$

$\leftarrow O(n)$

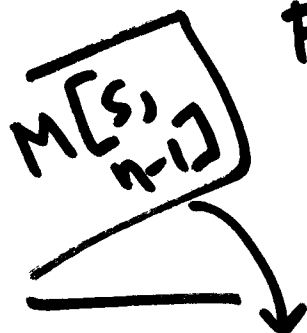
$$2. \text{For } i = 1 \dots n-1 \leftarrow n$$

$$\text{For } u \in V \leftarrow n$$

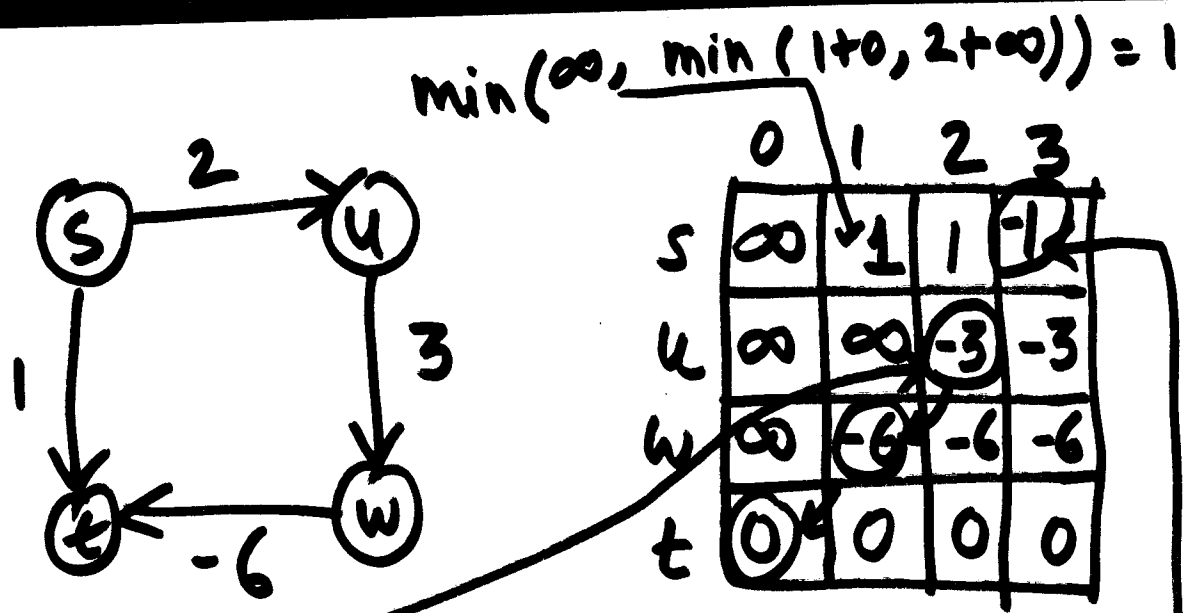
$$M[u, i] = \min \{ M[u, i-1], \min_{(u, w) \in E} c_{uw} + M[w, i-1] \}$$

$O(nu)$

" $O(n)$ "



Q: How to figure out least cost of shortest
 $s-t$ path (given M)



Shortest
s-t
path

$$\min(\infty, 3 + (-6))$$

$$\min(1, \min(1+0, 2+(-3)))$$



Run time:

$$\rightarrow O(n^3)$$

$$\rightarrow O(n \cdot \sum_{u \in V} n_u) = O(nm)$$

$\nwarrow O(m)$

Ex: Algo to reconstruct shortest path

Q: Space: $O(n^2)$. Can you reduce this?

$O(n)$ space. Obs: i^{th} column only depends on $i-1^{\text{th}}$ column

Longest path problem

Simple

Read
[KT, 1.8]