

Q_j : optimal soln for $\{1, \dots, j\}$

$$\text{OPT}(j) = \text{value of } Q_j = \sum_{i \in Q_j} u_i$$

Case 1: $j \in Q_j$

$$Q_j = j; Q_{p(j)} \quad \text{OPT}(j) = u_j + \text{OPT}(p(j))$$

Case 2: $j \notin Q_j$

$$Q_j = Q_{j-1} \quad \text{OPT}(j) = \text{OPT}(j-1)$$

$$\text{OPT}(j) = \max \left\{ \overset{\textcircled{1}}{u_j + \text{OPT}(p(j))}, \overset{\textcircled{2}}{\text{OPT}(j-1)} \right\}$$

given $\text{OPT}(1), \text{OPT}(2), \dots, \text{OPT}(j-1)$.

(By defn. $p(j) \leq j-1$)

→ you know $\text{OPT}(p(j))$ & $\text{OPT}(j-1)$

$$\left. \begin{array}{l} \textcircled{1} \geq \textcircled{2} \Rightarrow j \in Q_j \\ \textcircled{1} < \textcircled{2} \Rightarrow j \notin Q_j \end{array} \right\} \begin{array}{l} \text{Ex: given } \text{OPT}(1) \dots \text{OPT}(j) \\ \Rightarrow \text{can compute } Q_j \end{array}$$

Claim: If we compute $\text{OPT}(1) \dots, \text{OPT}(n)$

→ we can compute in $O(n)$ time $Q = Q_n$.