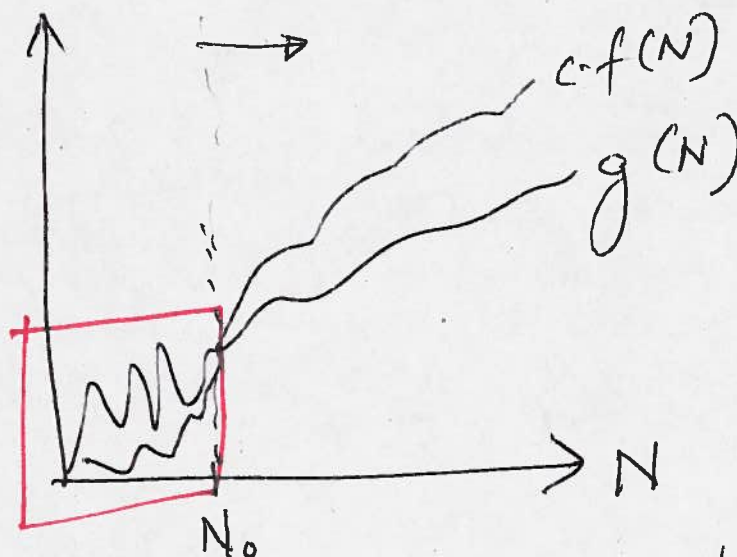


Def: $g(N)$ is $O(f(N))$
(or g is $O(f)$)

if $\exists$ a $N_0 \geq 0$ s.t.
& a constant c

$\forall N \geq N_0$

$$g(N) \leq c \cdot f(N)$$



Q: $pN^2 + qN + r \leq r$
 $O(N^2) \leq qN^2$

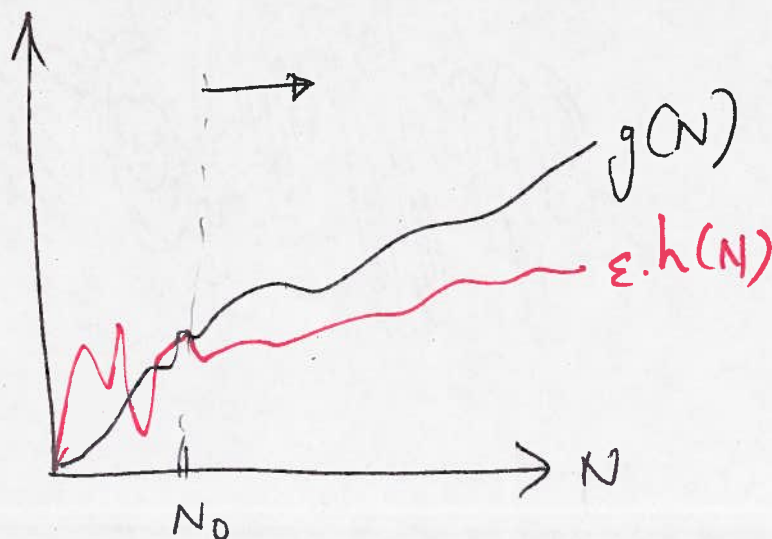
$\leq (p+q+r)N^2$
 $\hookrightarrow O(N^3) \checkmark$
 $\hookrightarrow O(N) \times$

Def: $g(N)$ is $\Omega(h(N))$
(g is $\Omega(h)$)

if $\exists N_0 \geq 0$
& $\epsilon > 0$ s.t.

$\forall N \geq N_0$

$$g(N) \geq \epsilon \cdot h(N)$$



$pN^2 + qN + r$
 $\Omega(N^2)$

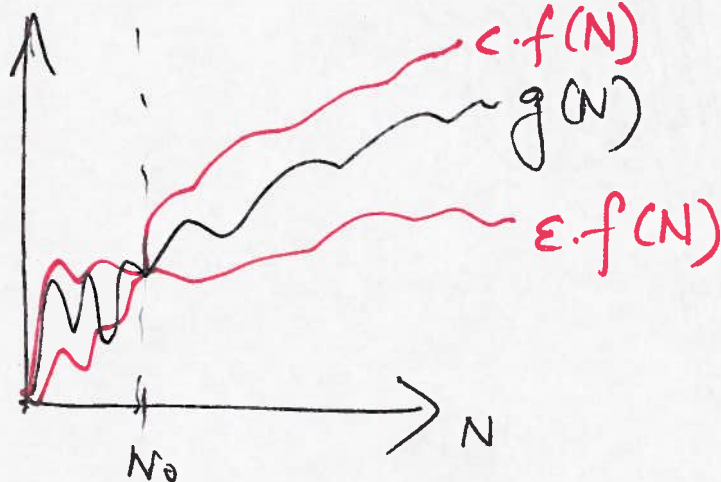
$\epsilon = p, N_0 \geq 0$

$\Omega(N^3) \times$

$\Omega(N) \checkmark$

Def: $g(N)$ is $\Theta(f(N))$

if $g(N)$ is both $O(f(N))$
and $\Omega(f(N))$



$pN^2 + qN + r$
 $\rightarrow \Theta(N^2)$
 $\rightarrow \Theta(N^3) \times$
 $\rightarrow \Theta(N) \times$

Useful properties

① Transitivity g is $O(h)$ & h is $O(f) \Rightarrow g$ is $O(f)$

Similarly, Ω, Θ

② Sum of functions.

g is $O(h)$ & f is $O(h)$, $g+f$ is $O(h)$

$(f_1, \dots, f_k \text{ are } O(h)) \quad f_1 + \dots + f_k \text{ is } O(h)$
 $k \rightarrow \text{is constant}$

also true for Ω & Θ

③ g is $O(h_1)$, f is $O(h_2)$
 $g \cdot f$ is $O(h_1 \cdot h_2)$

$T(n) = \max_{\substack{\text{all inputs } I \\ \text{of size } n}} \text{steps taken by the algo on } I.$

Case 1: $T(n) \leq U(n)$ [$O(U)$]

Show: $\forall$ inputs I of size n ,
steps $\leq U$. [$O(U)$]

$$(\equiv \max_{i=1..m} a_i \leq U)$$

$$\uparrow \uparrow$$
$$\forall i \ a_i \leq U)$$

Case 2: $T(n) \geq L(n)$ [$\Omega(L)$]

$\exists$ an input I of size n
steps $\geq L$ [$\Omega(L)$]

$$(\equiv \exists i \text{ s.t. } a_i \geq L)$$

FIND ($a_1, \dots, a_n; v$) /# Output: i if $a_i = v$
 -1 if no such $i \exists$.

runs $\leq n$
 ① $\rightarrow$ for ($i = 1 \dots n$) $\sim$ #itr
 ② if ($a_i == v$) return i . $\leftarrow O(1)$
 ③ return -1 . $\leftarrow O(1)$

$\left. \begin{array}{l} \text{②} \\ \text{③} \end{array} \right\} \leq c \cdot \#itr$
 $O(n \cdot 1) = O(n)$

$$T(n) = O(n)$$

$$\text{step ①} + \text{②} \rightarrow O(n)$$

$$\text{step ③} \rightarrow O(1) \leq O(n)$$

$$T(n) = (\text{step ①} + \text{②}) + \text{step ③}$$

$$= O(n) + O(1)$$

$$= O(n)$$

true for all
 inputs of size
 $(n+1)$

$$T(n) = \Omega(n)$$

for any $n \geq 1$

cons. $a_1 = n$

$a_2 = n-1$

$\vdots$

$a_n = 1$

$v = 1$