

Tue Jan 27

[Continue from 3rd slide of overhead notes]

④

Sets and Logic

$$197 + 48 = 245$$

$$197 + 48 + 48 = 293$$

$$197 + 48 + 48 + 48 = 341$$

Introducing - - -

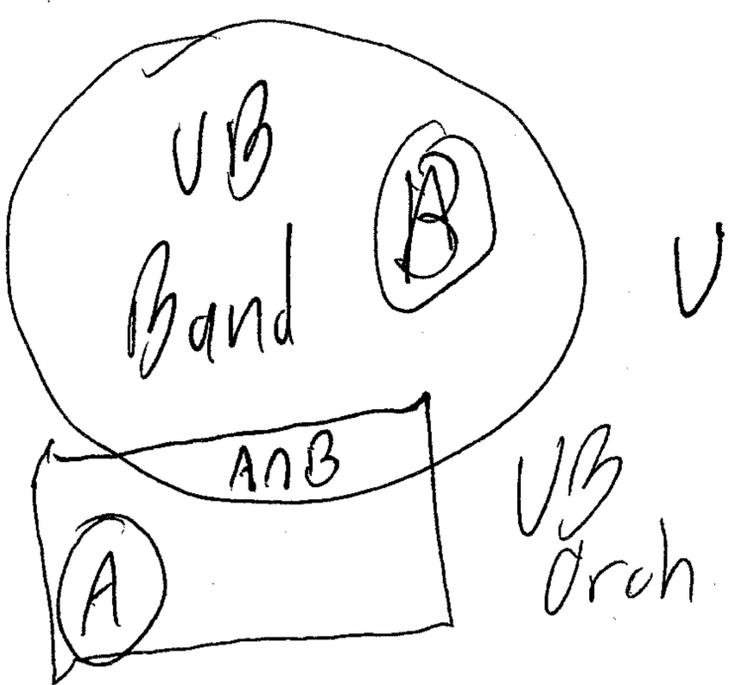
The UB Band + The Union College Band

Person X

$X \in \text{UB Band}$ OR $X \in \text{Union Band}$



$X \in \text{UB}$ \cup $X \in \text{Union}$ \cup $X \in \text{Union}$



Union Coll Band

$\sim B$ \bar{B}
Complement of set.
 \bar{B} B'

In Orch, Not in Band.

$$= A \cap (\sim B)$$

$$A - B$$

$A \setminus B$ "Backslash Set Minus"



People in one but not the other ⁽²⁾

$$= (A \setminus B) \cup (B \setminus A)$$

Alternative Notations: $A \oplus B$, $A \Delta B$

$$A \Delta B \Delta B \stackrel{?}{=} A$$

$$A \oplus B \oplus B = A$$

6 people
 4 in orch
 3 in Band
 (1 in neither)

	1	2	3	4	5	6
^A Orch =	1	1	0	1	0	1
Band	1	0	0	1	1	0

Orch = {1, 2, 4, 6}

Band = {1, 4, 5}

neither = 3

~~Band~~ ~~Orch~~

~~Band~~ ~~A orch~~

Orch Δ Band

" $A \Delta B \Delta B$

= A

Band Orch	1	1	0	1	0	1
Band A orch	1	0	0	1	1	0
Orch Δ Band	0	1	0	0	1	1
" $A \Delta B \Delta B$	1	0	0	1	1	0
= A	1	1	0	1	0	1

Max

Bitwise And = Bitwise Min

Bitwise XOR

Vector

Addⁿ

Mod 2

= orch Again

Continue from Thu 1/29 overhead notes --

(3)

$$\Sigma = \{0, 1\} : \text{object of type set(char) = Alph.}$$

$$\text{Sigma} = \{ '0', '1' \} \quad \begin{array}{l} '0' : \text{char} \\ "0" : \text{string} \end{array}$$

$$\text{Bin} = \mathcal{B} = \{ "0", "1" \} : \text{object of type set(string)}$$

$$\mathcal{B}_2 = \{ "00", "01", "10", "11" \} \quad \text{Empty string: } \epsilon \text{ or } \lambda$$

$$= \{00, 01, 10, 11\} = \Sigma^2 = \{0, 1\}^2$$

$$\mathcal{B}_3 = \{000, 001, 010, 011, 100, 101, 110, 111\} = \Sigma^3 = \{0, 1\}^3$$

The set of all finite binary strings is written Σ^* $\begin{array}{l} * \equiv \text{zero or more.} \\ + \equiv \text{one or more.} \end{array}$

$$\mathcal{B}_2 = \{ \epsilon, 0, 1, 00, 01, 10, 11 \} = \Sigma^{\leq 2}$$

Pairs are any elements (a, b) of a Cartesian Product A x B

$$(a, b) \in A \times B \text{ means } a \in A \text{ \& } b \in B.$$

When $A = B$, we can abbreviate $A \times A$ as A^2 . Thus

$$\Sigma^2 = \Sigma \times \Sigma = \{ (0, 0), (0, 1), (1, 0), (1, 1) \} \text{ collapsed as } \{00, 01, 10, 11\}$$

Triples (a, b, c) belong to $A \times B \times C \equiv (A \times B) \times C \equiv (A \times (B \times C))$

$$A \times A \times A \equiv A^3 \text{ and } \Sigma^{\leq 3} = \Sigma^3 \cup \Sigma^2 \cup \Sigma \cup \{ \epsilon \}$$

Σ^0

Relations: A ^{binary} relation R between sets A and B is any subset of $A \times B$. (4)

Empty relation: \emptyset

Full relation: $A \times B$ iBelf

If $B = A$, we call it a binary relation on A .

Example: $R_{id} = \{(a, a) : a \in A\}$. Does this equal $A \times A$? NO ^{Generally}

$A = \Sigma = \{0, 1\}$. $R_{id} = \{(0, 0), (1, 1)\}$. $A \times A$ includes $\begin{matrix} (0, 1) & (1, 0) \\ 01 & 10 \end{matrix}$

A relation R is a function if for all $a \in A$ there is (exactly) one $b \in B$ such that $(a, b) \in R$. $R_{id} =$ identity function

"at most" defines a partial function whose actual domain can be a subset of A .

Usually: $f: A \rightarrow B$ means

$$\begin{aligned} \text{Dom}(f) &= A \\ \text{Ran}(f) &\subseteq B \end{aligned}$$

But sometimes we lie and allow $\text{Dom}(f) \subsetneq A$.

Then f is a partial function.

Defn: f is a partial function if for all $a \in A$ there is at most one $b \in B$ such that (a, b) belongs to the relation of f .

Is \emptyset a partial fn?
 [from any A to any B]

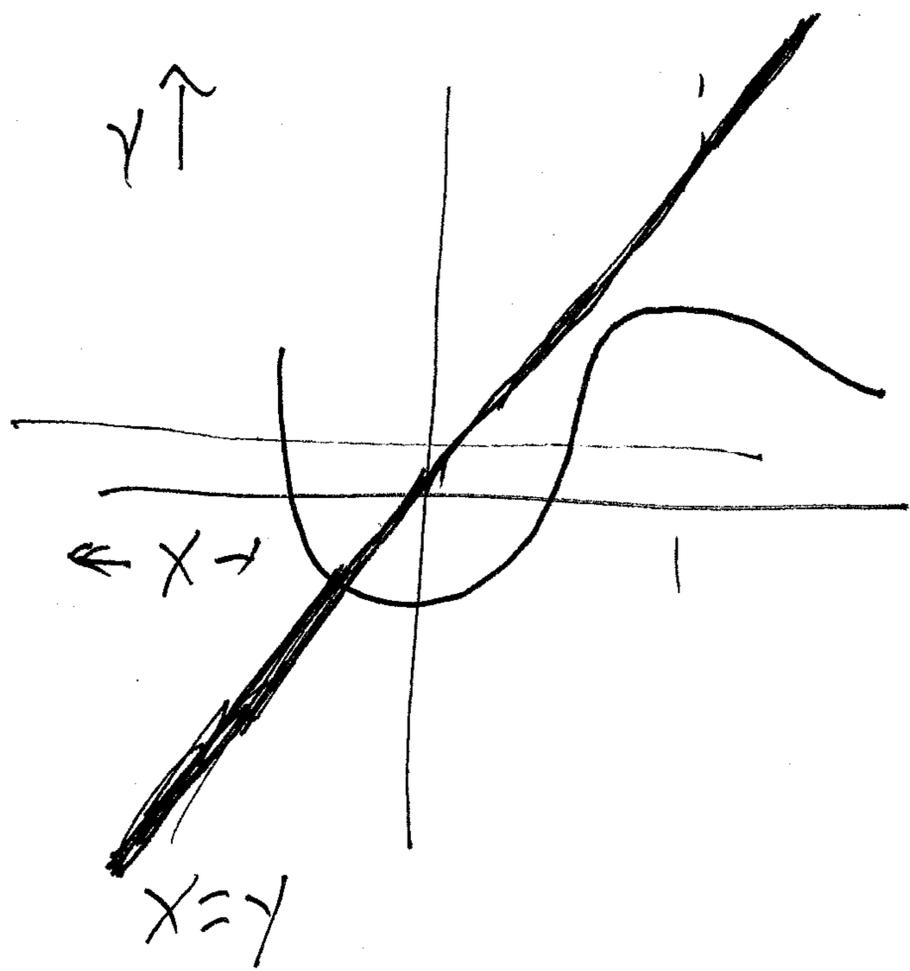
But: $\emptyset: \emptyset \rightarrow \emptyset$ as a proper total function!

R defines a total function if $(\forall a \in A) (\exists b \in B) : (a, b) \in R$.

R is onto its "declared range" if $(\forall b \in B) (\exists a \in A) : (a, b) \in R$.

R is 1-1 if ~~$(\forall a \in A) (\exists \text{at most one } b \in B) : (a, b) \in R$~~

$(\forall b \in B) (\exists \text{at most one } a \in A) : (a, b) \in R$.



Defn: $R \subseteq A \times A$ is an equivalence relation if it is: (5)

Reflexive: $\forall a \in A, (a, a) \in R$.

Symmetric: $\forall a, b \in A, (a, b) \in R \Leftrightarrow (b, a) \in R$.

Transitive: $\forall a, b, c \in A:$
 $(a, b) \in R \wedge (b, c) \in R \rightarrow (a, c) \in R$.

$R(\text{story } v1, \text{story } v2) \equiv$ Versions 1 and 2 are the same up to 95% of the words in either.

$R'(\text{story } 1, \text{story } 2) \equiv$ Stories 1 & 2 differ only in their headlines. at most.

Yes, this is an equivalence relation.

Prove: All horses have the same color.

\equiv All finite sets of horses have the same color.

For all $n \geq 0$, all sets of n horses have the same color.

Basis: ($n=1$) Any set of one horse has just one color.

Ind: Given $n \geq 2$, assume (Ind. Hyp.) all sets of $n-1$ horses have the same color.

Let any set S of horses be given. List S out as $\{H_1, H_2, \dots, H_n\}$

Take $S_1 = \{ \text{first } n-1 \text{ horses} \} = \{ H_1, H_2, \dots, H_{n-1} \}$

$S_2 = \{ \text{last } n-1 \text{ horses} \} = \{ H_2, H_3, \dots, H_n \}$

QED!
 (???)